



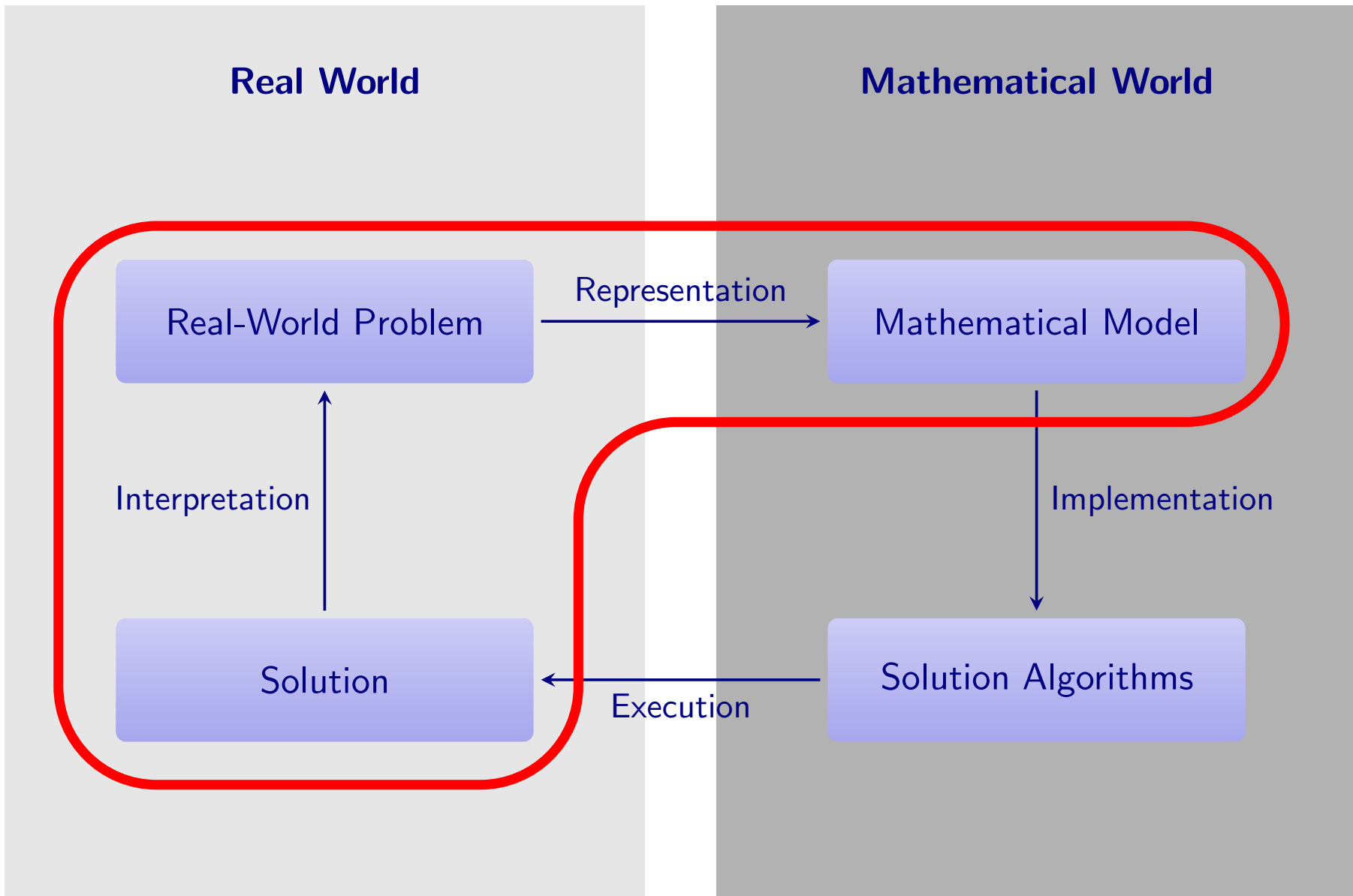
Mathematical Tools for Engineering and Management

Lecture 12

18 Jan 2012



- ▷ Models, Data and Algorithms
- ▷ Linear Optimization
- ▷ Mathematical Background: Polyhedra, Simplex-Algorithm
- ▷ Sensitivity Analysis; (Mixed) Integer Programming
- ▷ MIP Modelling
- ▷ MIP Modelling: More Examples; Branch & Bound
- ▷ Cutting Planes; Combinatorial Optimization: Examples, Graphs, Algorithms
- ▷ TSP-Heuristics
- ▷ Network Flows
- ▷ Shortest Path Problem
- ▷ Complexity Theory
- ▷ **Nonlinear Optimization**
- ▷ Scheduling (Jan 25)
- ▷ Lot Sizing (Feb 01)
- ▷ Summary (Feb 08)
- ▷ Oral exam (Feb 15)



▷ Production Planning in Automobile Industry



Product	Beetle	Cabrio
Revenue	\$10000	\$20000
Manufacturing	5h	3h
Assembly	4h	7h
Raw material	400kg	400kg

Plant capacity and available raw materials:

- Manufacturing capacity: 50h
- Assembly capacity: 70h
- Raw material: 4500kg

▷ Production Planning in Automobile Industry



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▷ More realistic: Price of Cabrio depending on the demand

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 - ➔ Objective (total revenue): $10000x_b + 150000\sqrt{x_c} - 20000x_c$



▷ Model (non-linear program):

$$\begin{array}{llll} \text{maximize} & 10x_b + 150\sqrt{x_c} - 20x_c & & \\ \text{subject to} & 4x_b + 4x_c \leq 45 & & \text{(total raw material)} \\ & 5x_b + 3x_c \leq 50 & & \text{(time in manufacturing)} \\ & 4x_b + 7x_c \leq 70 & & \text{(time in assembly)} \\ & x_b, x_c \geq 0 & & \text{(non-negativity)} \end{array}$$

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▷ Optimal solution: $(x_b, x_c) = (5, 6.25)$

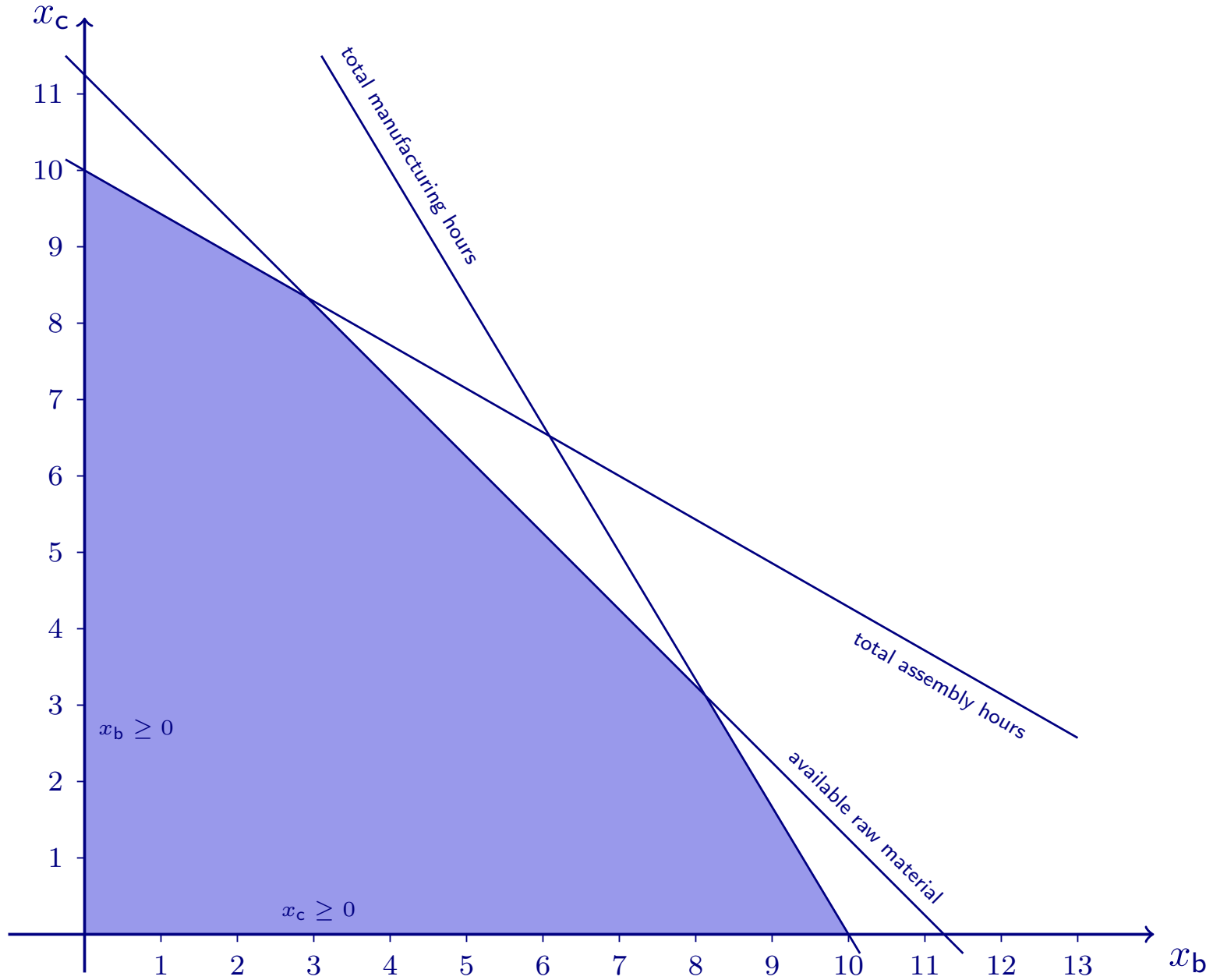
▷ Model (non-linear program):

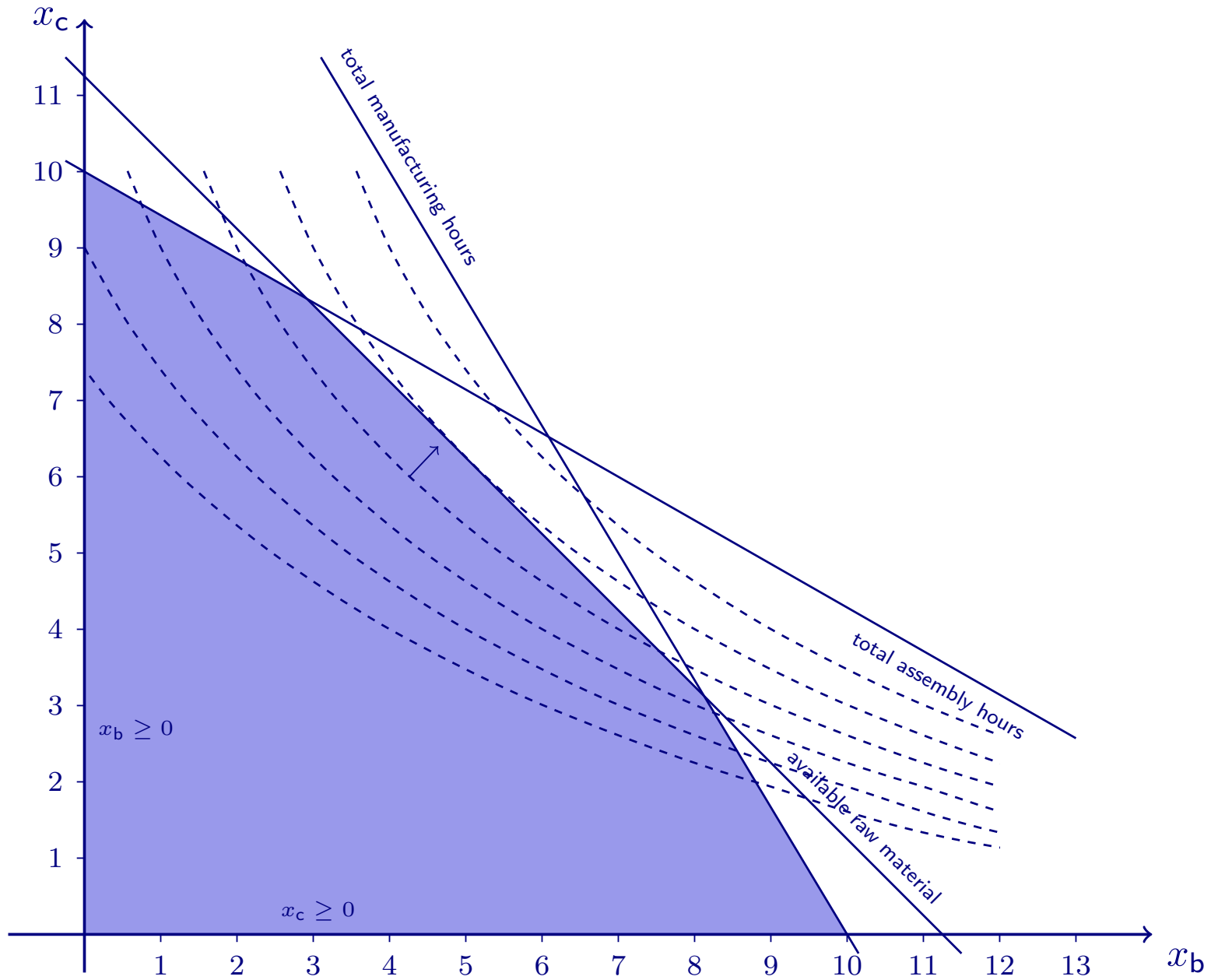
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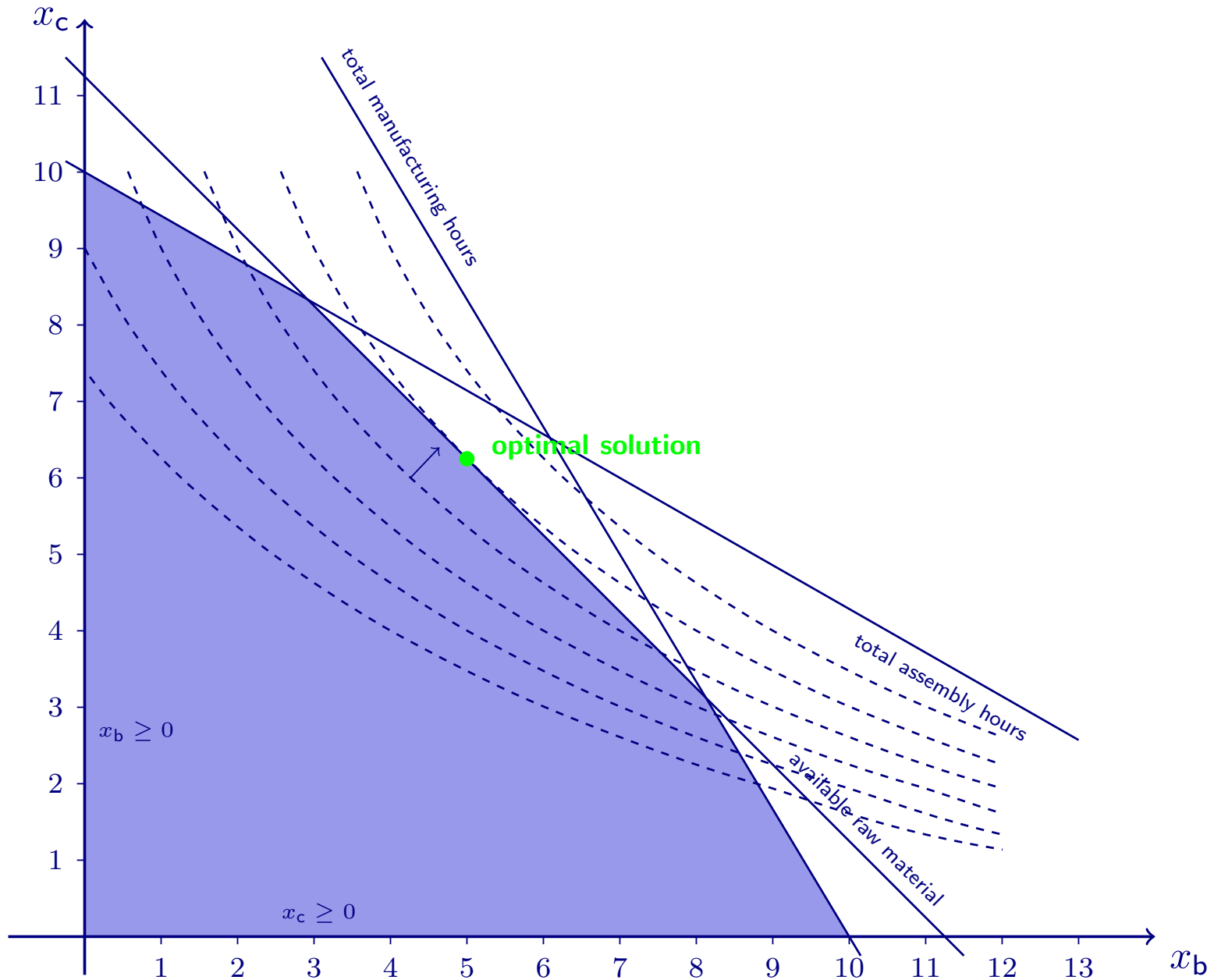
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➔ Price for one cabrio at this demand: 60000

➔ Profit for one cabrio: 40000









▷ Linear optimization

➔ Linear objective **and** linear constraints

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➔ Sum of linear terms: $\dots + a_i \cdot x_i + \dots$
 ↑ ↑
 parameter variable

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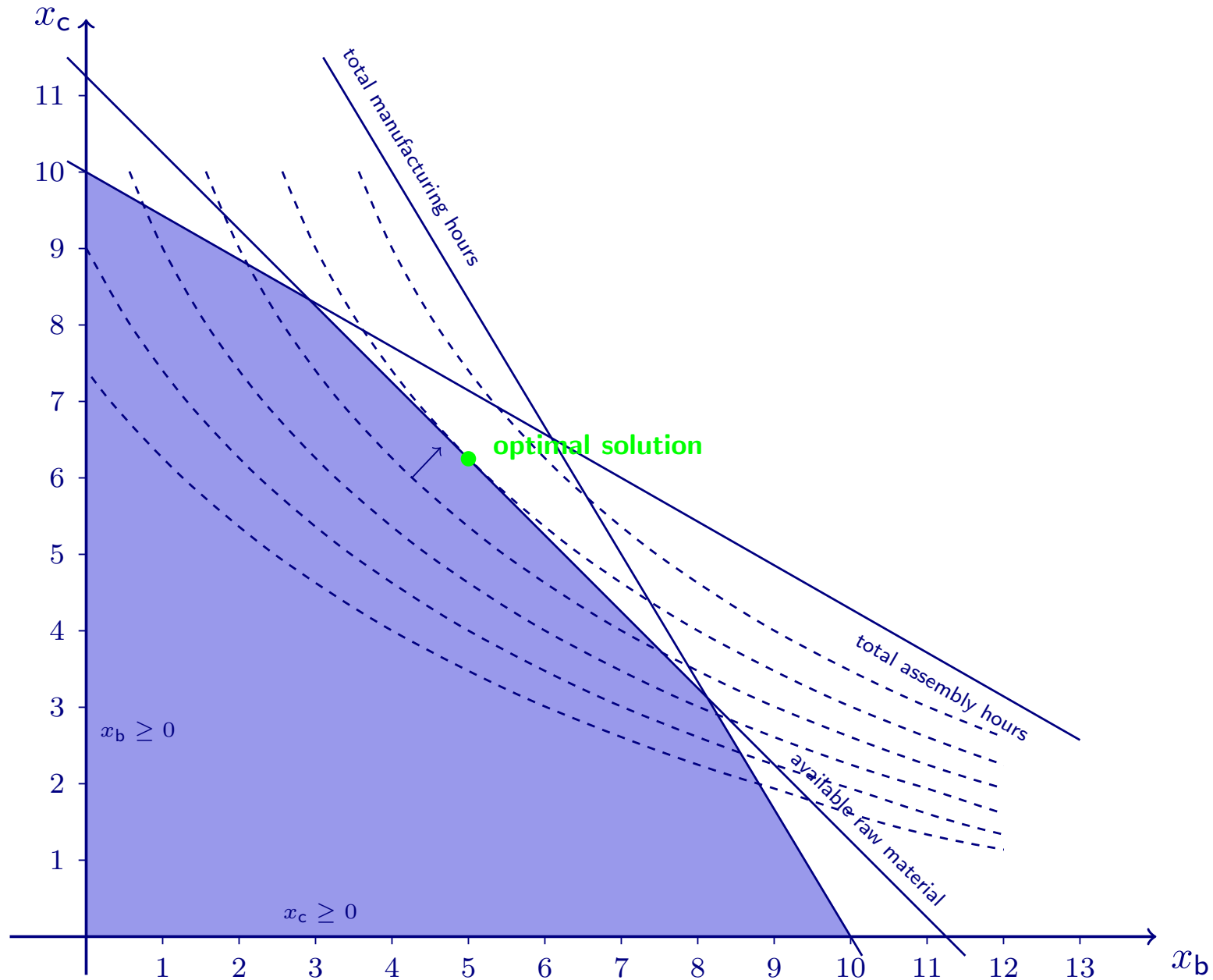
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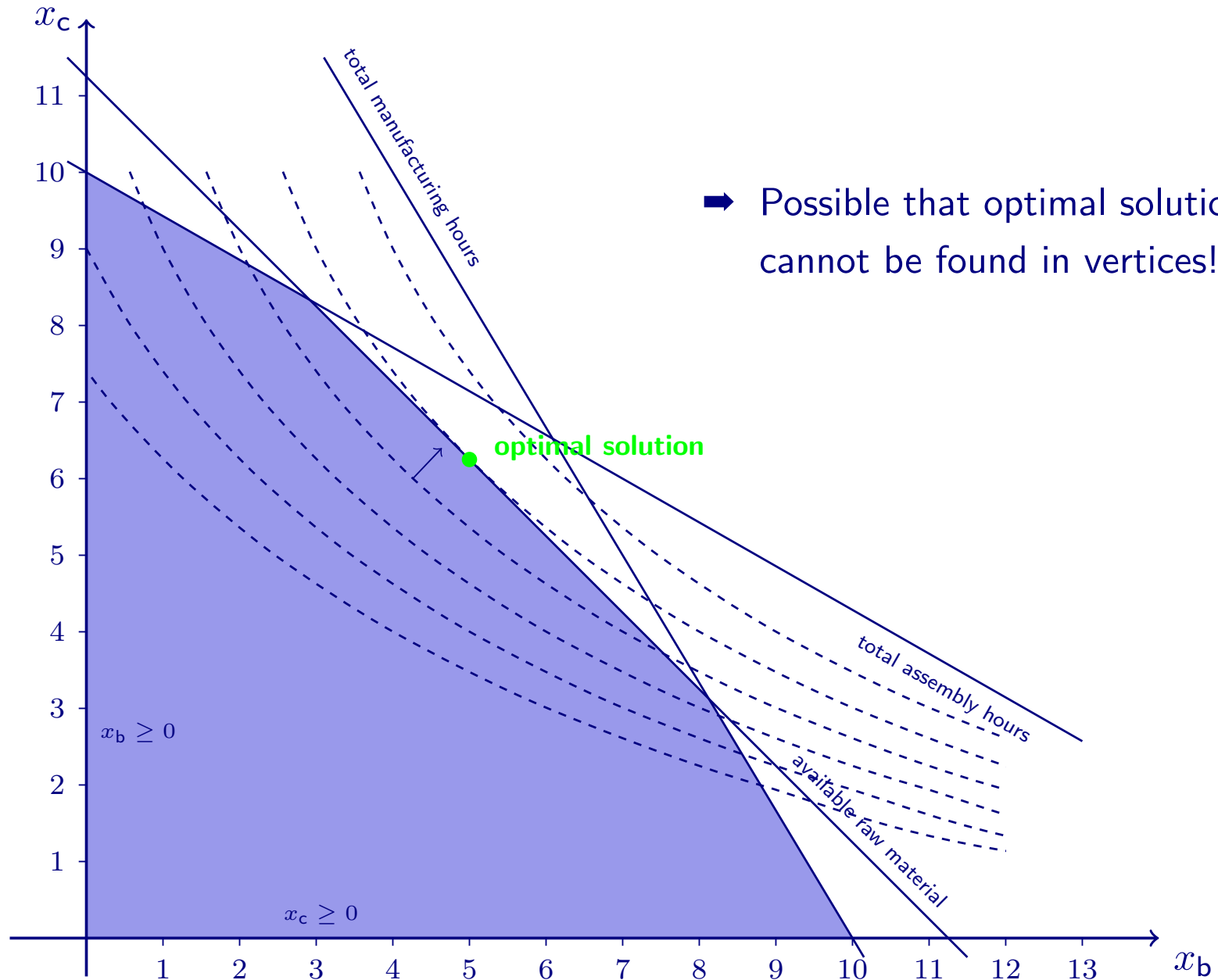
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- Absolute values or maxima/minima: $|x_i|, \max x_j$
- Terms including elementary functions: $\sin x_i, 2^{x_i \cdot x_j}, \frac{1}{\sqrt{x_i}}, \log(x_i + x_j^{x_k})$



