TECHNISCHE UNIVERSITÄT BERLIN Institut für Mathematik ADM III – Advanced Methods for Integer Linear Programming Summer Term 2007

Prof. Dr. h.c. Martin Grötschel Andreas Bley Benjamin Hiller

# Exercise sheet 11

Deadline: Thursday, July 12th, 2007, 08:30 h in MA-313

### Exercise 1:

#### 4 points

4 points

Let *E* be a finite set and  $E_1, \ldots, E_k$  a partition of *E*. Furthermore, let  $b_1, \ldots, b_k$  be integers satisfying  $1 \le b_i \le |E_k|$ . Then  $(E, \mathcal{I})$  with  $\mathcal{I}$  defined by

$$\mathcal{I} := \{ I \subseteq E \mid |I \cap E_i| \le b_i, 1 \le i \le k \}$$

is a matroid, called the *partitioning matroid*.

Determine the facets of the partitioning matroid based on the characterization of the facets of the matroid polytope given in the lecture.

## Exercise 2:

- a) Let  $b_1, \ldots, b_m$  be positive integers. Show that  $L(a) = L(b_1, \ldots, b_m)$ , where  $a = \text{gcd}(b_1, \ldots, b_m)$  is the greatest common divisor of  $b_1, \ldots, b_m$ .
- b) Give an algorithm that from a set of possibly dependent vectors  $a_1, \ldots, a_m \in \mathbb{Z}^n$  finds a basis  $b_1, \ldots, b_k$  of  $L(a_1, \ldots, a_m)$  such that  $b_i^T e_j = 0$  for j < i, where  $e_j$  is the *j*th unit vector.

*Hint:* Consider one dimension at a time and use the first part of this exercise.

## Exercise 3:

#### 1+2+1 points

A matrix  $A \in \mathbb{Z}^{m \times n}$  of full row rank is said to be in *integer normal form* if it is of the form [B, 0], where  $B \in \mathbb{Z}^{n \times n}$  is invertible and lower triangular. For every matrix  $A \in \mathbb{Z}^{m \times n}$  of full row rank there is a unimodular matrix U such that AU is in integer normal form.

Prove the following theorem:

**Theorem 1** Let  $A \in \mathbb{Z}^{m \times n}$  be a matrix of full row rank and let [B, 0] = AU be the integer normal form of A with a unimodular matrix U. Let  $b \in \mathbb{Z}^m$  and  $\mathcal{F} = \{x \in \mathbb{Z}^n \mid Ax = b\}$ .

- a)  $\mathcal{F}$  is nonempty if and only if  $B^{-1}b \in \mathbb{Z}^m$ .
- b) If  $\mathcal{F} \neq \emptyset$ , every element of  $\mathcal{F}$  is of the form

$$x = U_1 B^{-1} b + U_2 z, z \in \mathbb{Z}^{n-m},$$

where  $U_1$ ,  $U_2$  are submatrices of U such that  $U = [U_1, U_2]$ .

c)  $\mathcal{L} = \{x \in \mathbb{Z}^n \mid Ax = 0\}$  is a lattice and the column vectors of  $U_2$  constitute a basis of  $\mathcal{L}$ .

## Exercise 4:

The theorem of the last exercise suggests a way to solve certain integer programs via alternative bases of lattices. Consider the integer program

$$\max \quad c^T x x \in \mathcal{F} := \{ x \in \mathbb{Z}_{\geq 0}^n \mid Ax = b \},\$$

where A is an integer matrix of full row rank and b and c are integer vectors. Suppose  $x_0$  is an integer point satisfying  $Ax_0 = b$ . Then every  $x \in \mathcal{F}$  can be written as

$$x = x_0 + y$$
, for some  $y \in \mathbb{Z}^n$  s.t.  $Ay = 0, y \ge -x_0$ .

Let  $\mathcal{L} := \{y \in \mathbb{Z}^n \mid Ay = 0\}$  and consider the integer normal form [B, 0] of A obtained using the unimodular matrix U. Due to b) in the theorem, we have

$$\mathcal{L} = \{ y \in \mathbb{Z}^n \mid y = U_2 z, z \in \mathbb{Z}^{n-m} \},\$$

which allows us to reformulate the original IP as

$$\begin{array}{ll} \max & c^T U_2 z \\ & U_2 z \ge -x_0, \\ & z \in \mathbb{Z}^{n-m}. \end{array}$$

Since all bases of  $\mathcal{L}$  can be obtained from  $U_2$  via unimodular matrices, we get an alternative reformulation for any basis B of  $\mathcal{L}$ , namely

$$\begin{array}{ll} \max & c^T B z \\ & B z \geq -x_0, \\ & z \in \mathbb{Z}^{n-m}. \end{array}$$

For a suitable chosen basis the reformulation might be easier to solve than the orginal formulation. Solve the following IP via reformulations based on alternative bases.

$$\begin{array}{ll} \max & x_1 + x_2 + x_3 \\ & 3x_1 + 7x_2 + 10x_3 = 19, \\ & x_1, x_2, x_3 \in \mathbb{Z}_{\geq 0}. \end{array}$$

# 4 points