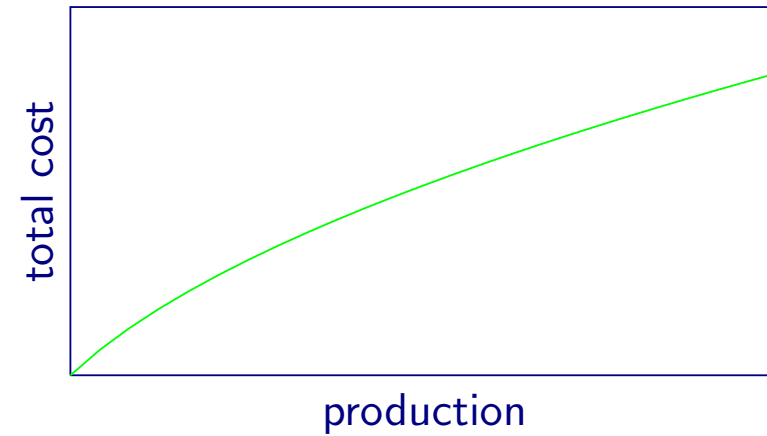




▷ Economy of scale

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➔ Production cost per item decrease with number x_i of produced items

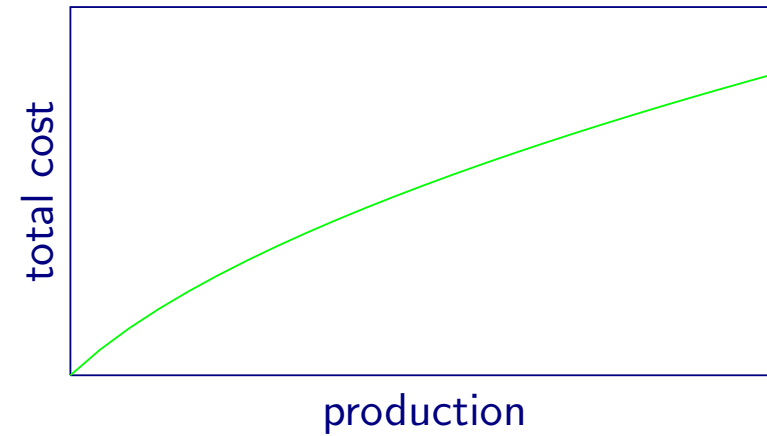


▷ Economy of scale

- ➔ Production cost per item decrease with number x_i of produced items
- ➔ Contribution to objective (examples):

$$\dots + \sqrt{x_i} + \dots$$

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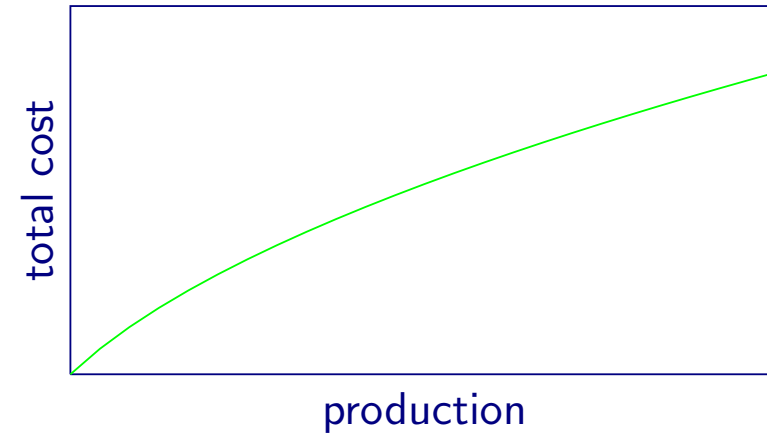
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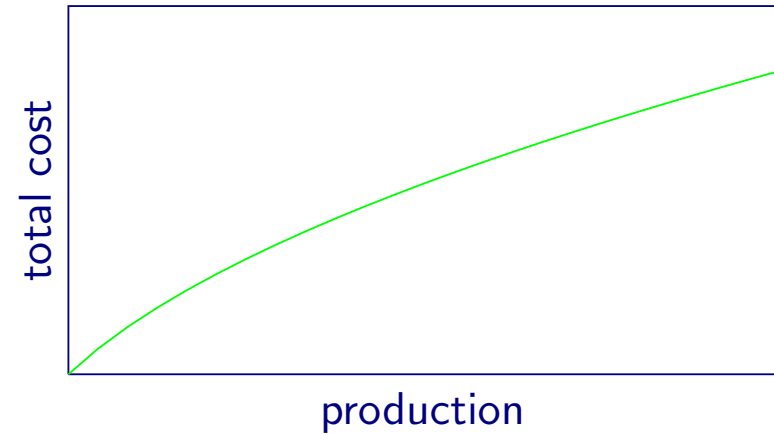
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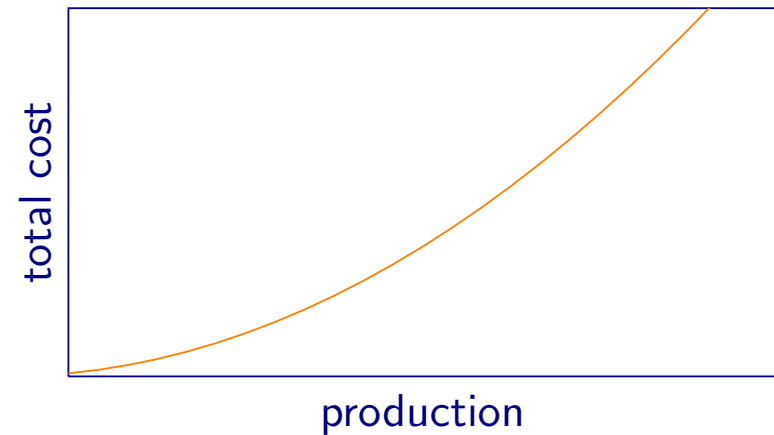
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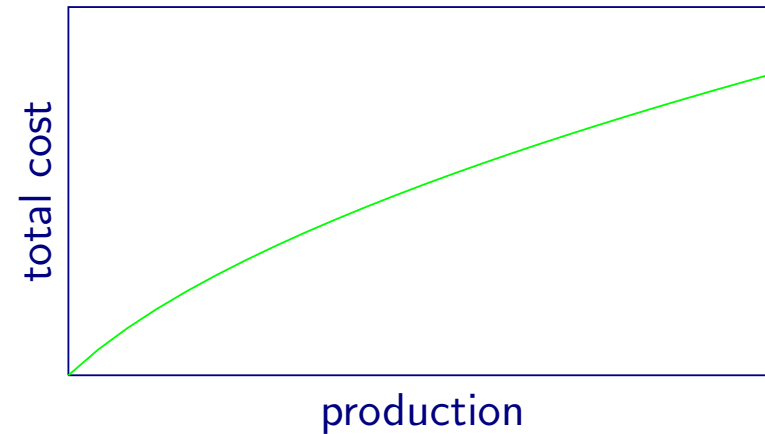


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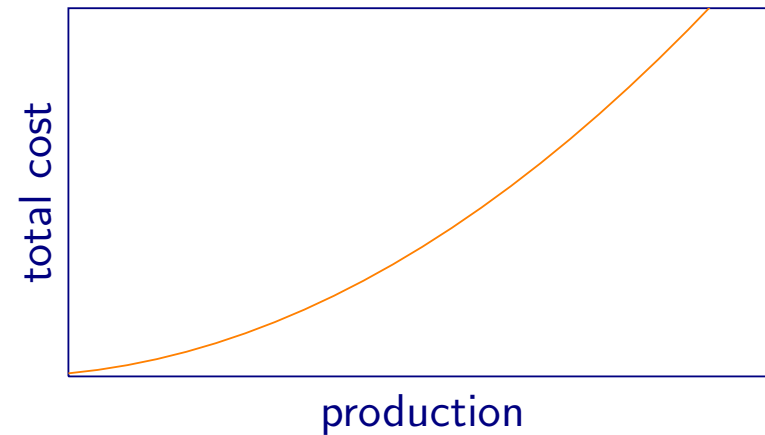


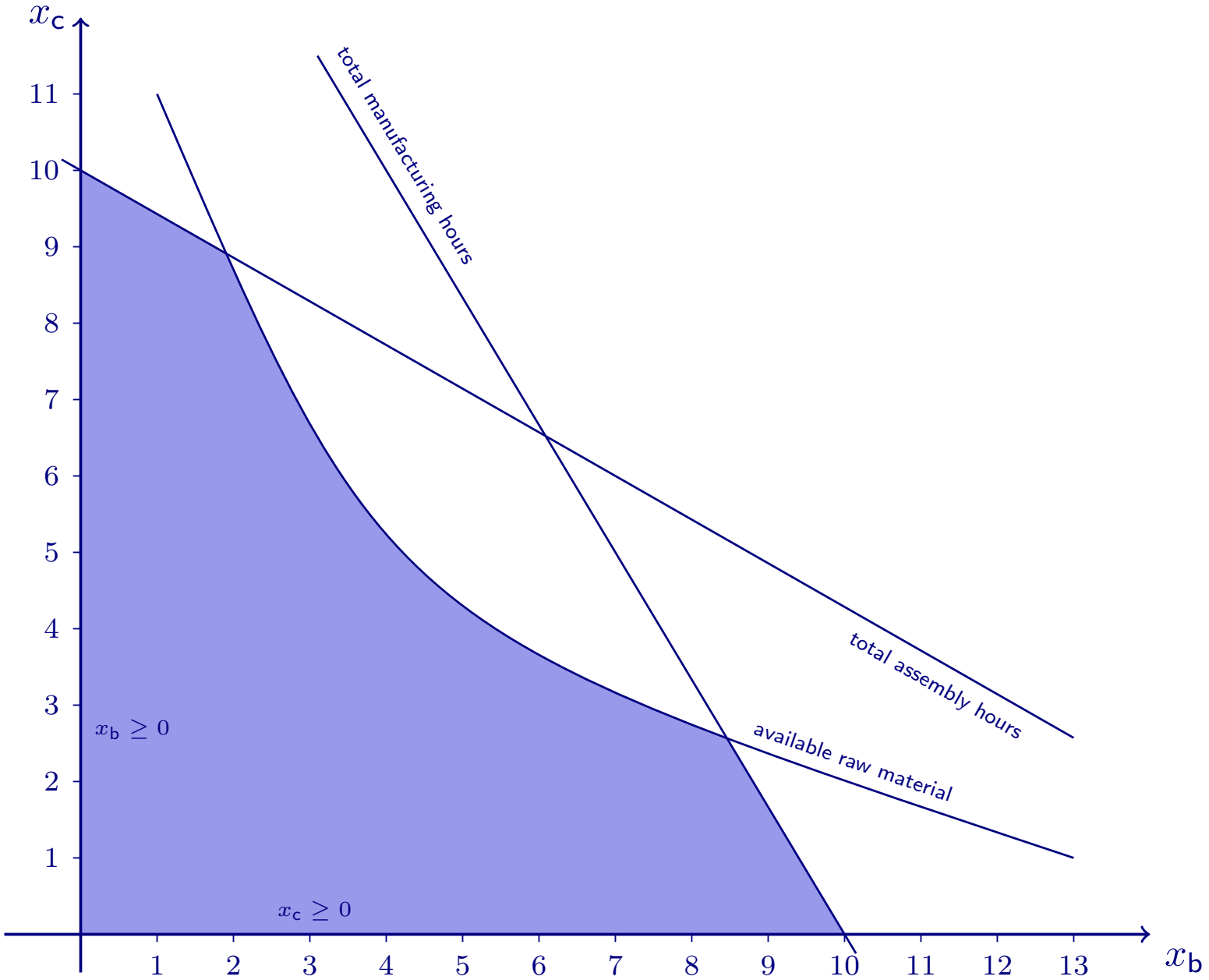
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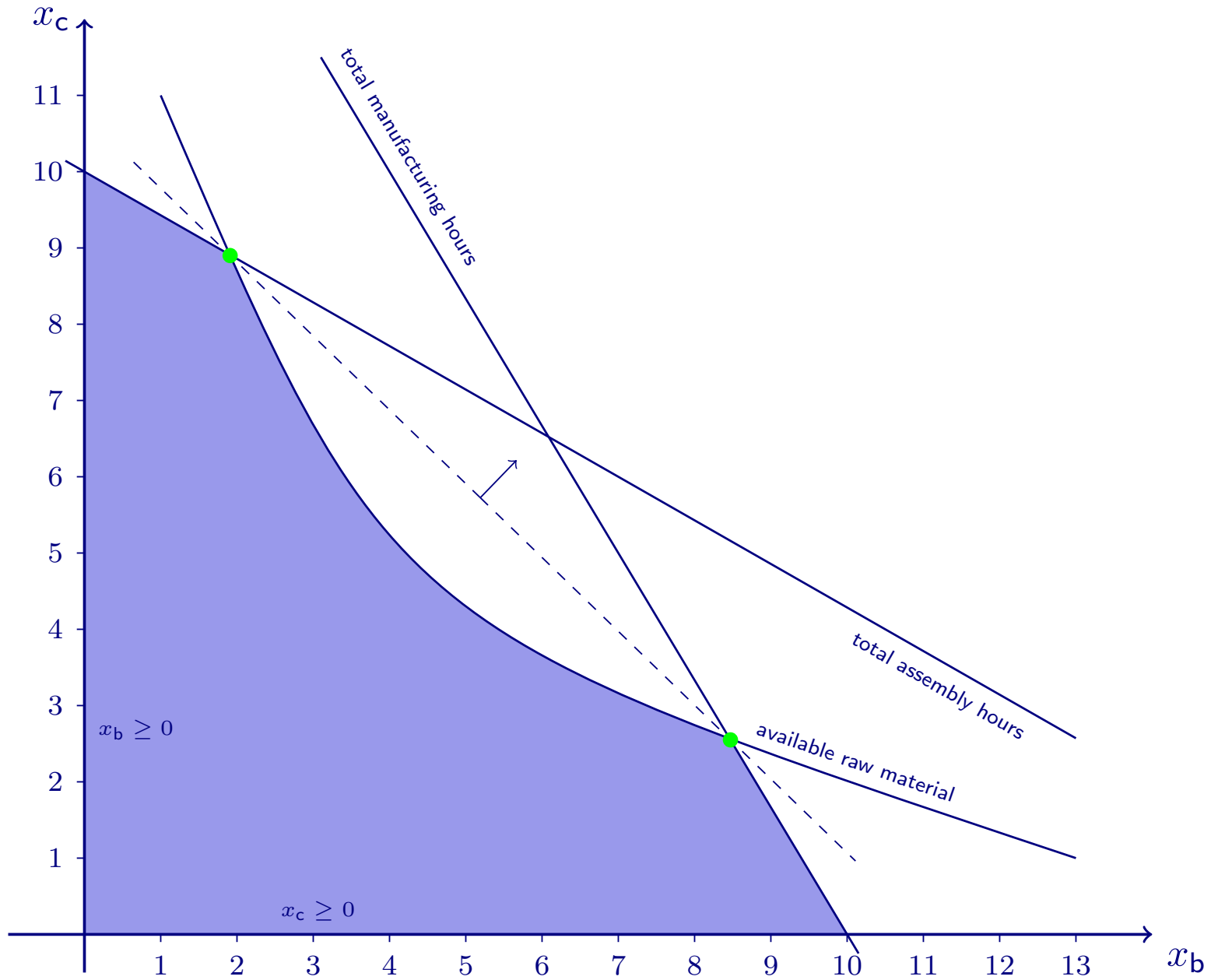
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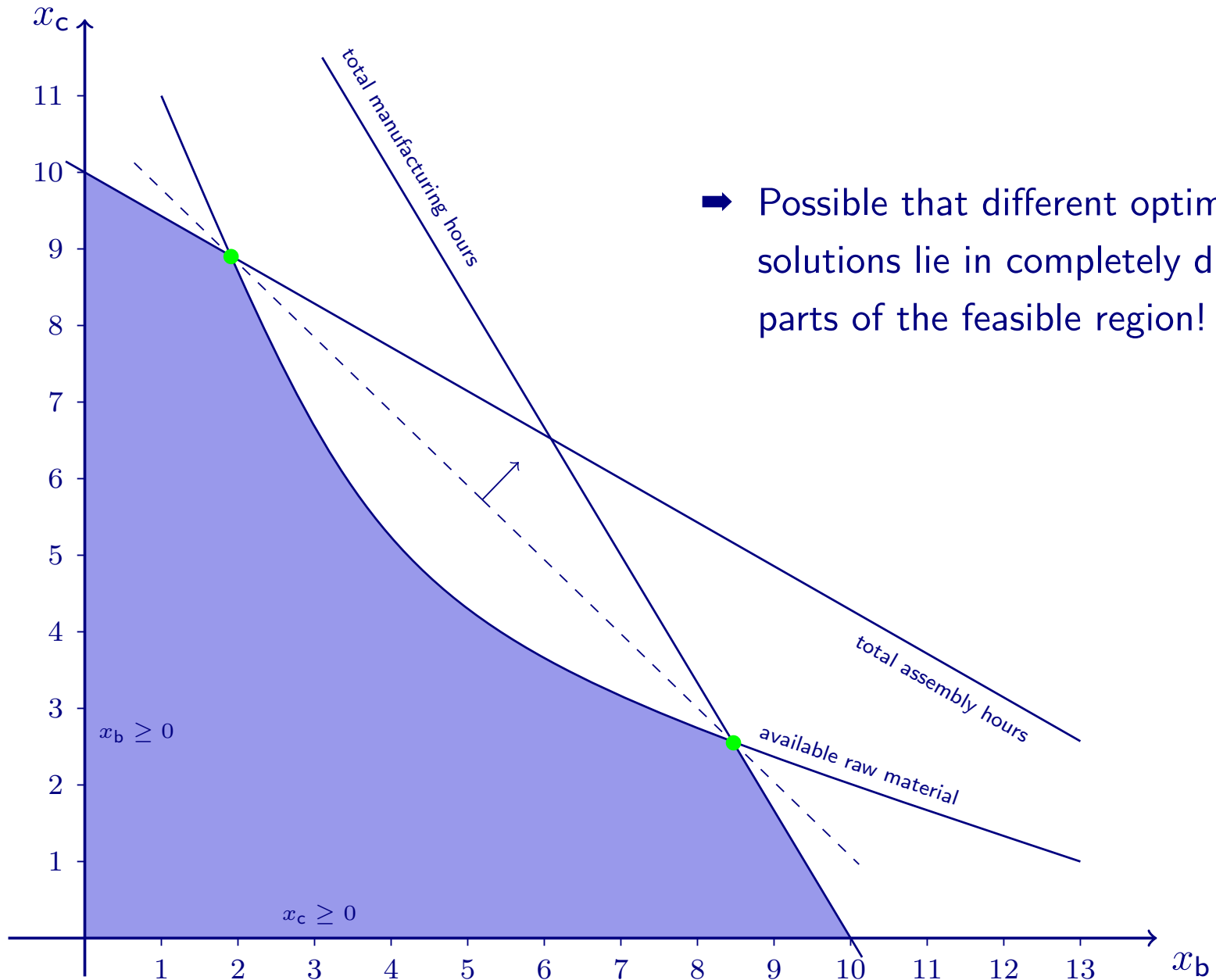
$$\dots + x_i^2 + \dots$$

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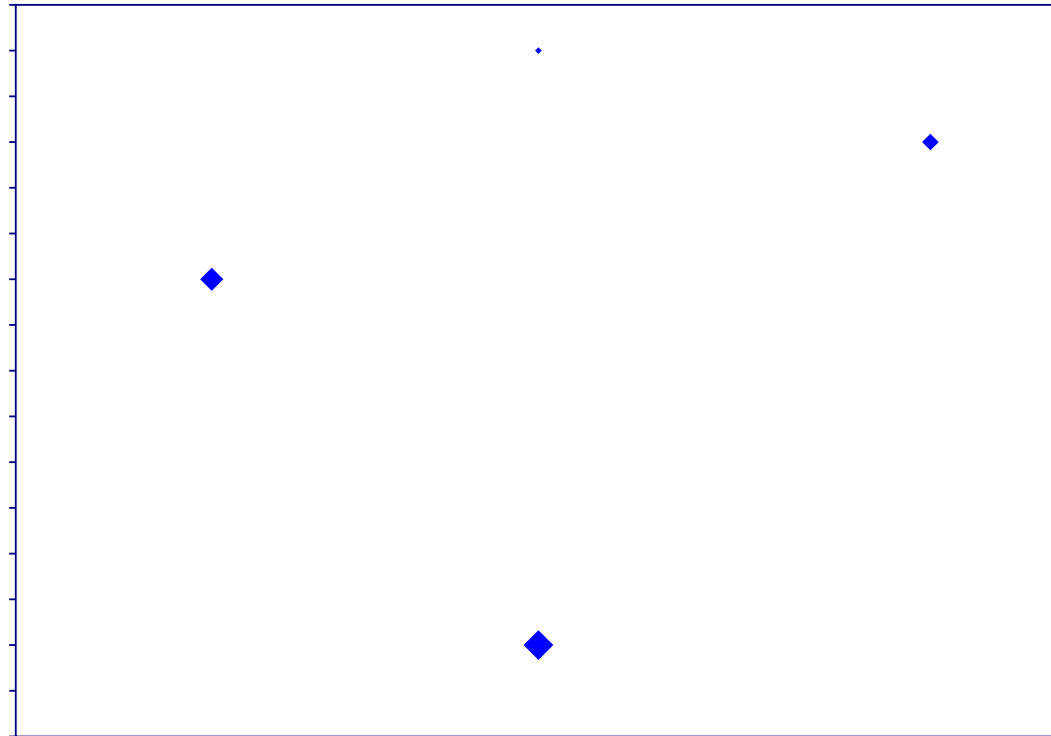


➔ Possible that different optimal solutions lie in completely different parts of the feasible region!

- ▷ Given: locations at specified coordinates in the plane
- ▷ Task: Find an optimal location for a central unit connecting every given location!

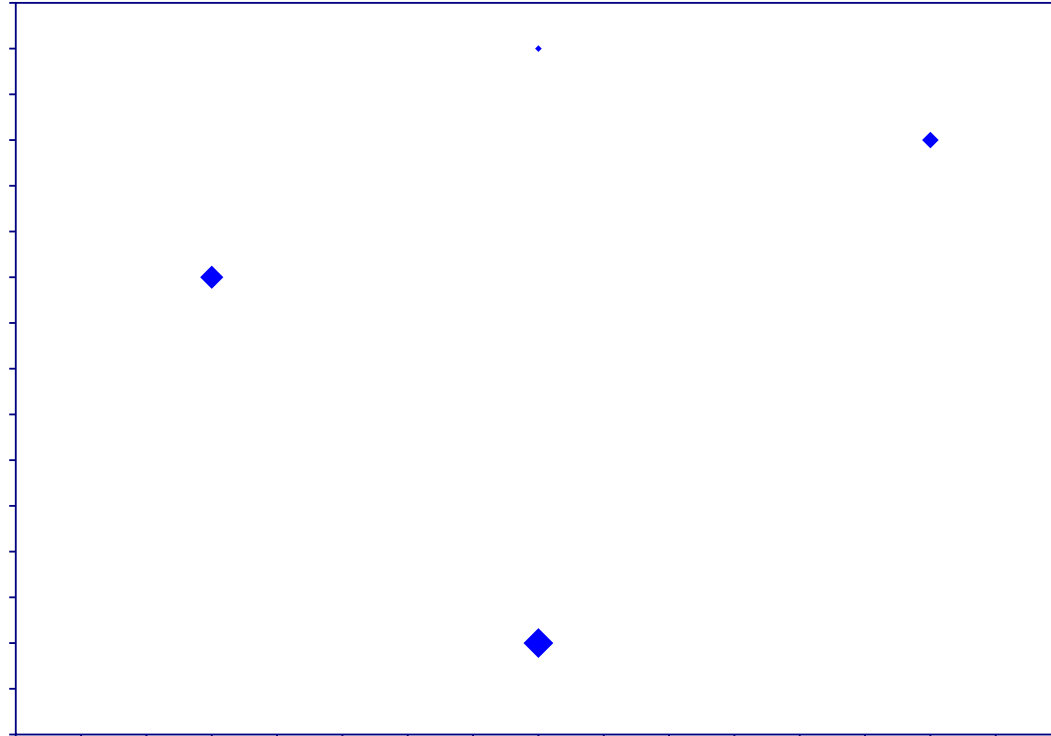
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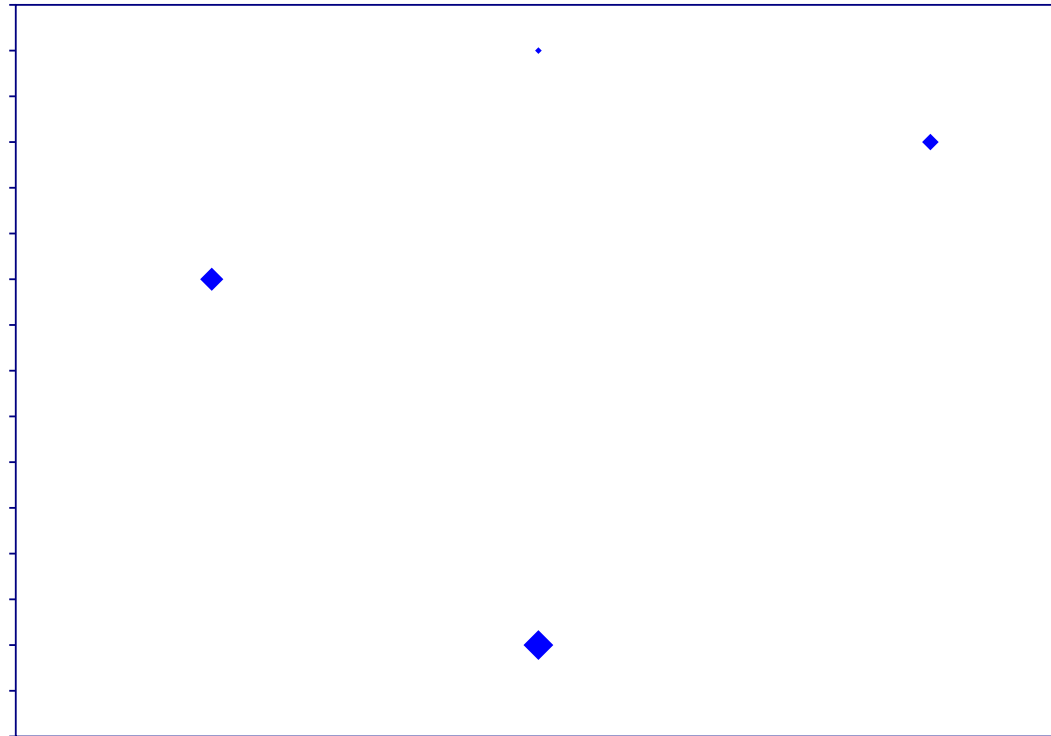
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Variables: x, y → coordinates of central unit

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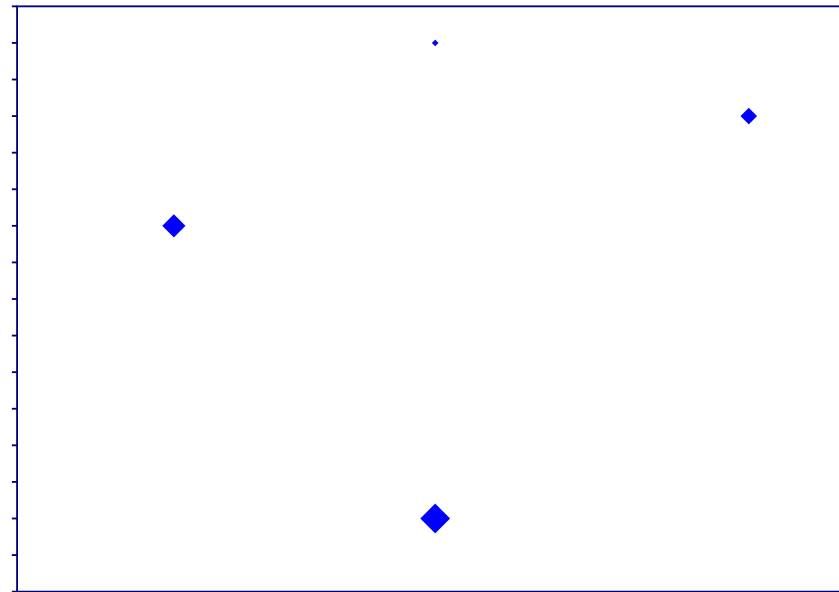
Variables: x, y → coordinates of central unit

Objective: minimize sum of connection costs to all given locations



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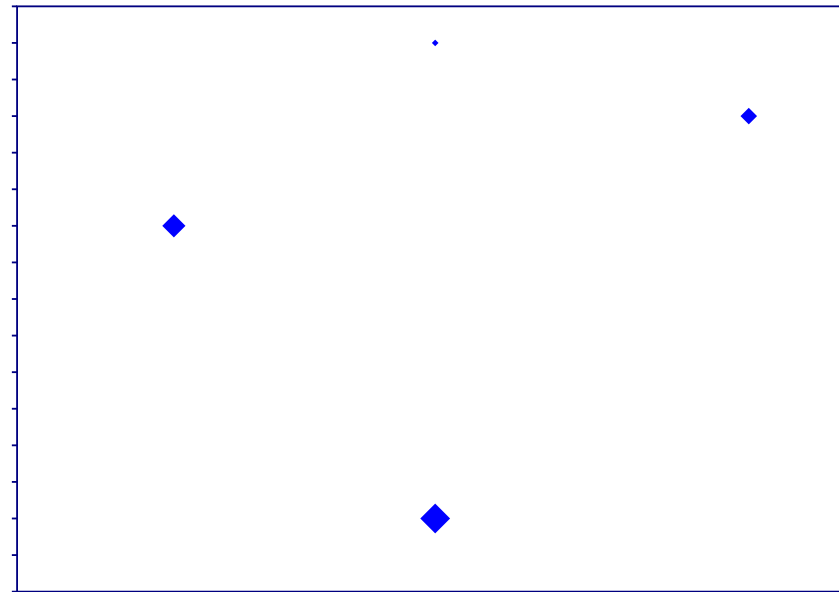




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→ Distance to **A**: $\sqrt{(x - 8)^2 + (y - 2)^2}$

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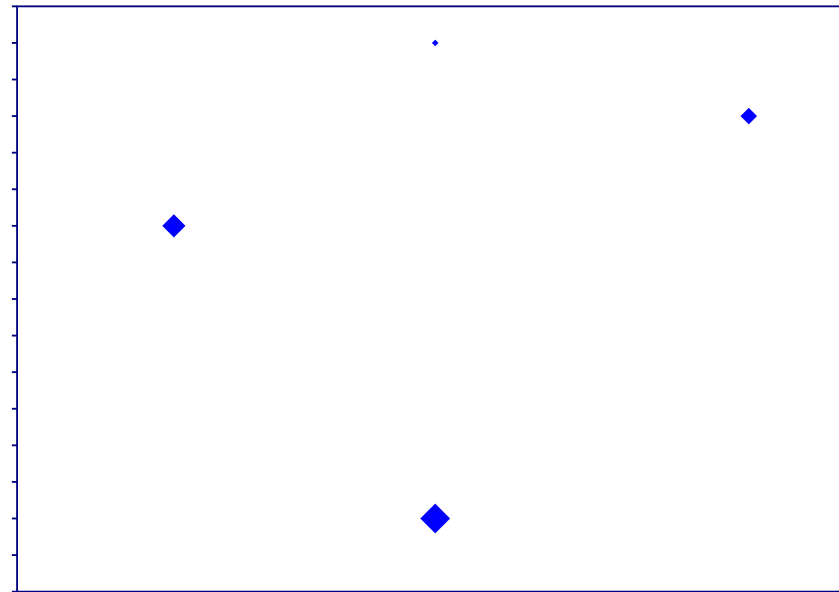


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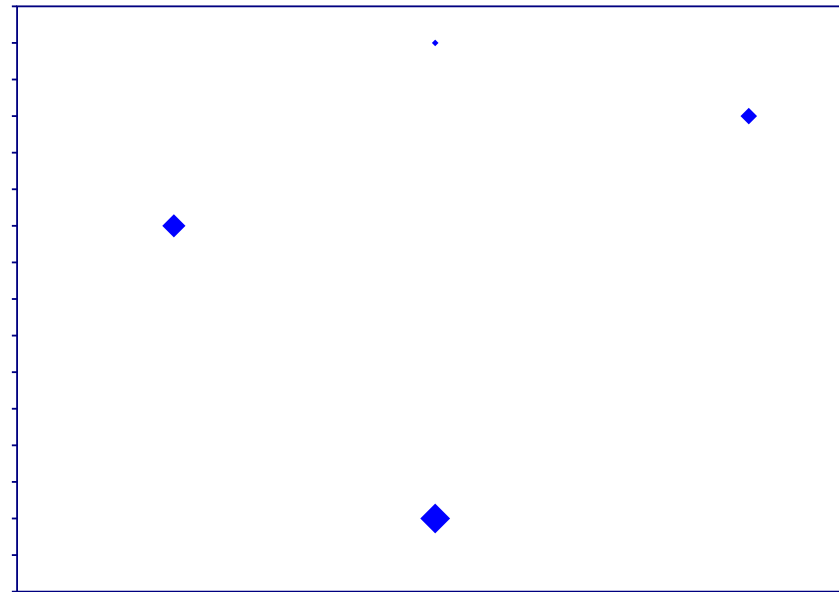
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→ Analogous for **B**, **C**, and **D**

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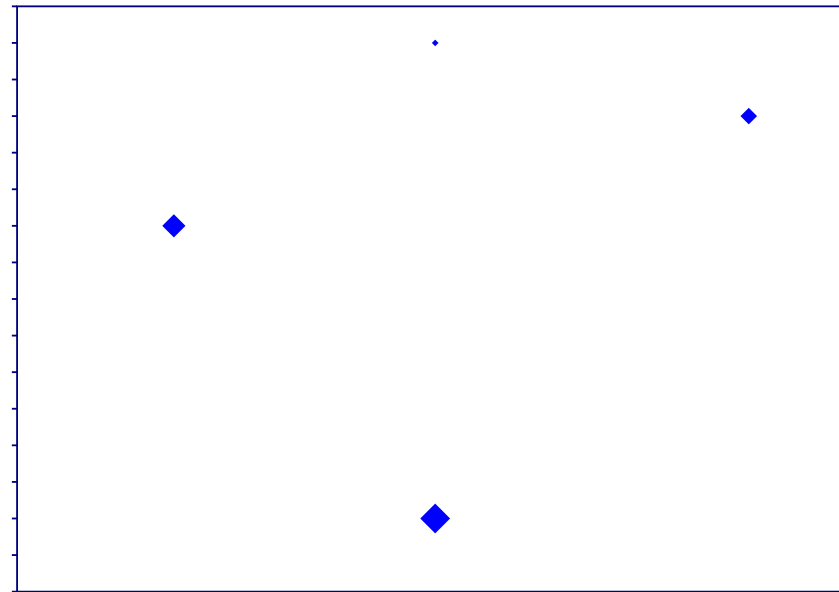
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→ **Objective:**
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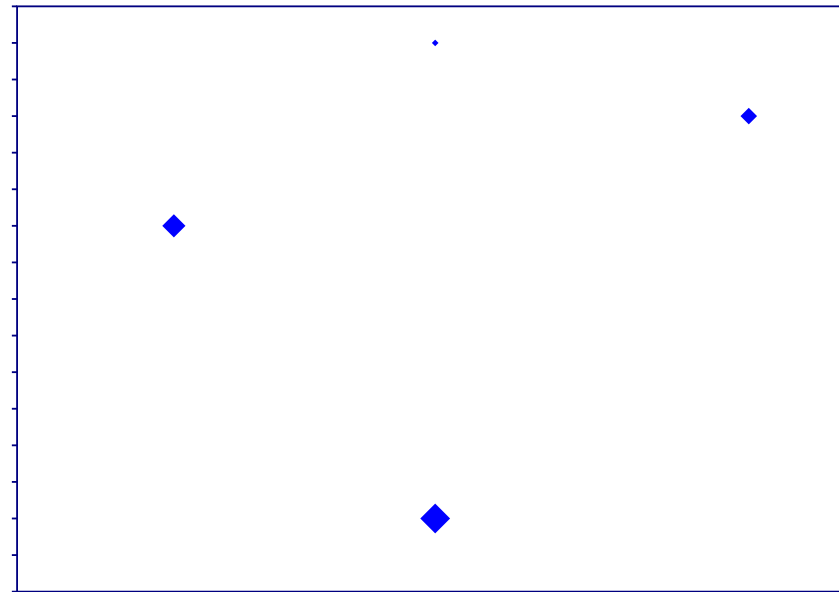


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V

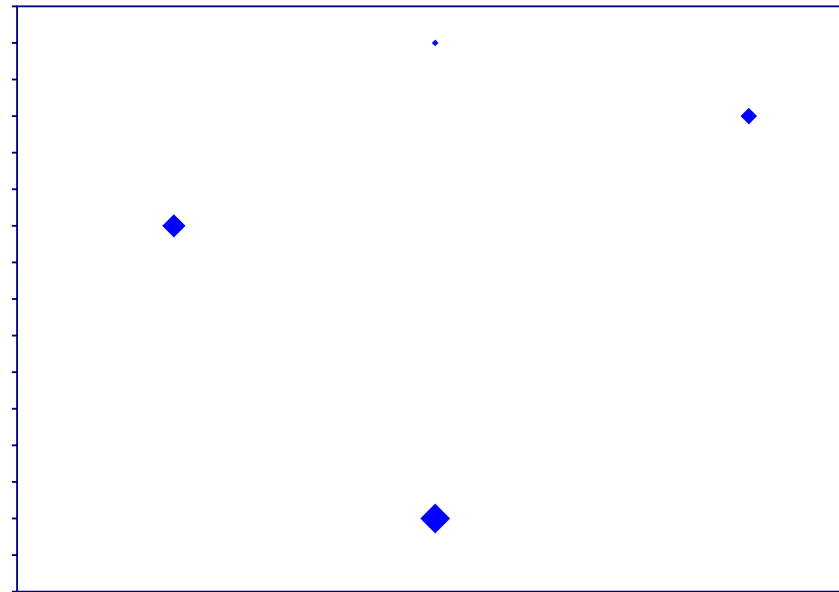
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No explicit constraints!

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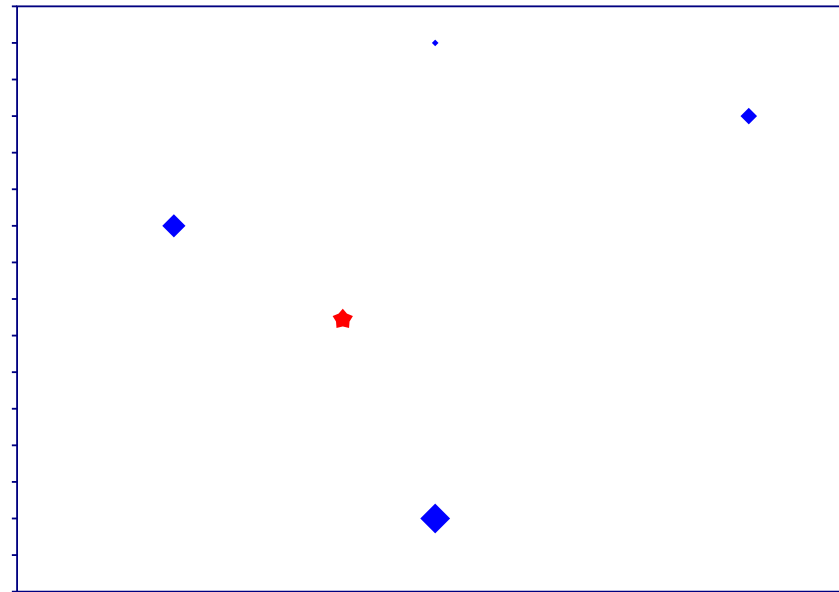
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No explicit constraints!

→ Optimal solution: $(x, y) = (6.25, 7.47)$ (unique!)

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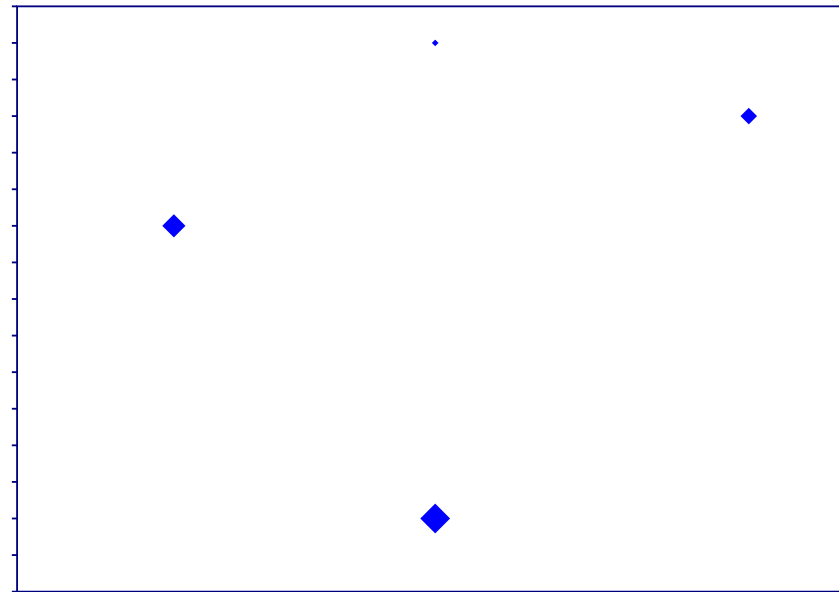
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Constraints

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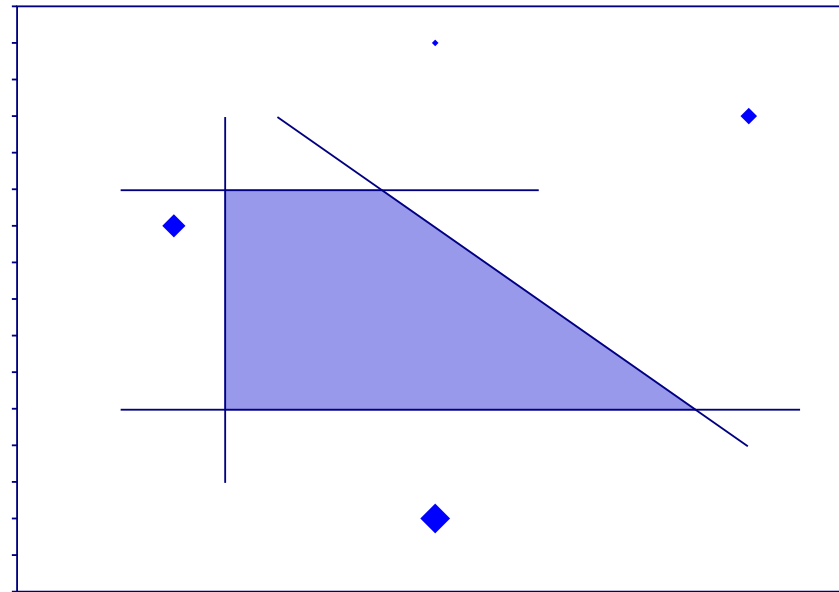
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Constraints: $x \geq 4, \quad y \geq 5, \quad y \leq 11, \quad x + y \leq 18$

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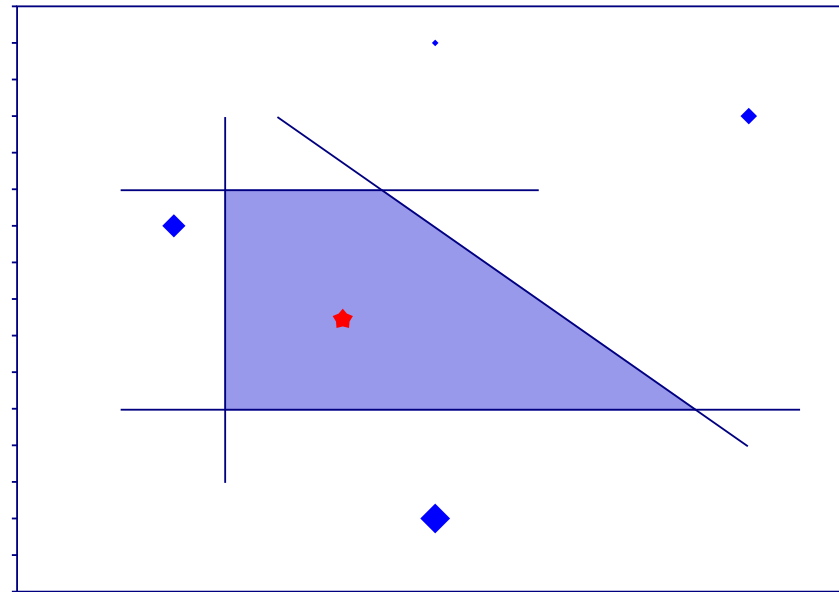
Objective:
$$\min \quad 9 \sqrt{(x - 8)^2 + (y - 2)^2} + 7 \sqrt{(x - 3)^2 + (y - 10)^2} \\ + 2 \sqrt{(x - 8)^2 + (y - 15)^2} + 5 \sqrt{(x - 14)^2 + (y - 13)^2}$$



Constraints: $x \geq 4, \quad y \geq 5, \quad y \leq 11, \quad x + y \leq 18$

→ Optimal solution: $(x, y) = (6.25, 7.47)$ (same as before!)

Location	coordinates	# connections
A	(8,2)	9
B	(3,10)	7
C	(8,15)	2
D	(14,13)	5





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➔ How to find an optimal solution...?



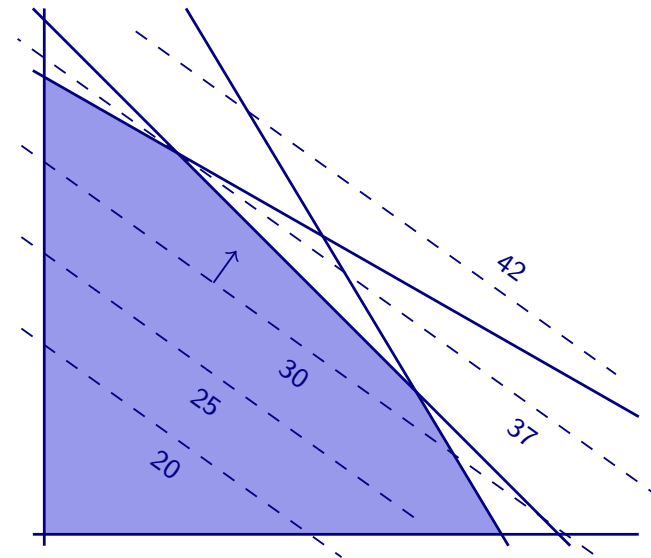
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- Linear objective
 - ➔ Level sets are straight lines
(in higher dimension: hyperplanes)
- Linear constraints
 - ➔ Feasible region is a polygon
(in higher dimension: polyhedron)

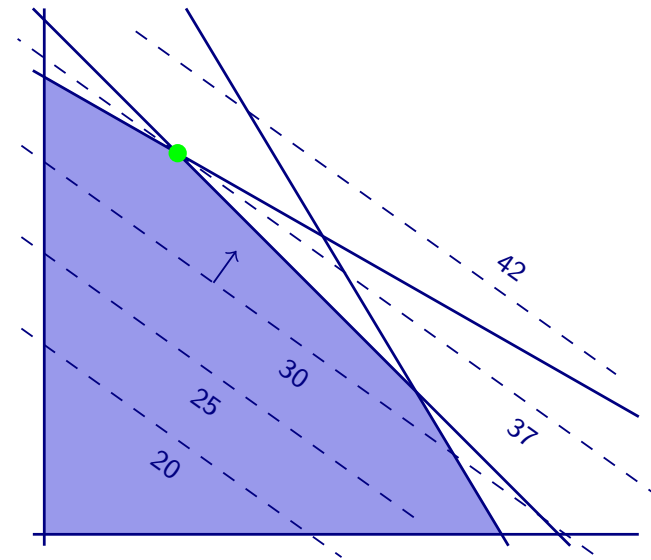
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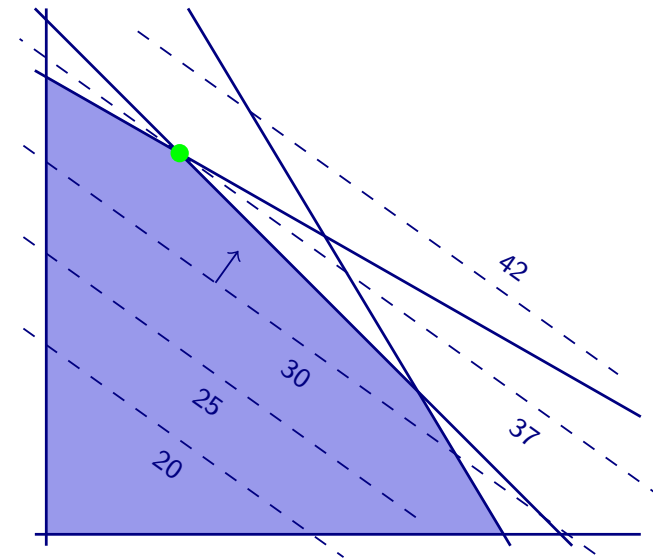
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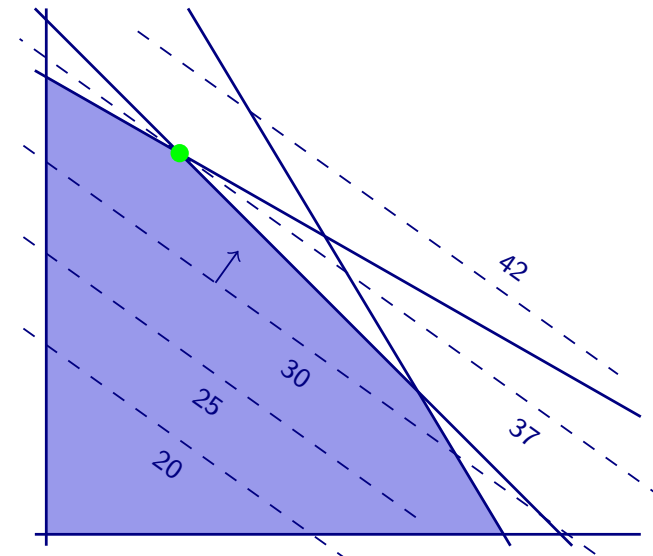
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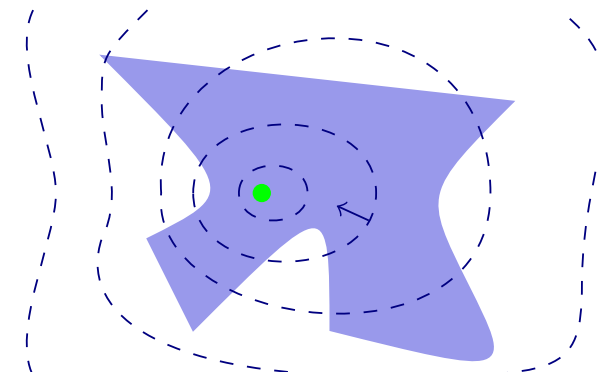
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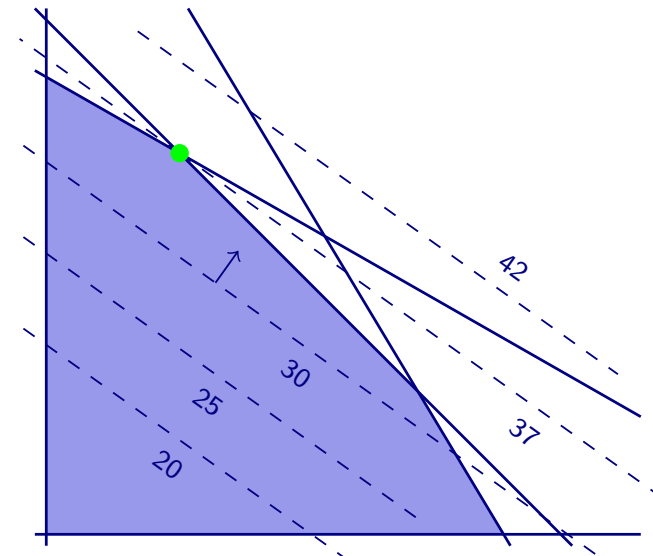
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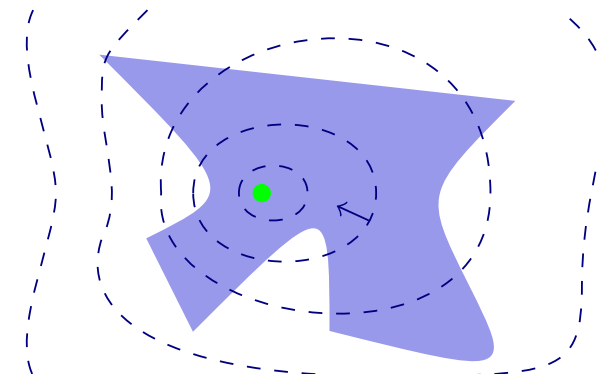
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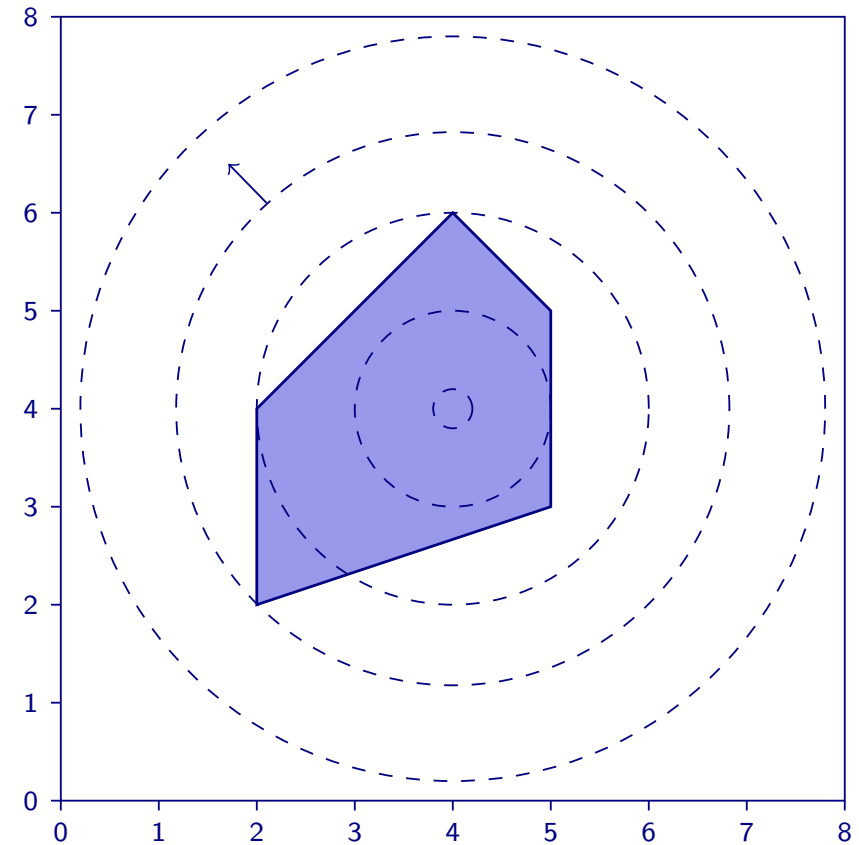
➔ Finding optimal solution can be difficult

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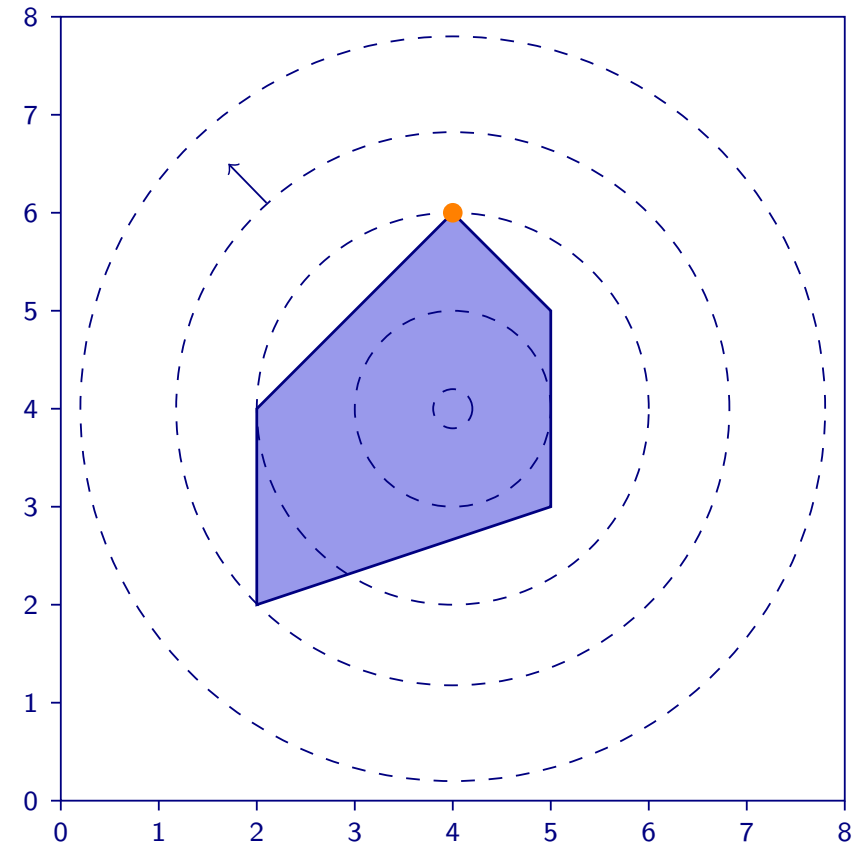
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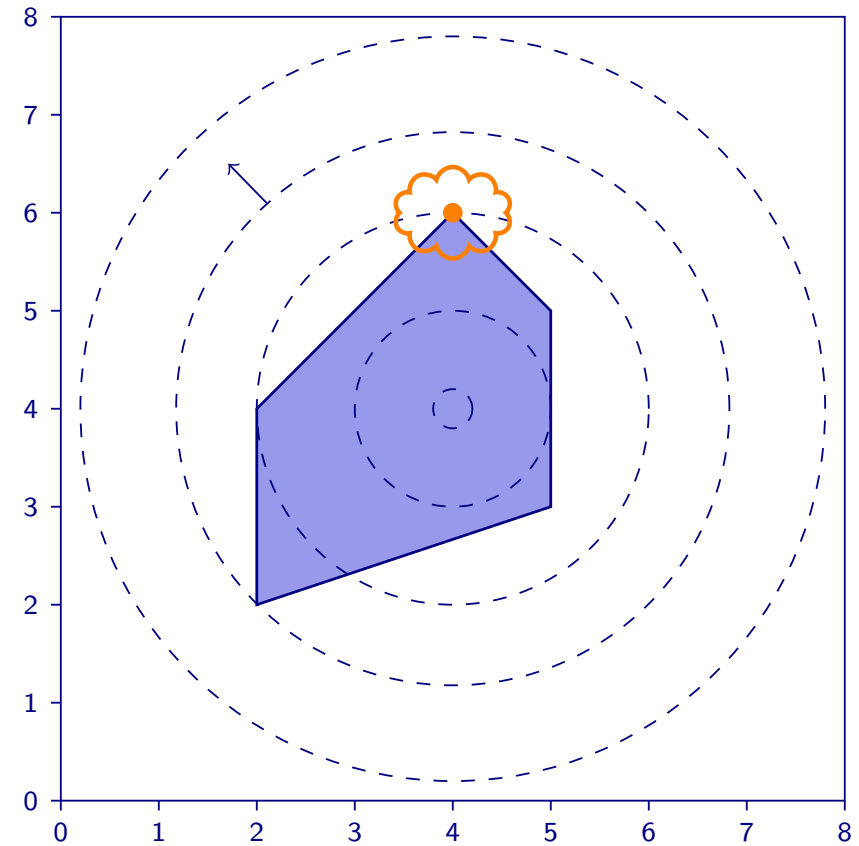
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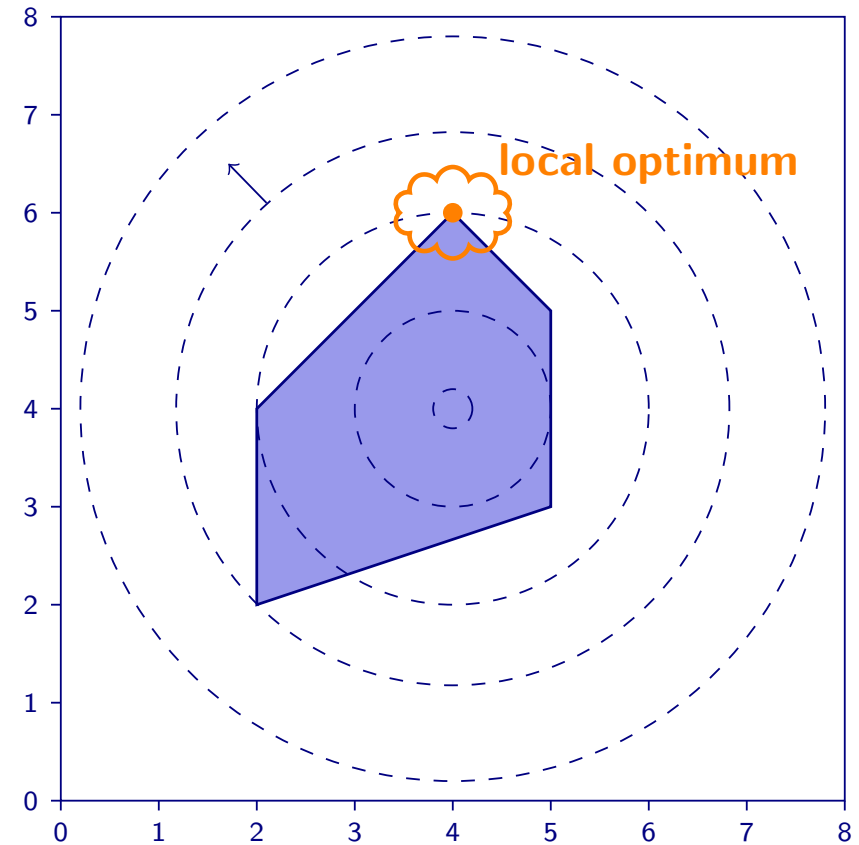
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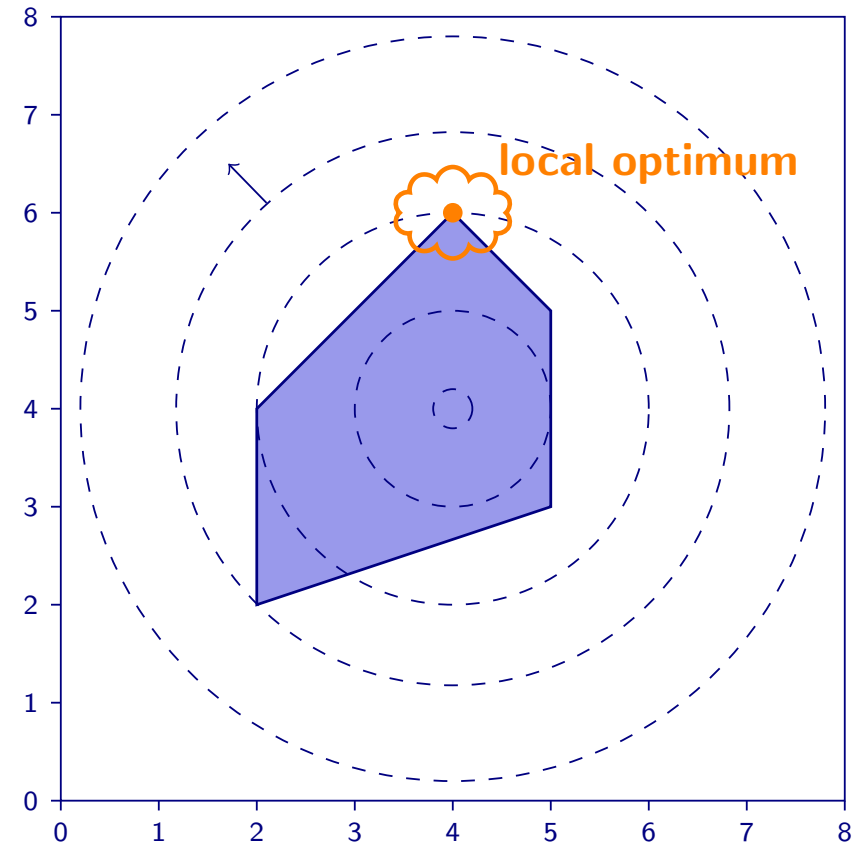
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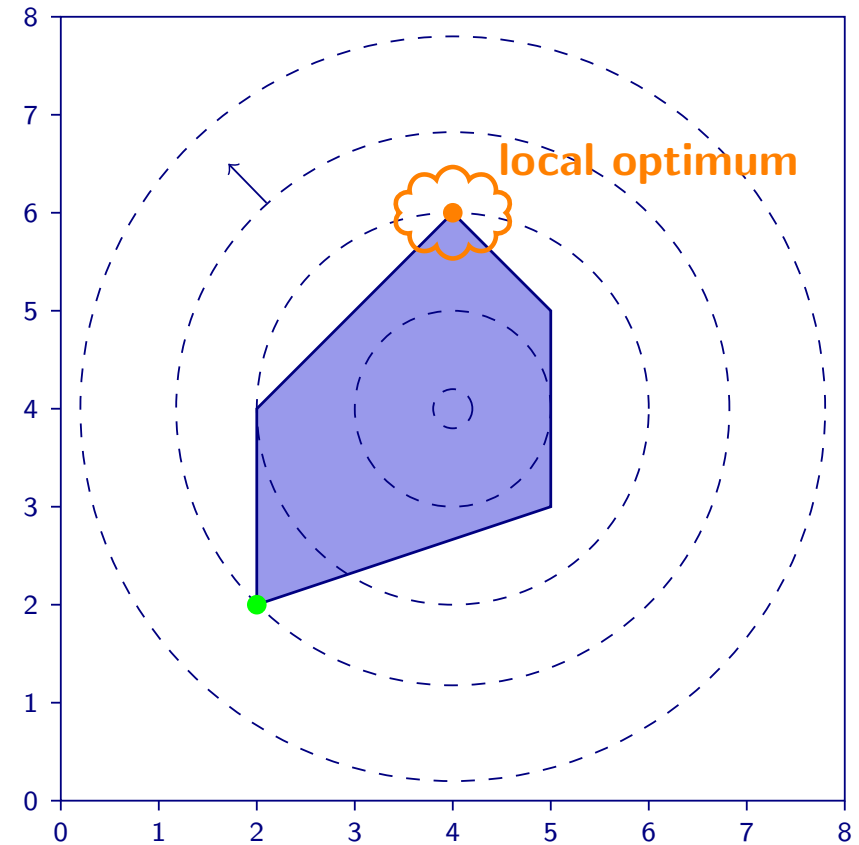
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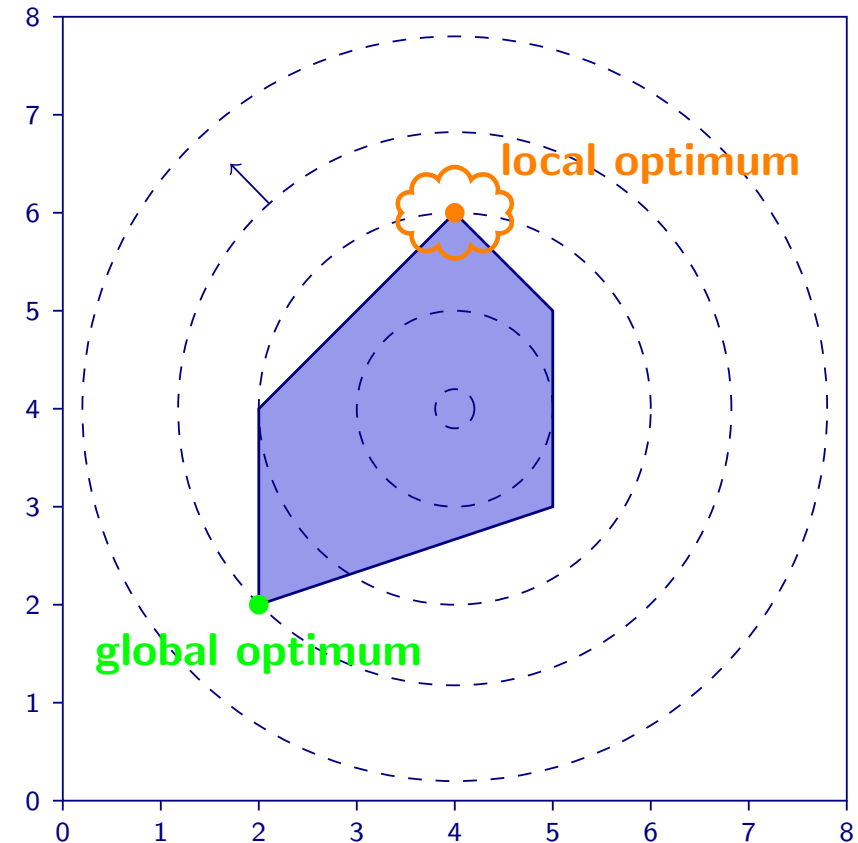
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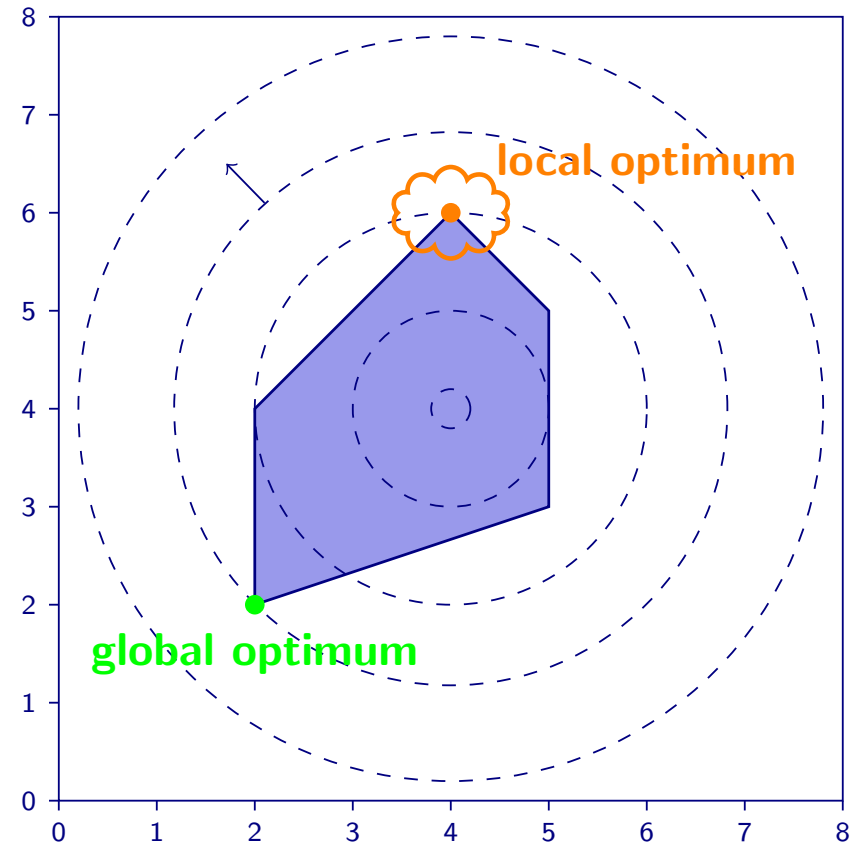
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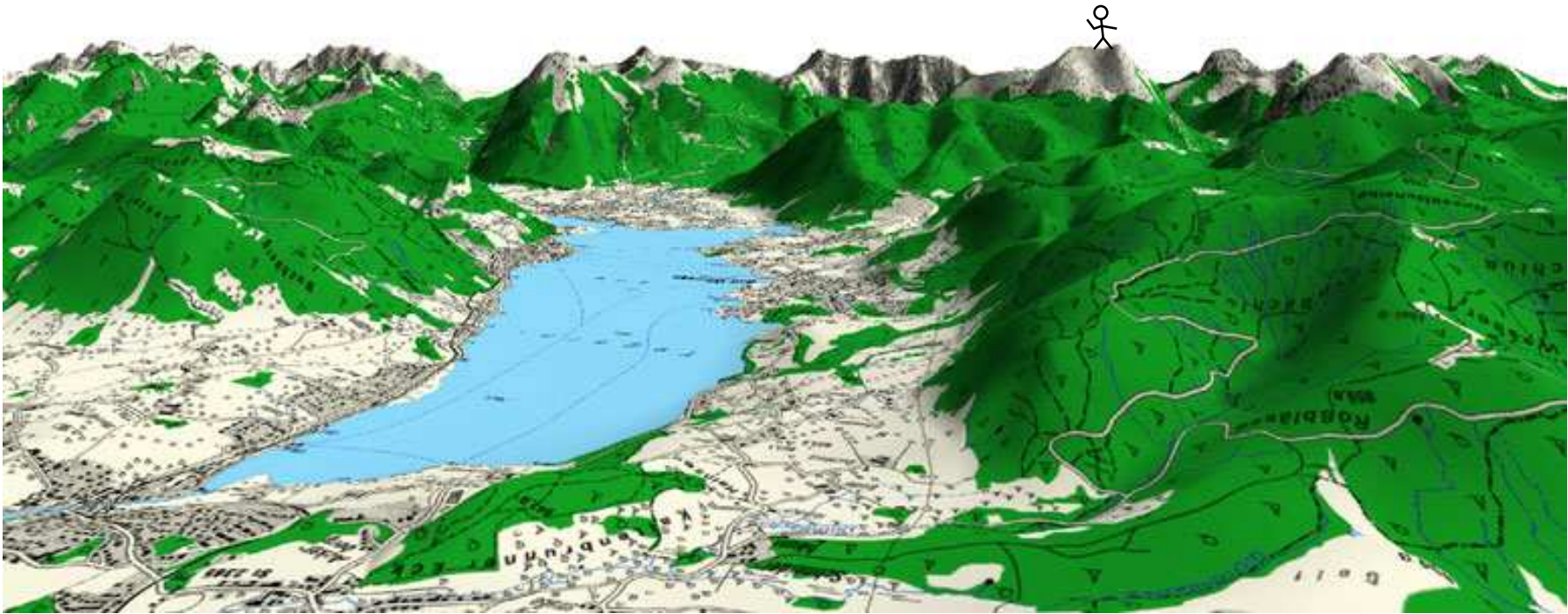
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- ▷ Non-linear optimization is like mountain-climbing in the fog



- ▶ Non-linear optimization is like mountain-climbing in the fog



- ▶ How do you know that you're on the highest mountain if you can't see the other peaks?

- ▶ Basic principle of interior point methods for finding a local maximum:

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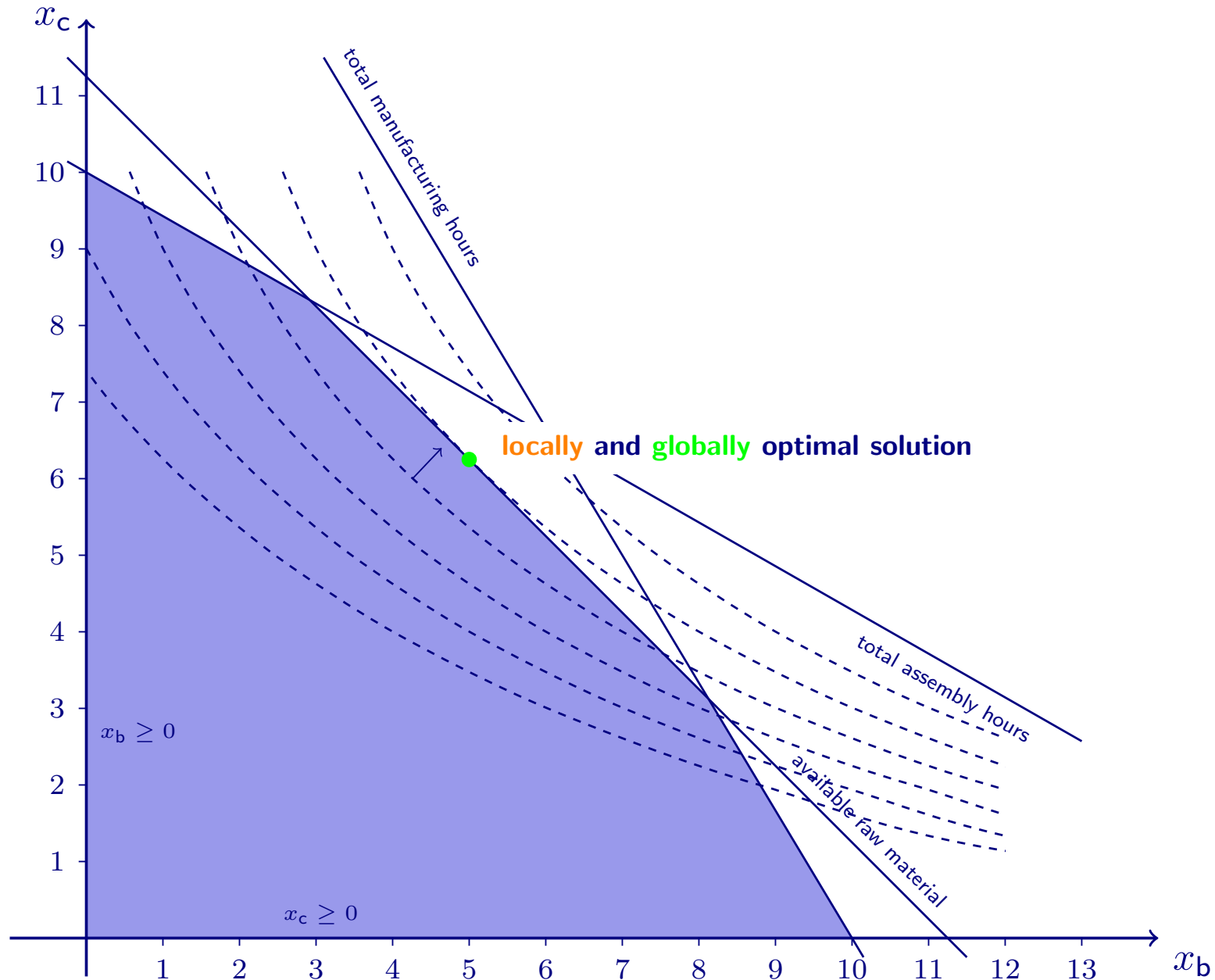
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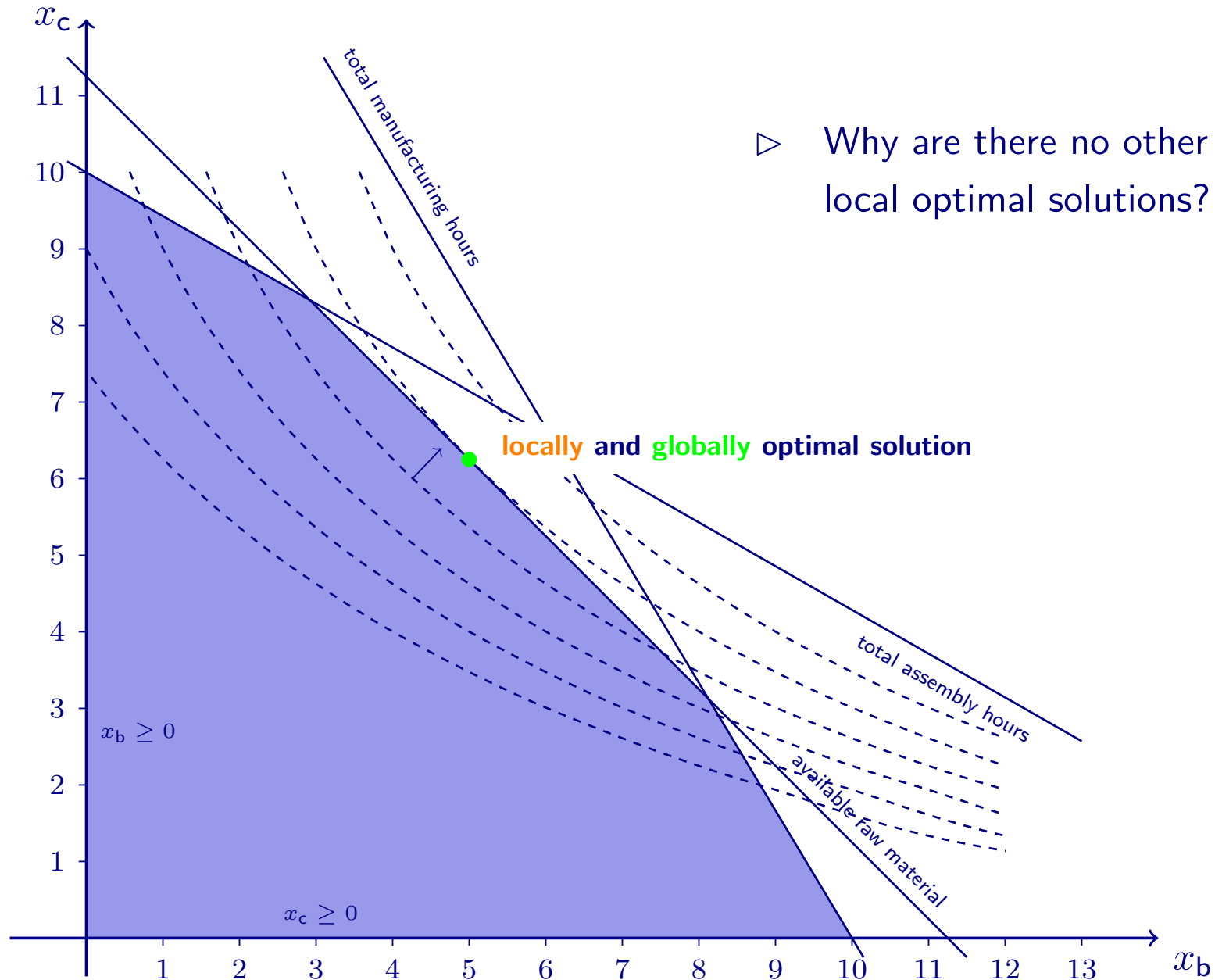
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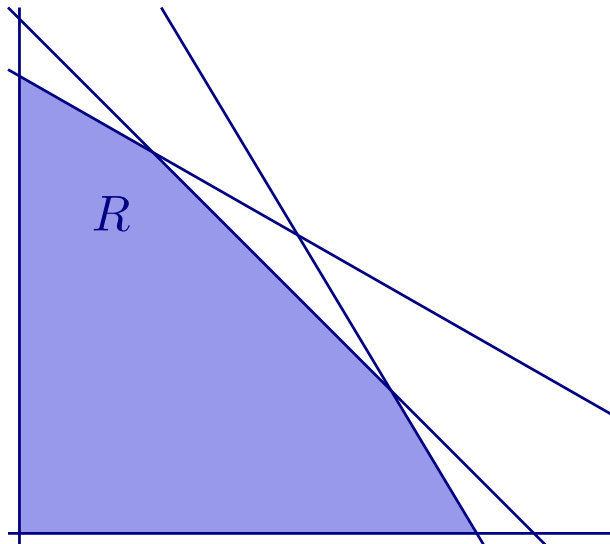
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- ▶ Big disadvantage: no optimality information (as gaps in branch & bound)!
- ➔ You have to rely on luck to get an optimal solution...



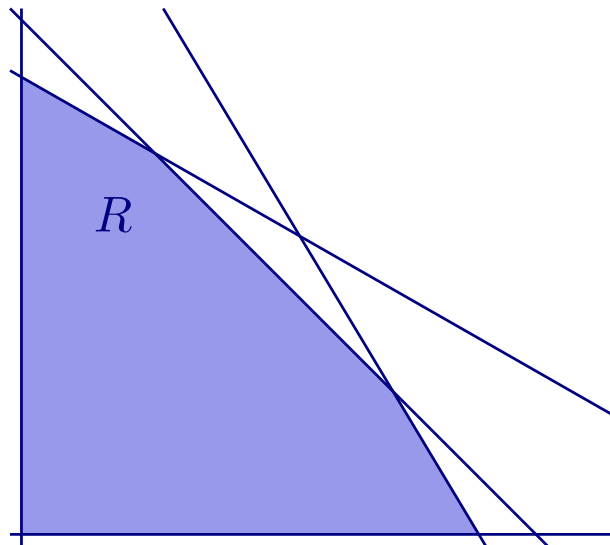


- ▷ Feasible region: $R = \{(x_b, x_c) \mid 4x_b + 4x_c \leq 45$
 $5x_b + 3x_c \leq 50$
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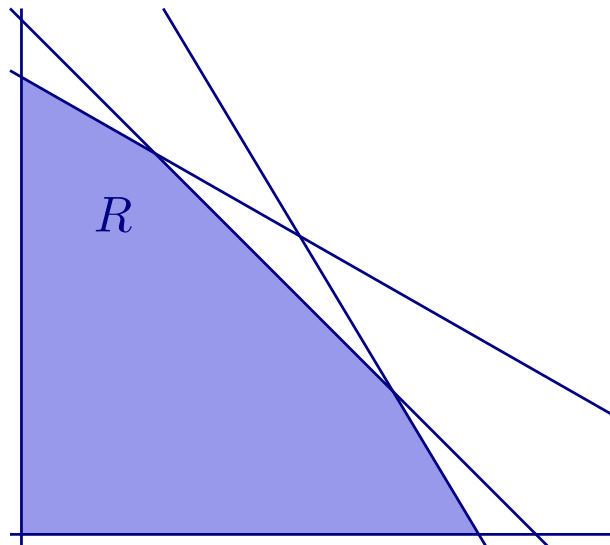
➔ is **convex**



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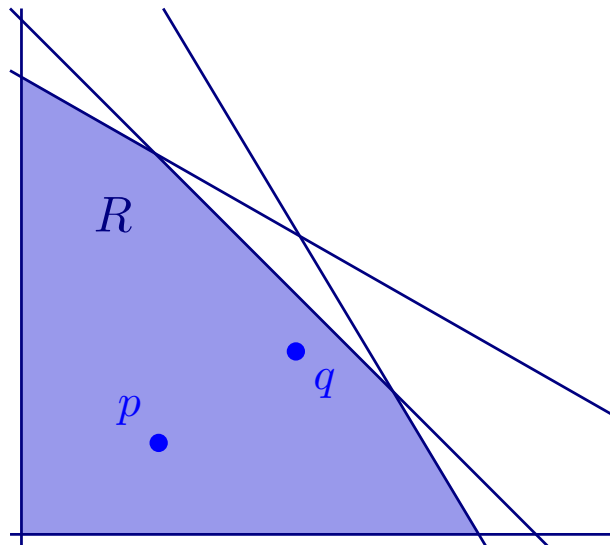
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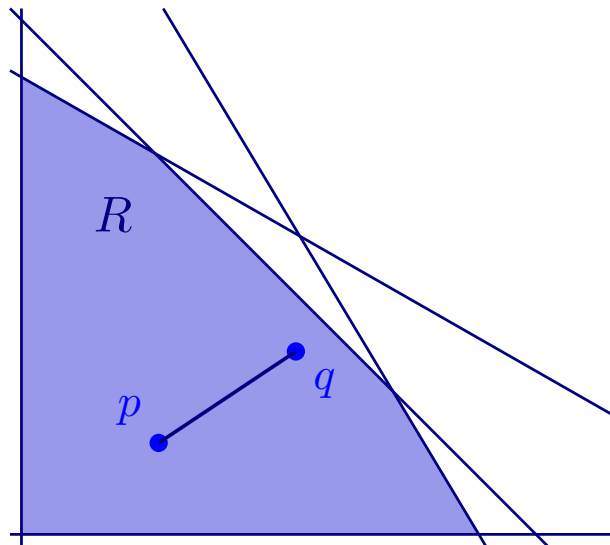
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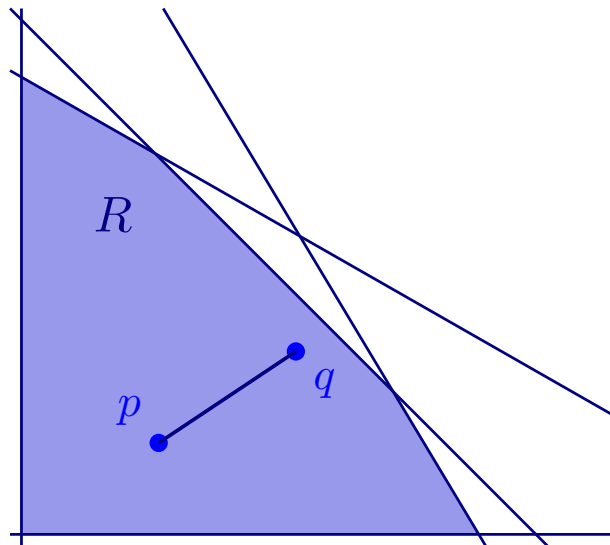


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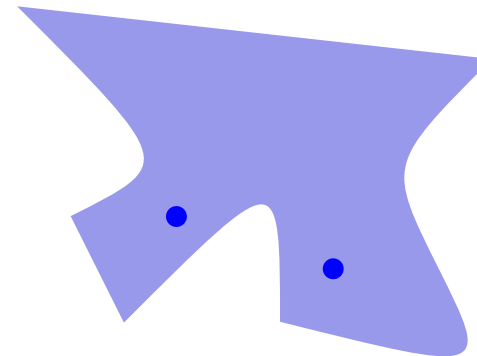
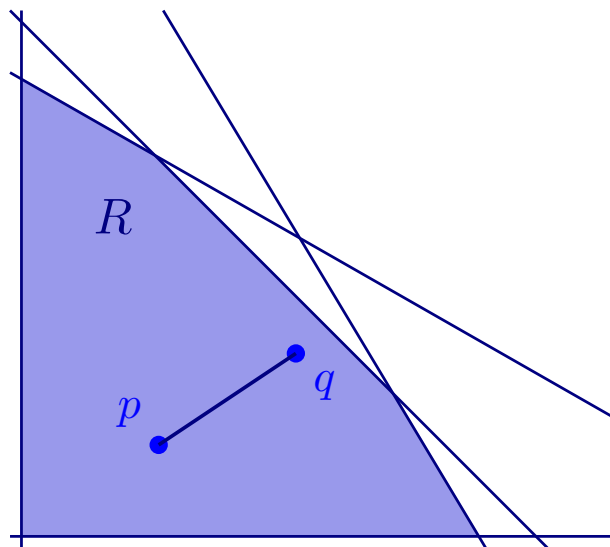


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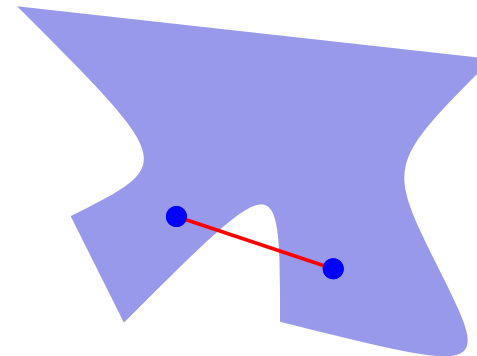
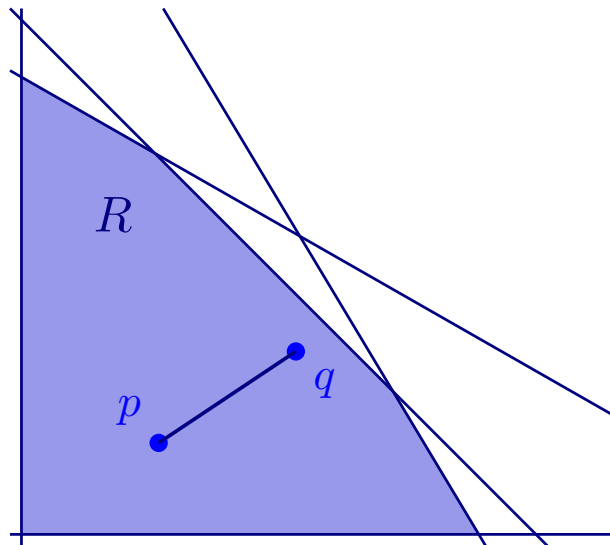


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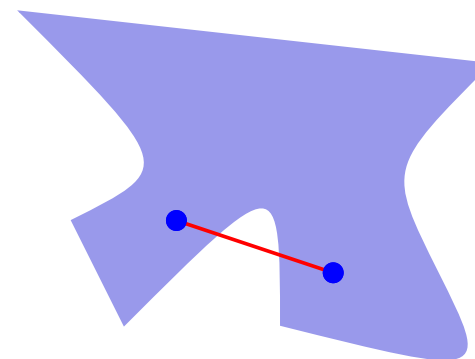
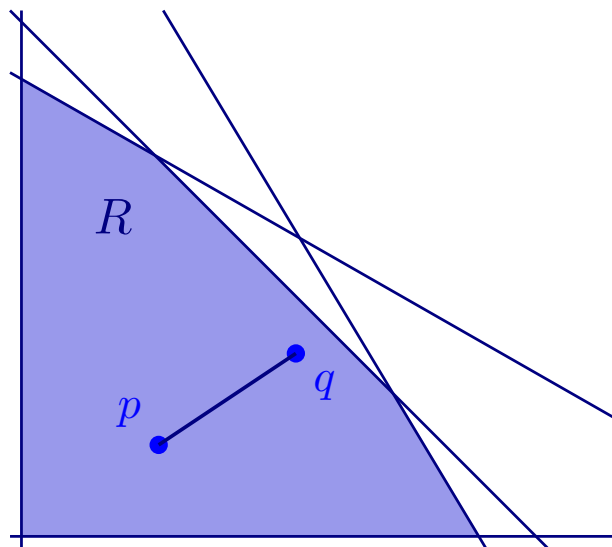


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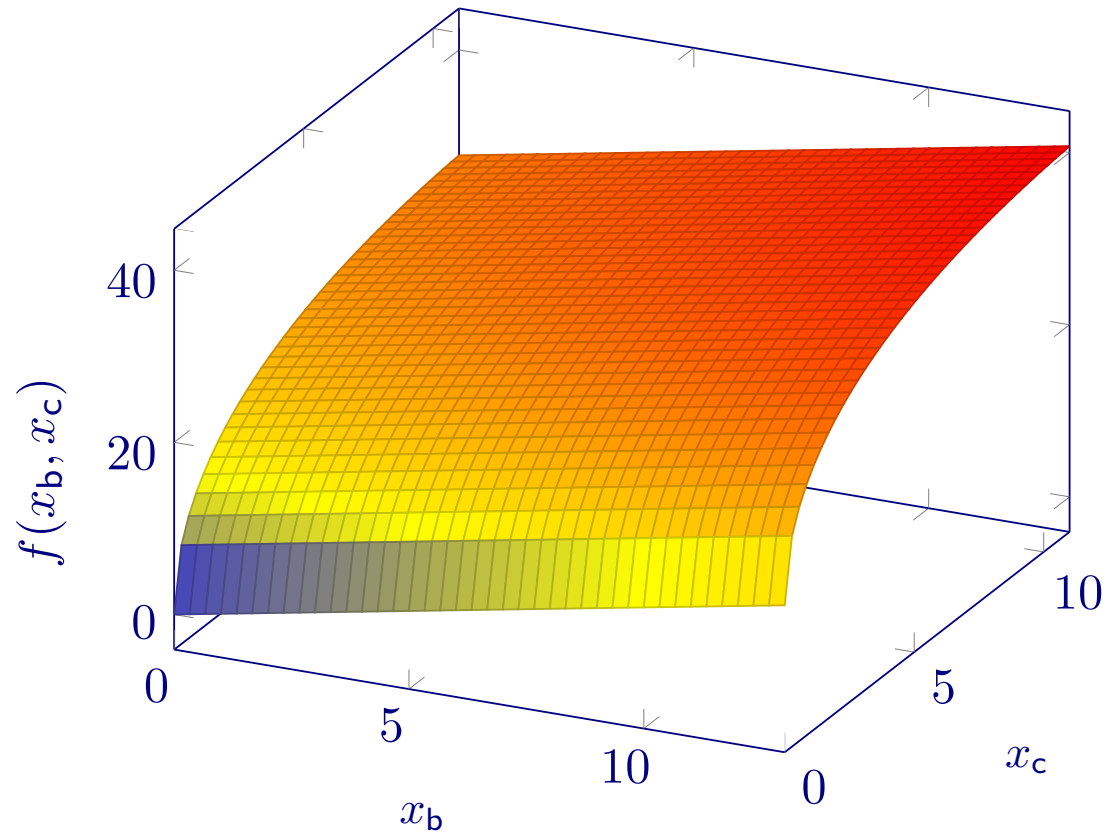
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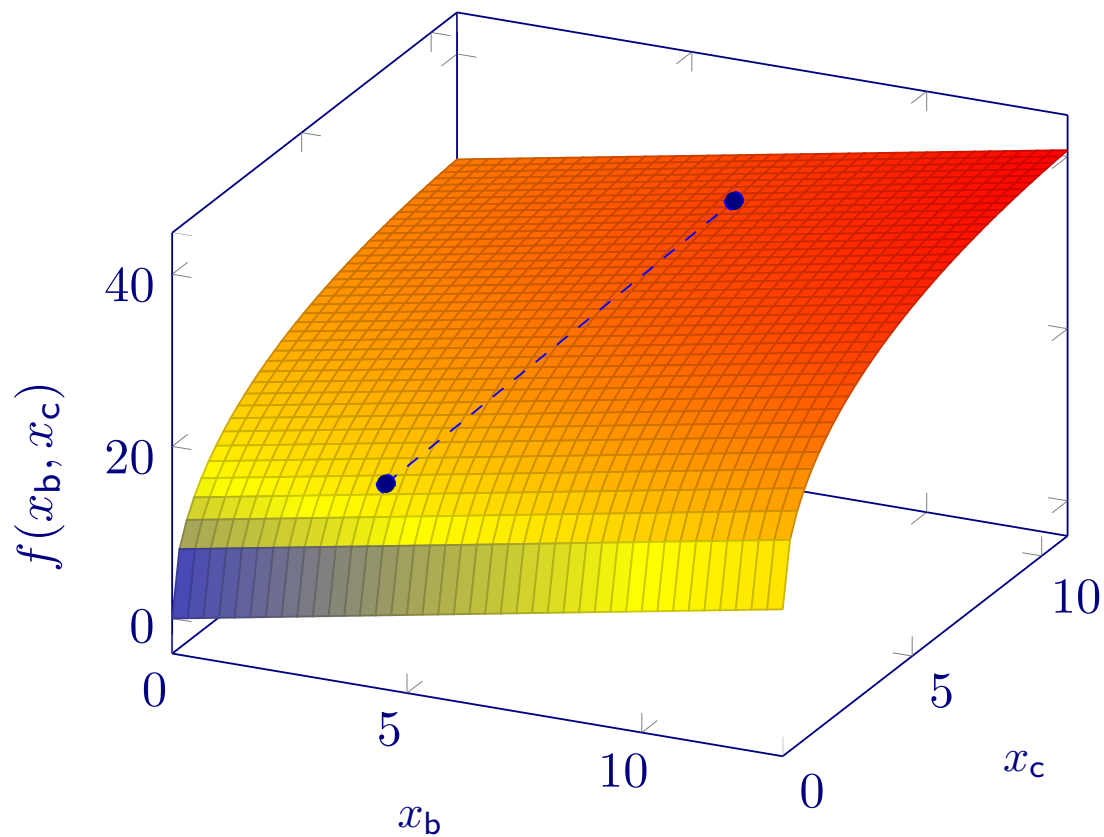


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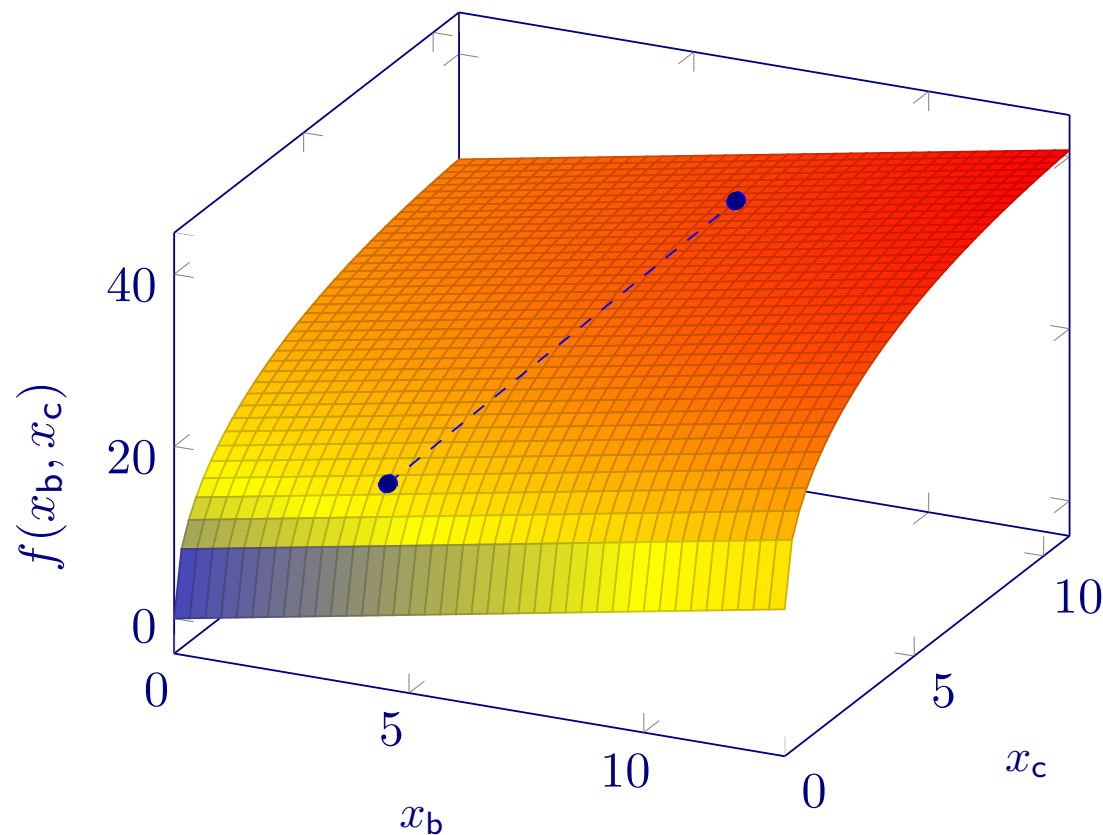
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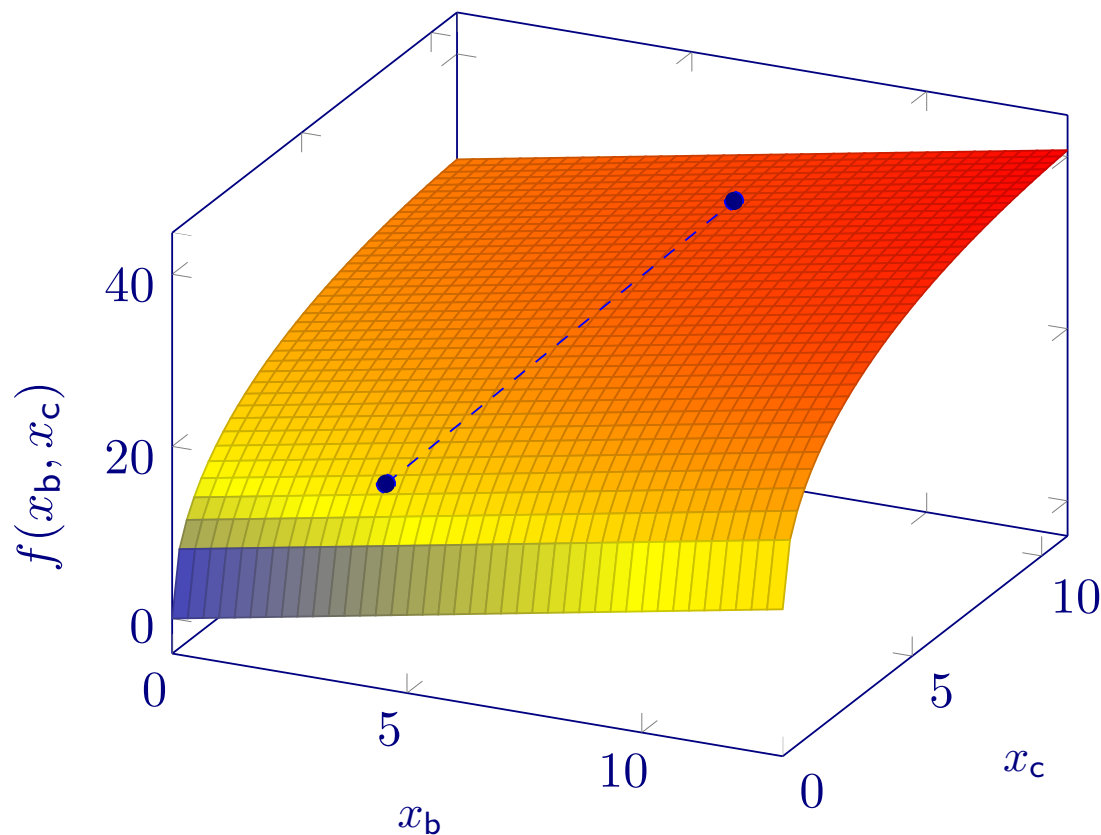


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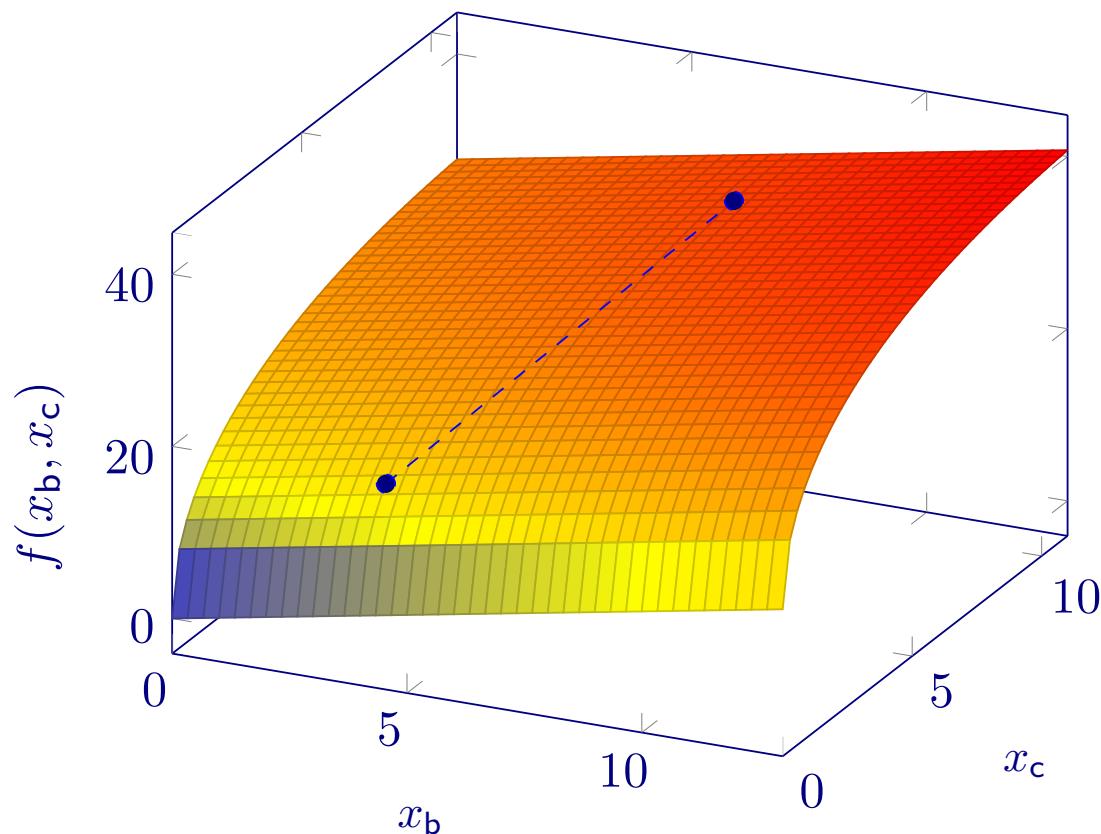
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- ▷ f is **concave**
i.e. every straight line between two points on the surface lies beneath the surface

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i.e. every straight line between two points on the surface lies beneath the surface

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 $f(\lambda p + (1 - \lambda)q) \geq \lambda f(p) + (1 - \lambda)f(q)$ for all $p, q \in R$ and $0 \leq \lambda \leq 1$.



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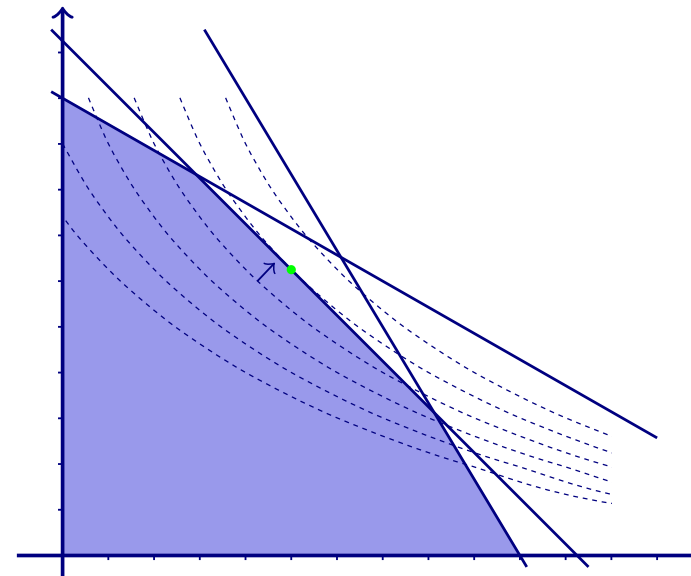
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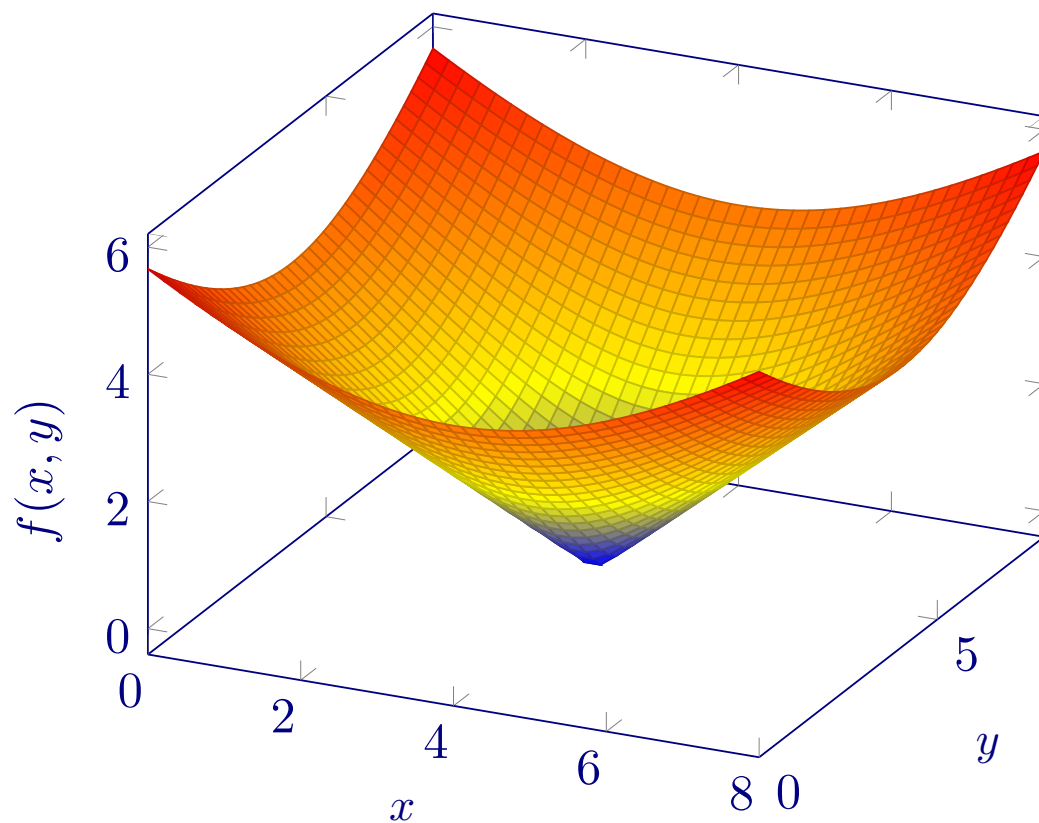
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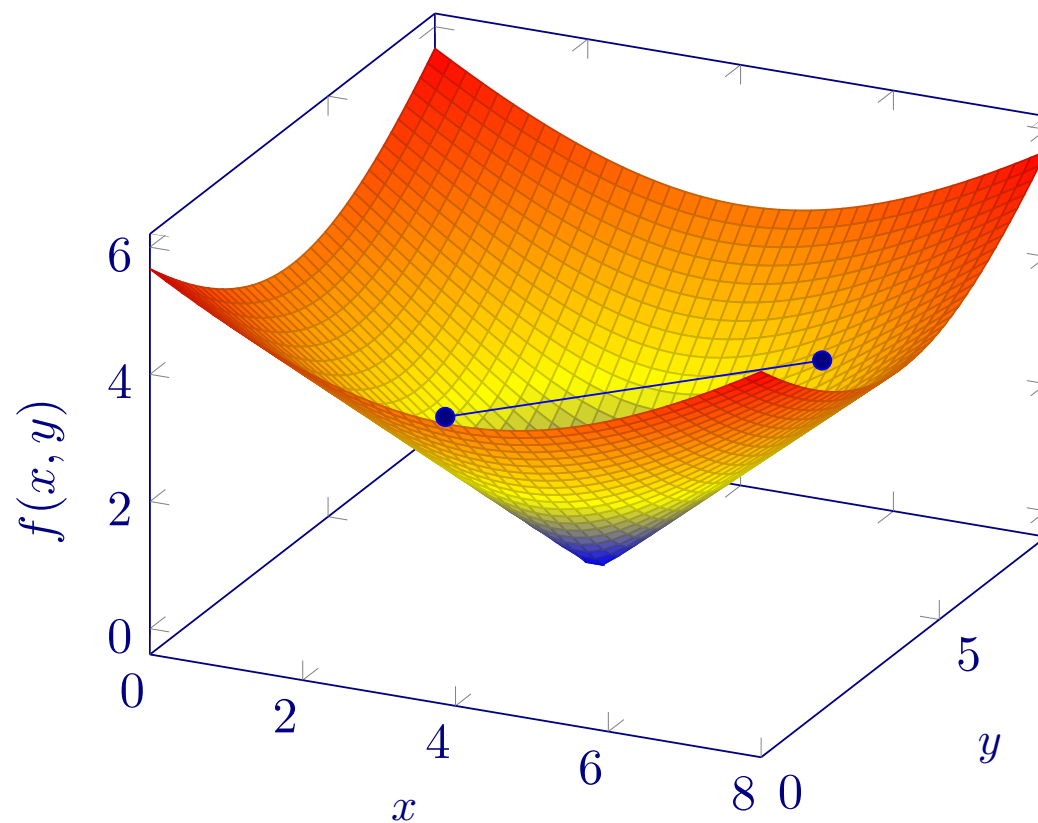
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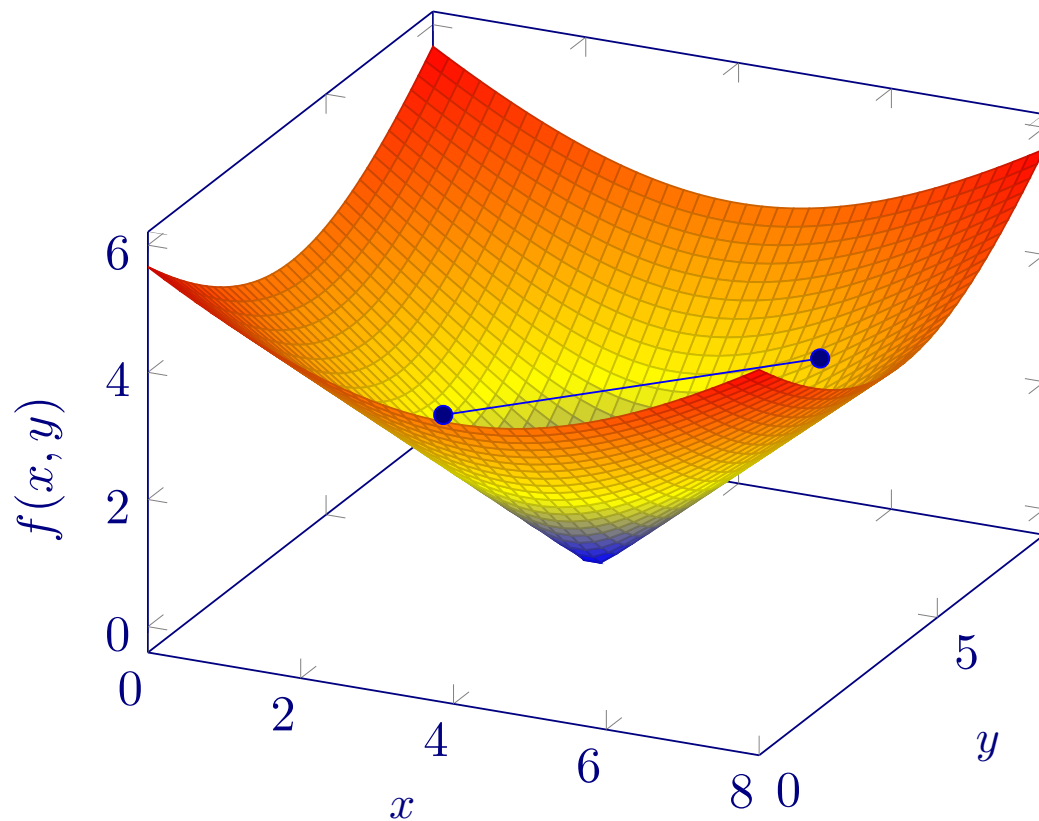
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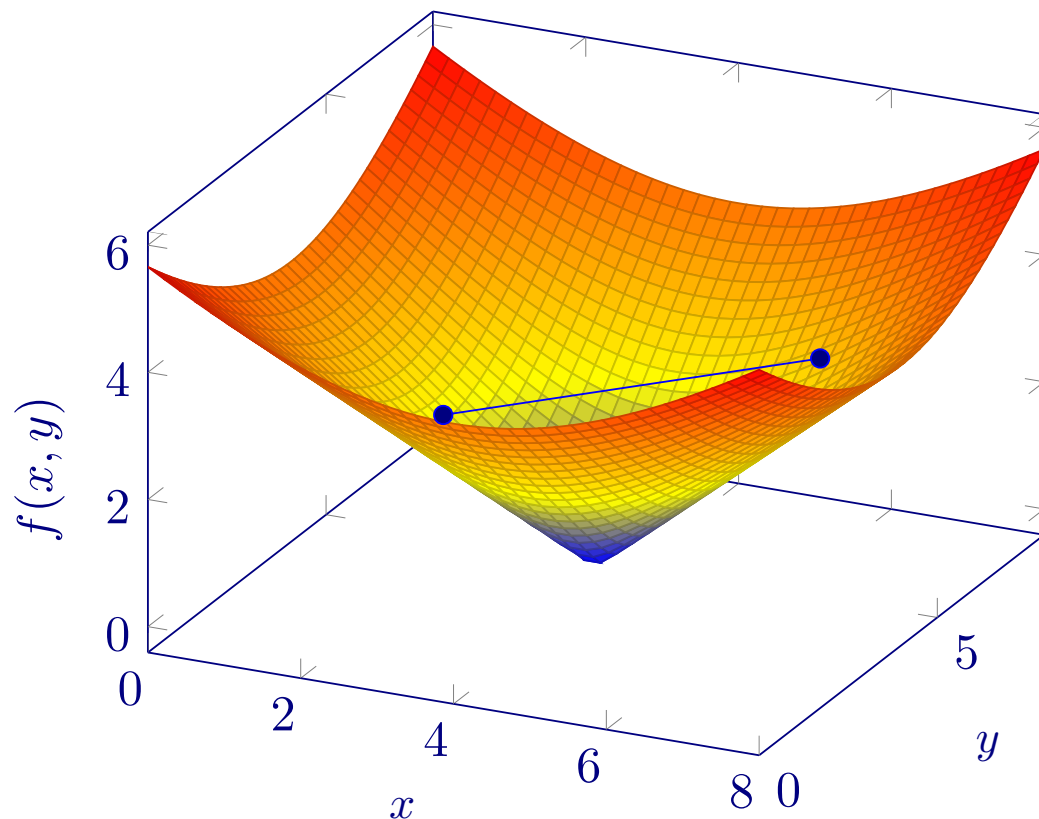


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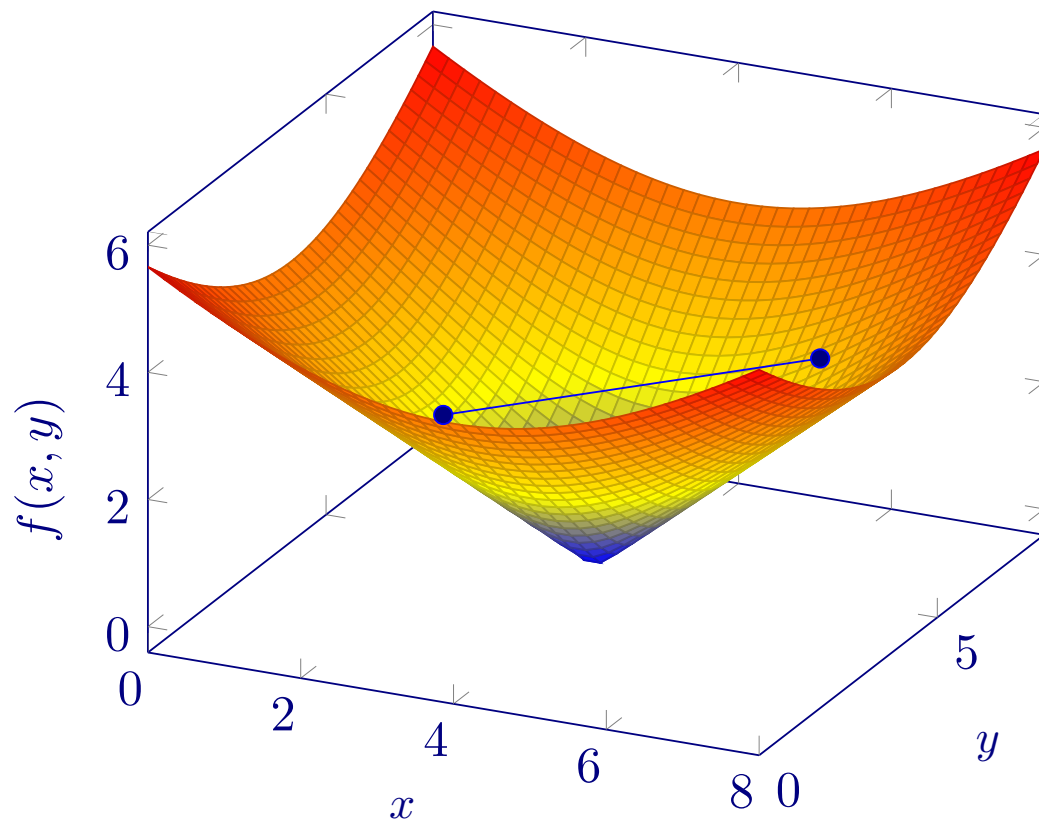
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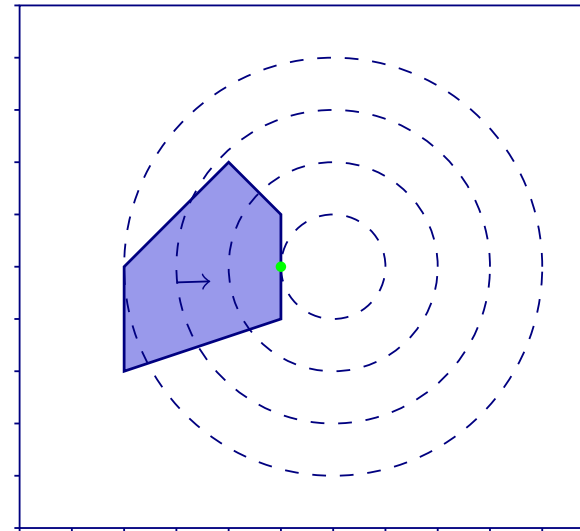
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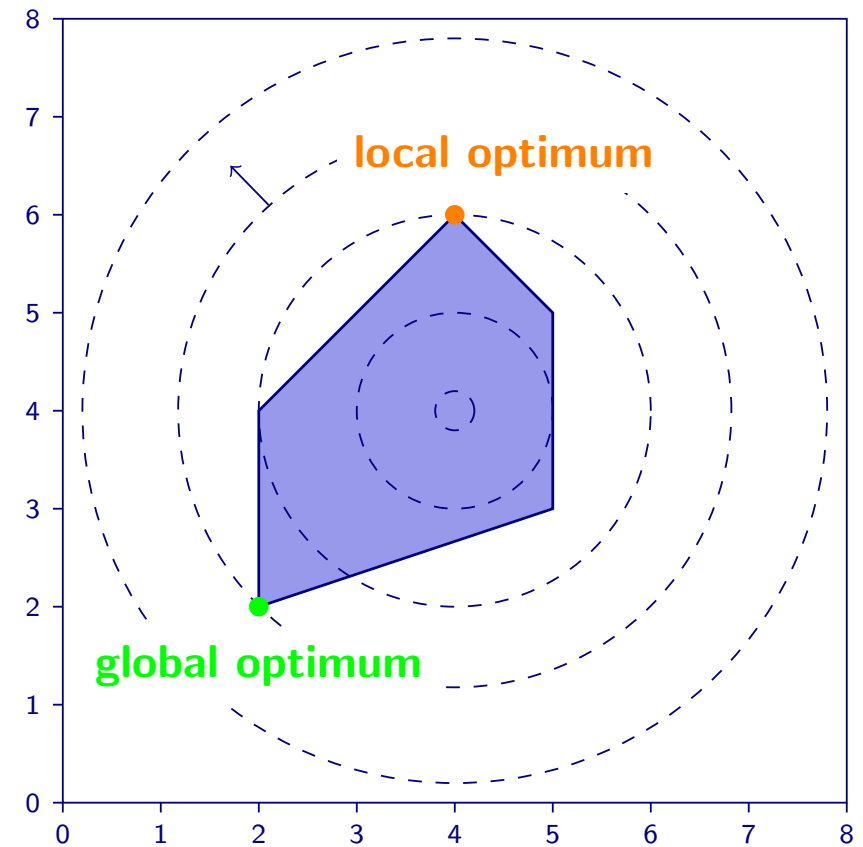
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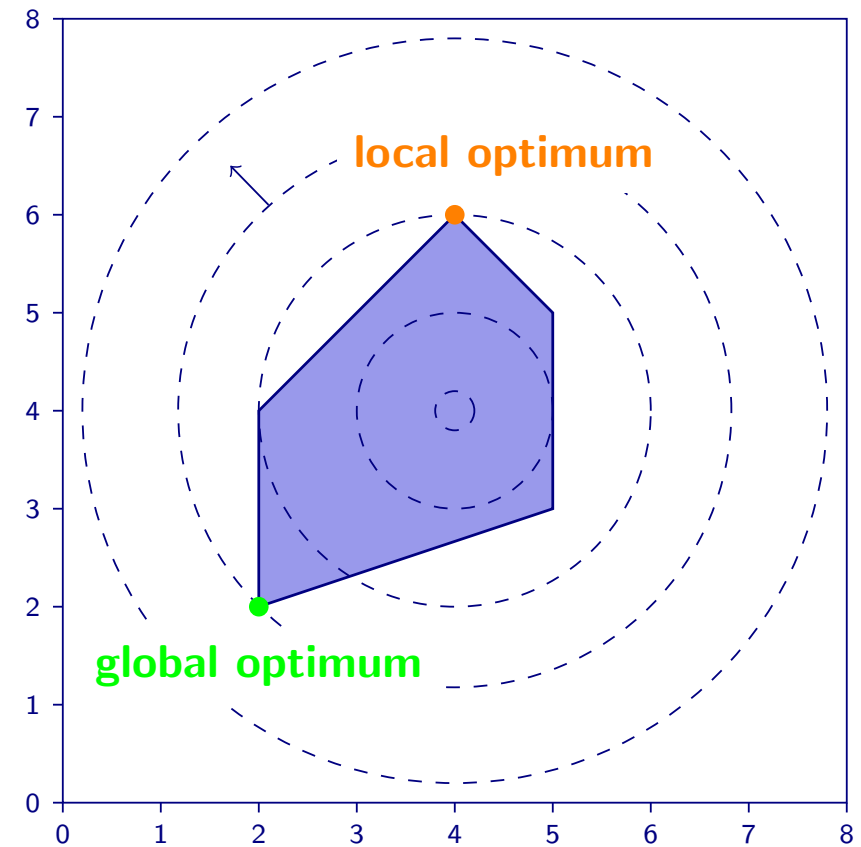
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$$\begin{aligned} \max \quad & \sqrt{(x-4)^2 + (y-4)^2} \\ \text{s.t.} \quad & x \geq 2 \\ & x \leq 5 \\ & -x + y \leq 2 \\ & x + y \leq 10 \\ & x - 3y \leq -4 \end{aligned}$$

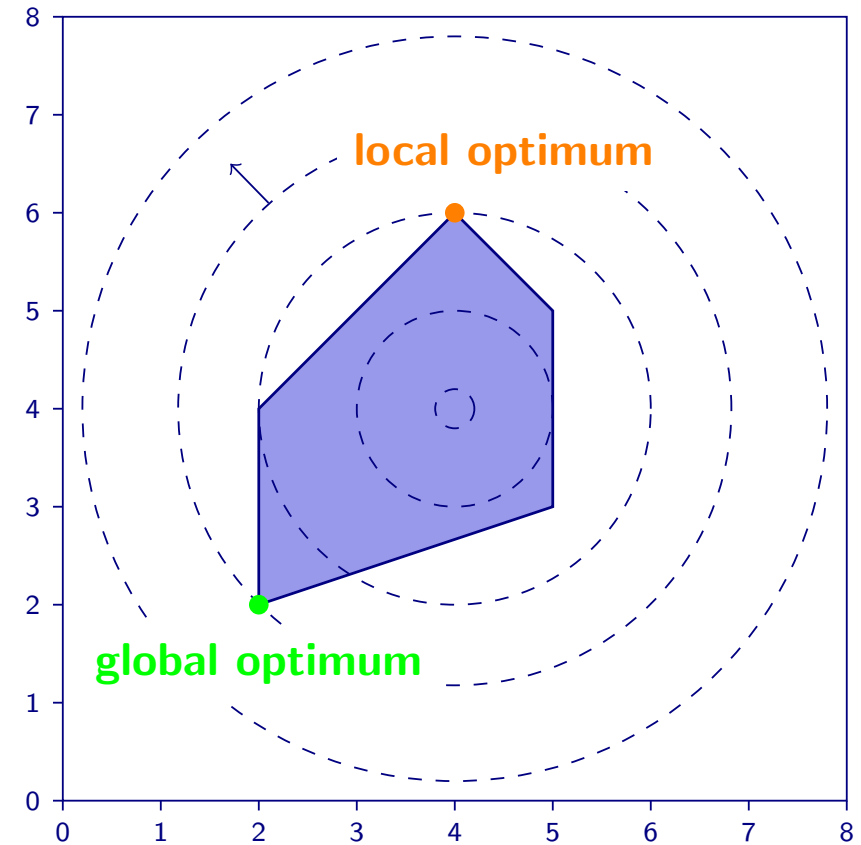


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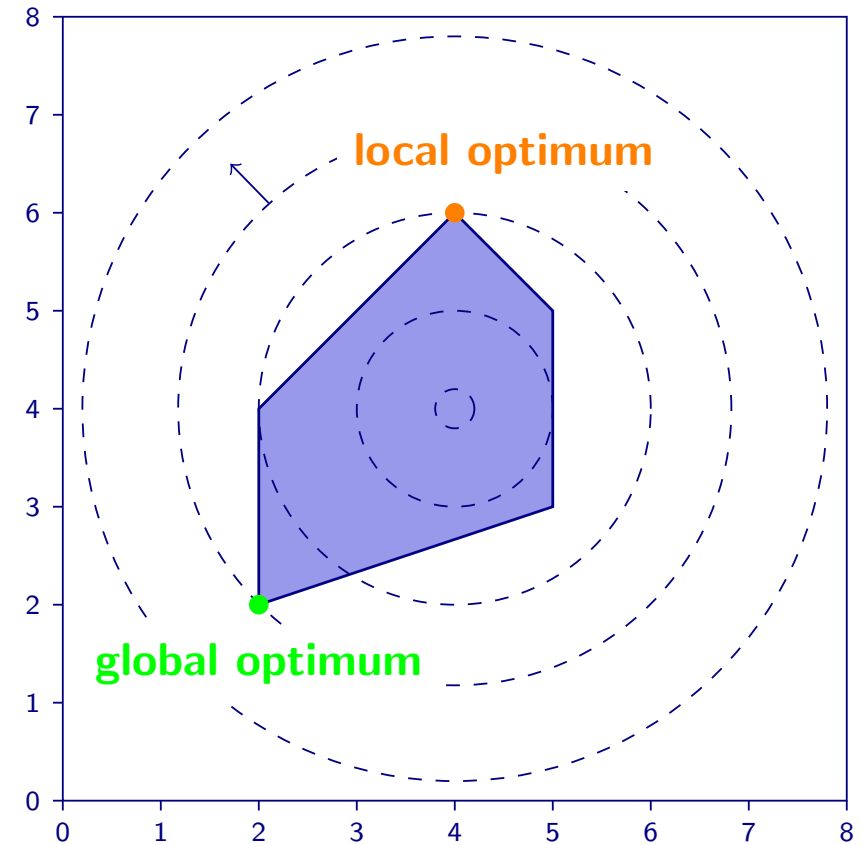
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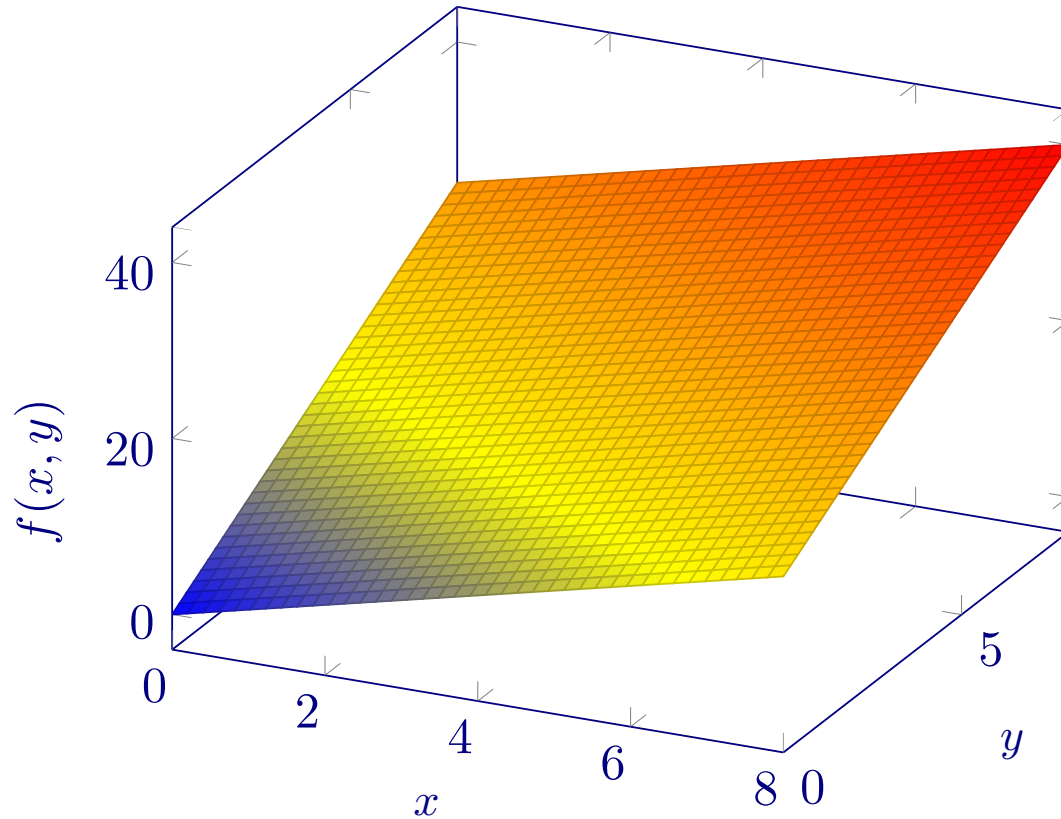
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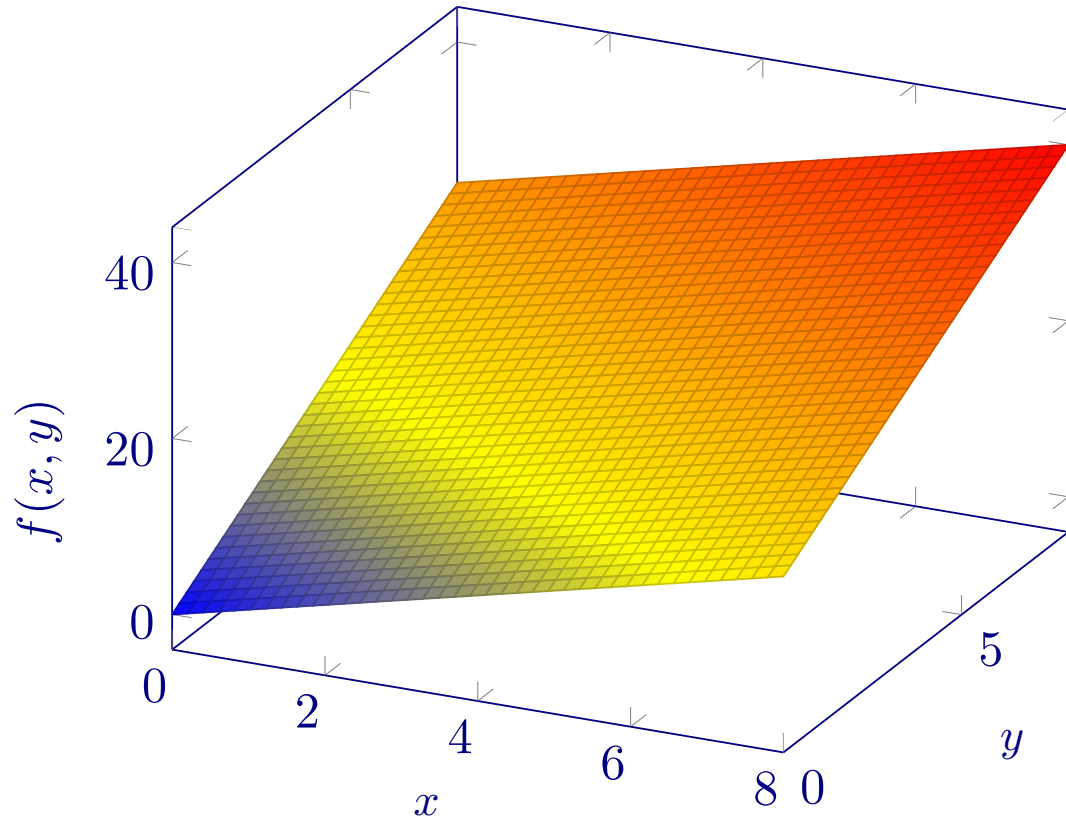


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- ▷ Maximization problem ☹️

- ▷ Linear objective function: $f(x, y) = 2x + 3y$

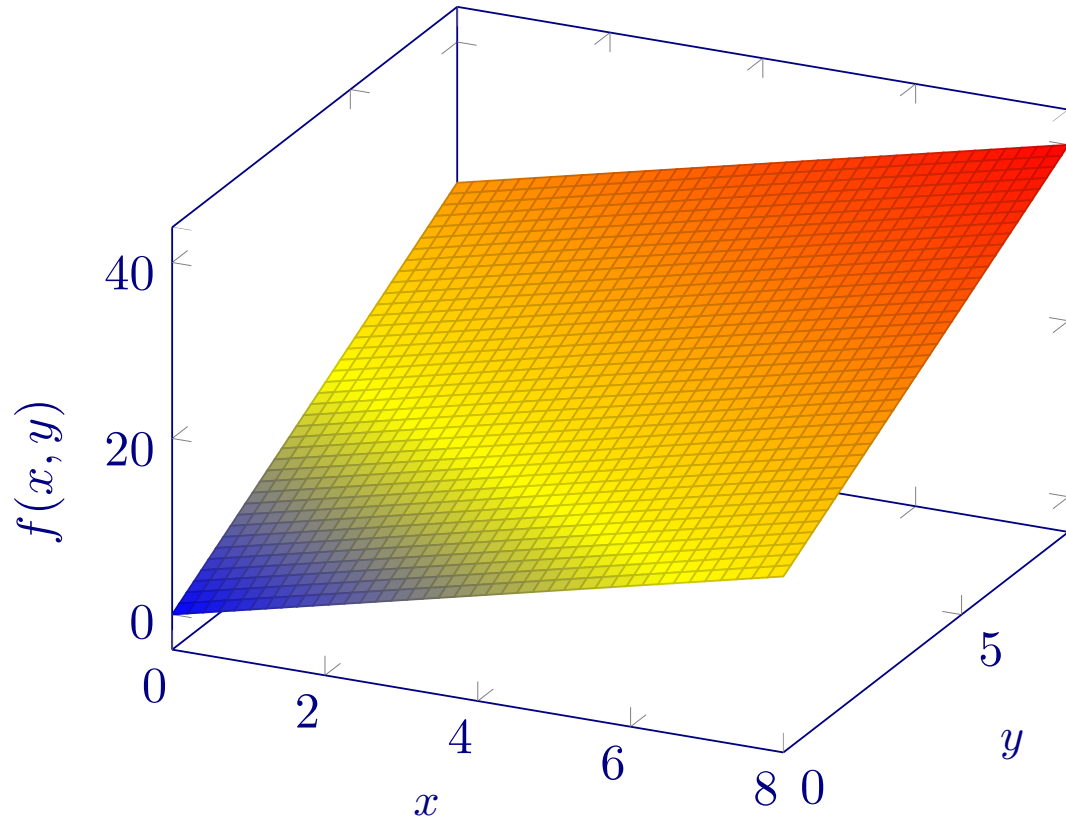


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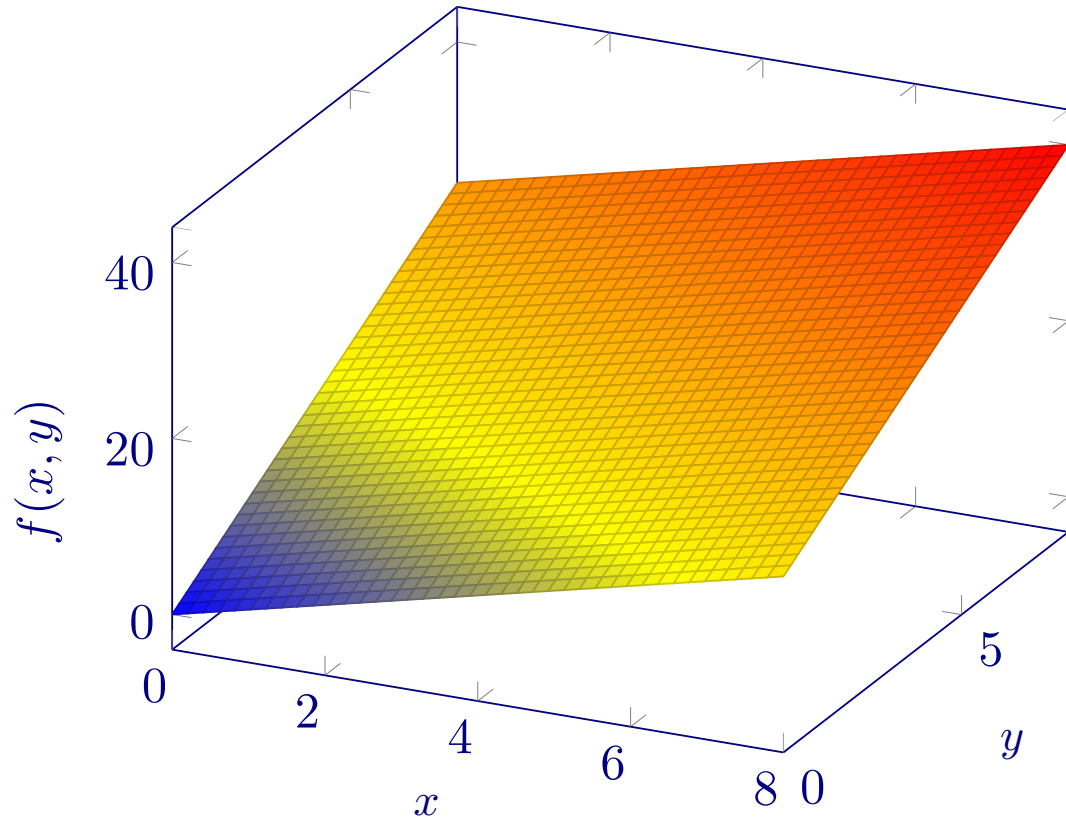
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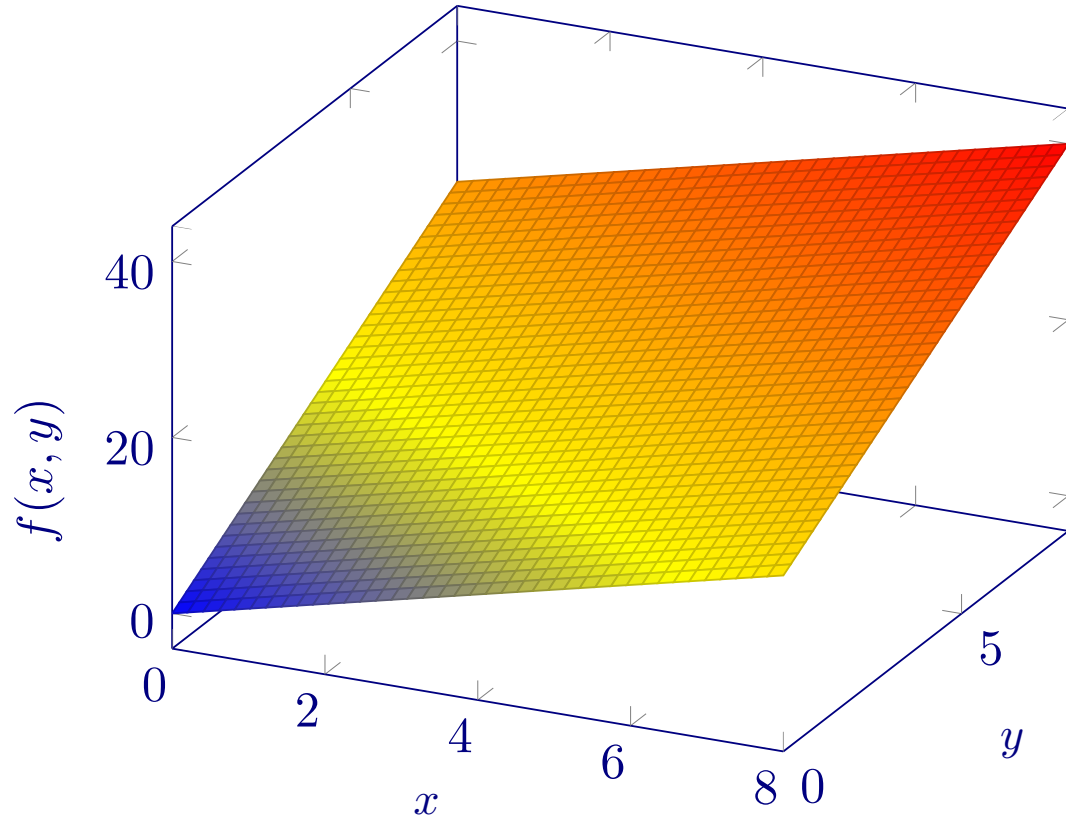
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➔ For linear programming local optima are always automatically global

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➔ Non-linear: minimize $\max_{k=1, \dots, \ell} f_k(x_1, \dots, x_n)$ (f_k are all linear)

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- ▶ Linear constraints and objective function to minimize is convex of the form

$$\sum_{i=1}^n c_i |x_i| \quad \text{with all } c_i \geq 0$$

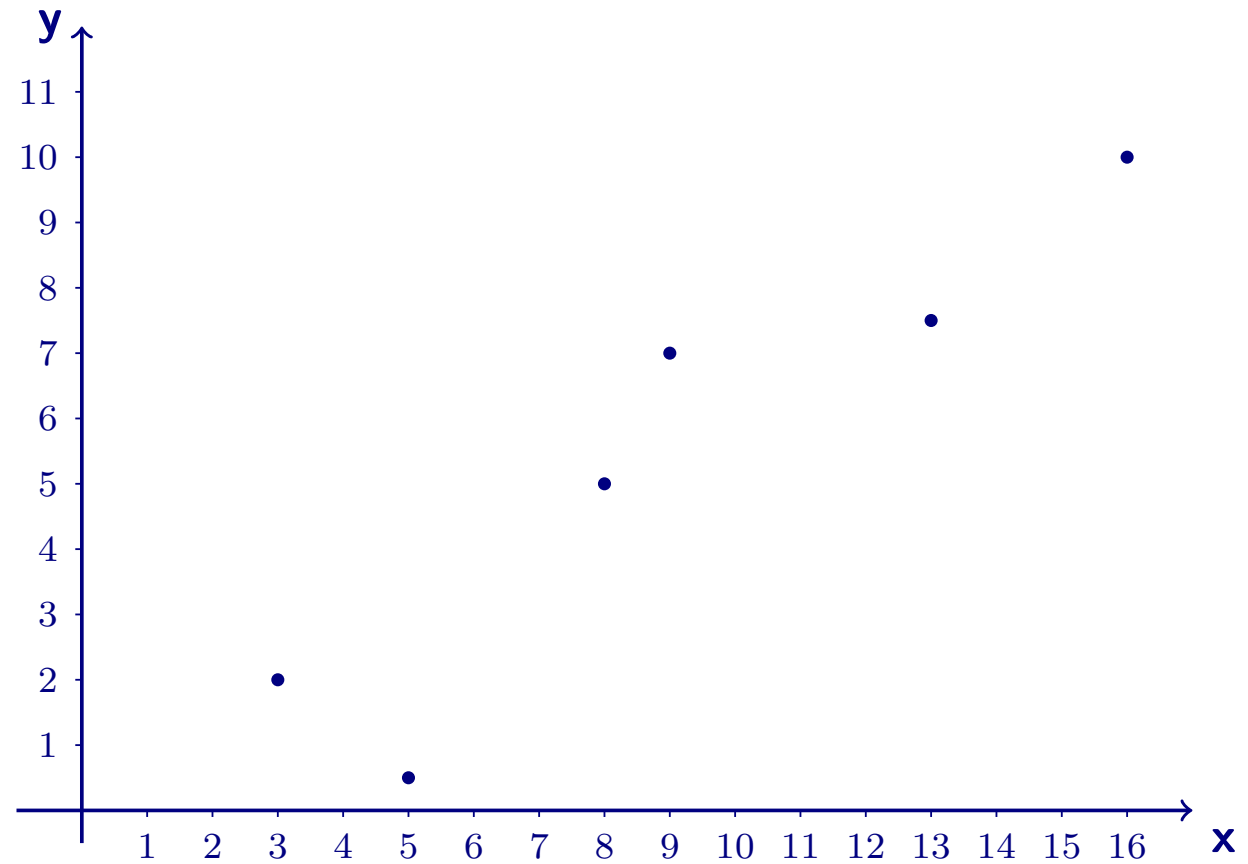
➔ Can be similarly rewritten into linear constraints



- ▶ Task: given a set of data points, find a line that “fits best” into the point set!

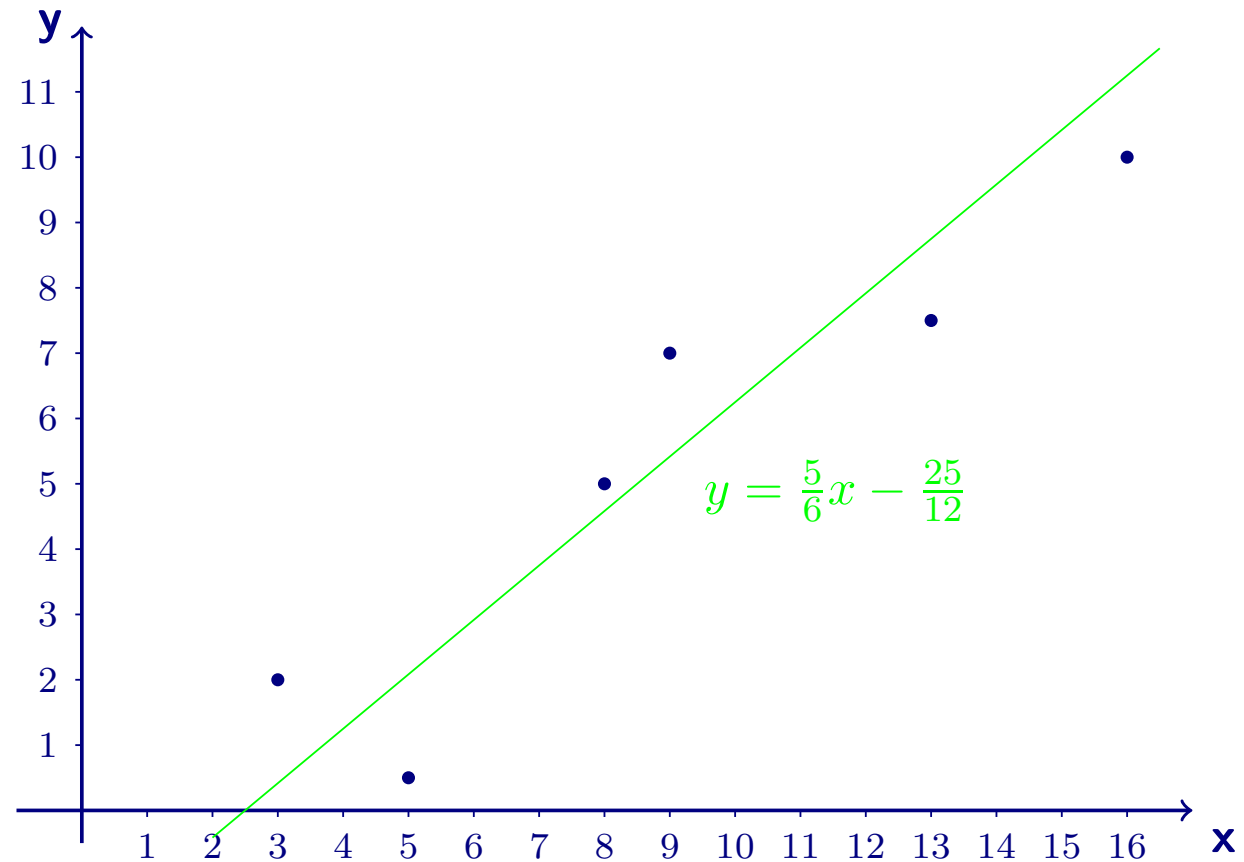
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x	y
3	2
5	0.5
8	5
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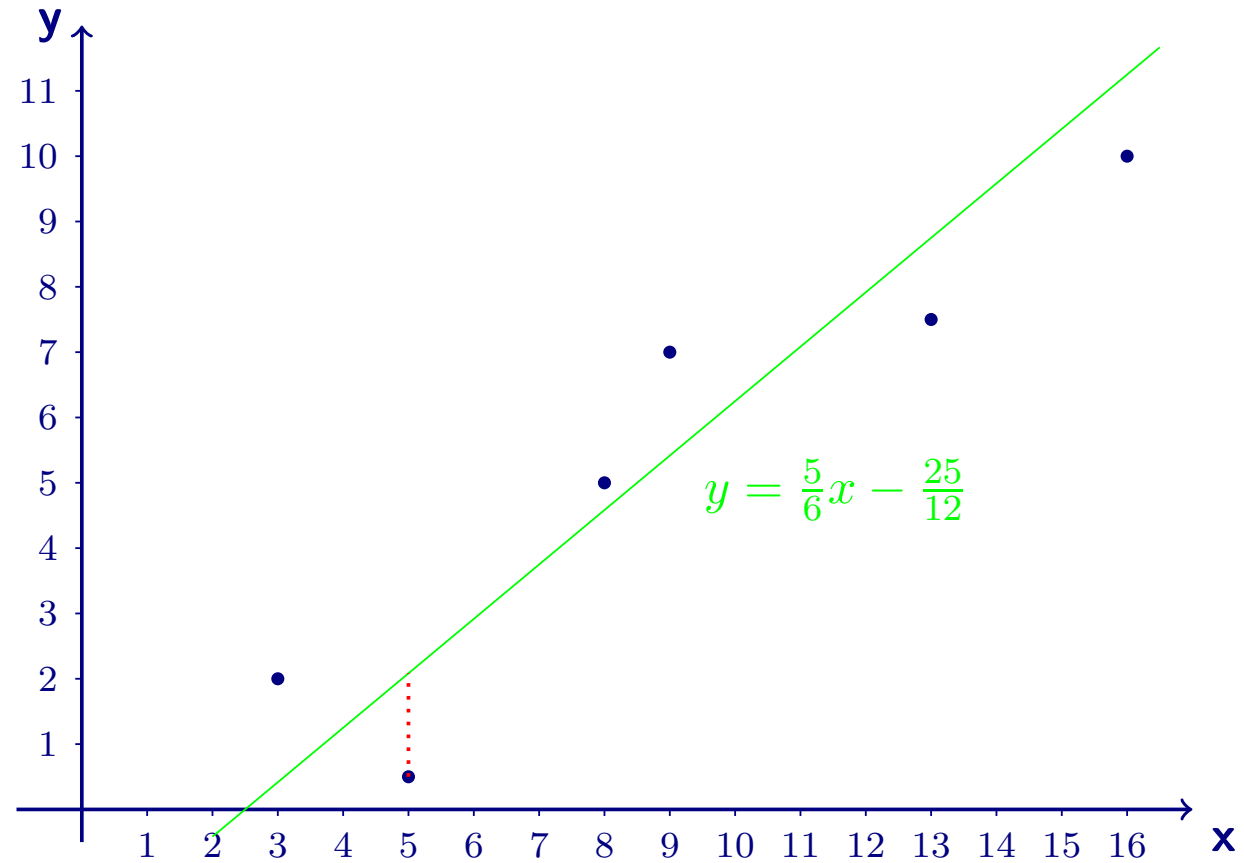
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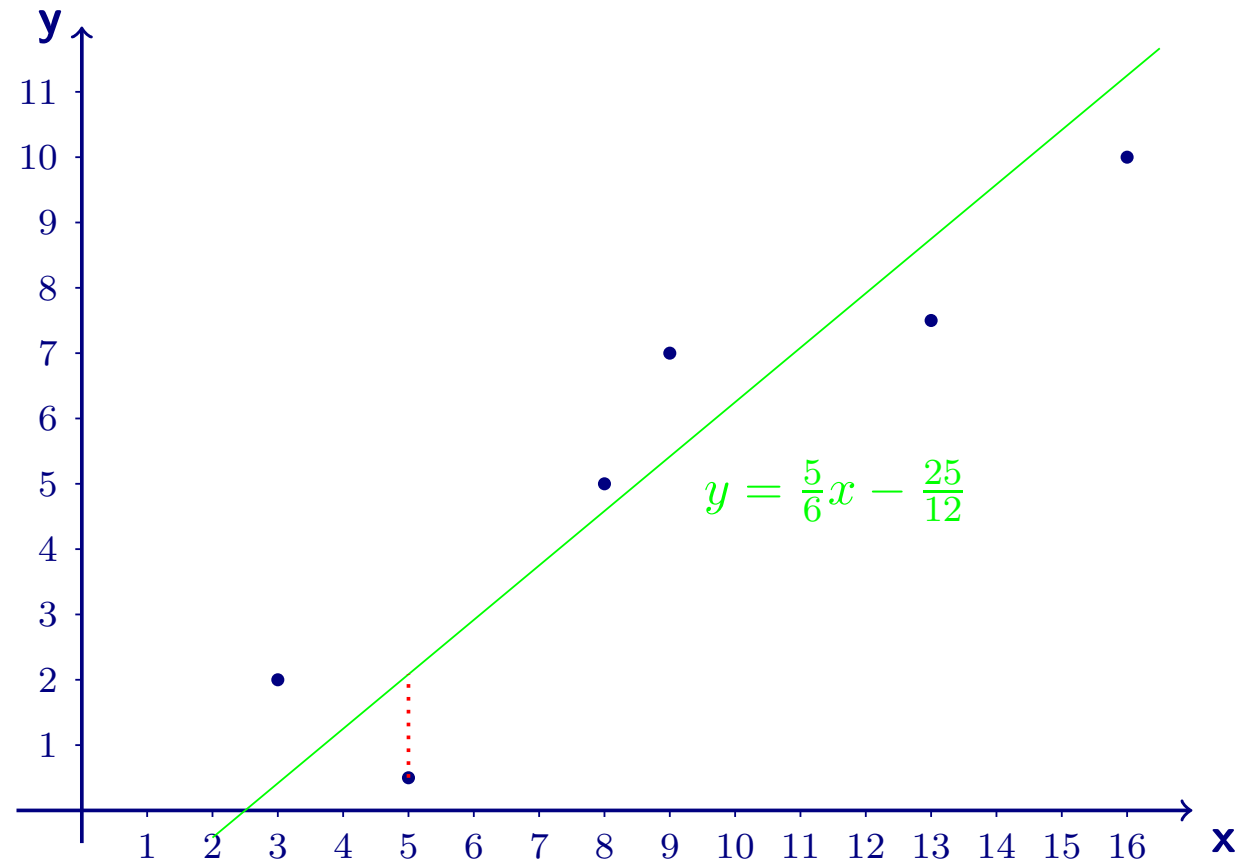
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- ➔ Minimize the largest occurring vertical distance between the wanted line and the data points!

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no constraints

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 &\text{subject to} && ax_i + b - y_i \leq z && (1 \leq i \leq n) \\
 &&& -ax_i - b + y_i \leq z && (1 \leq i \leq n) \\
 &&& a, b, z \in \mathbb{R}
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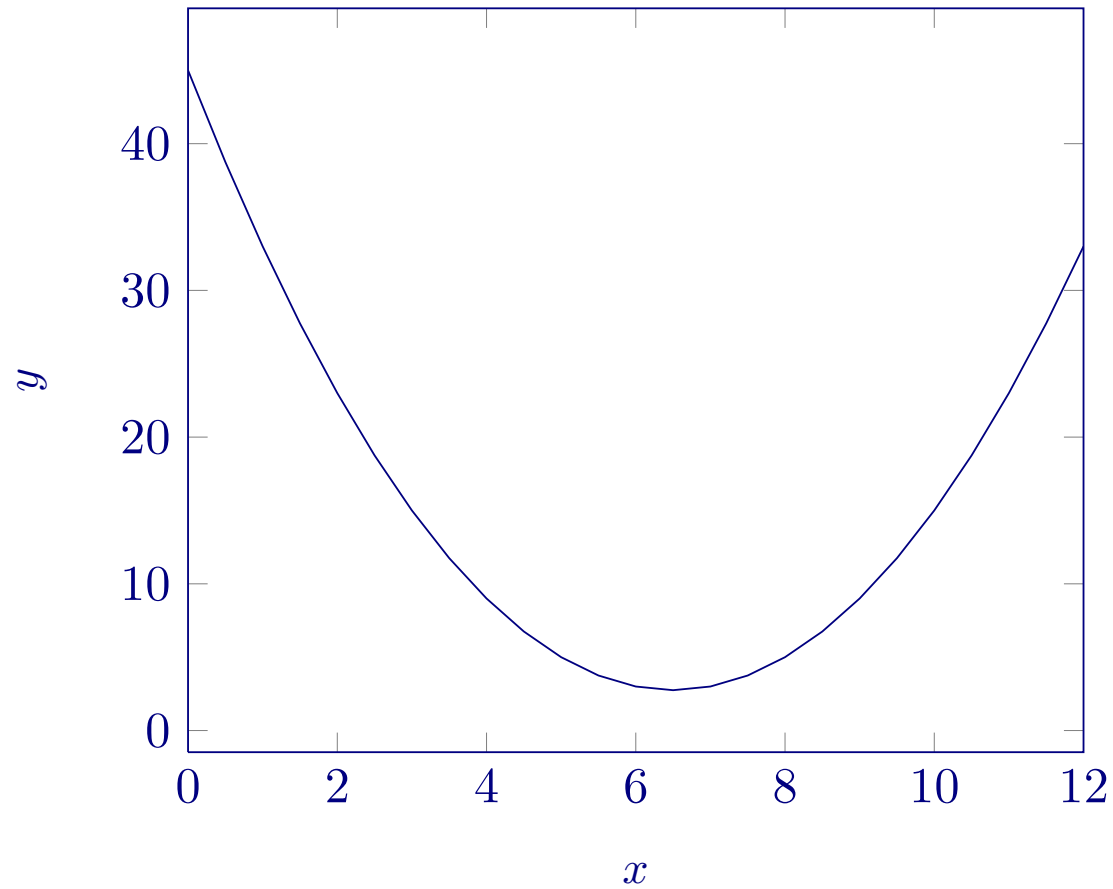
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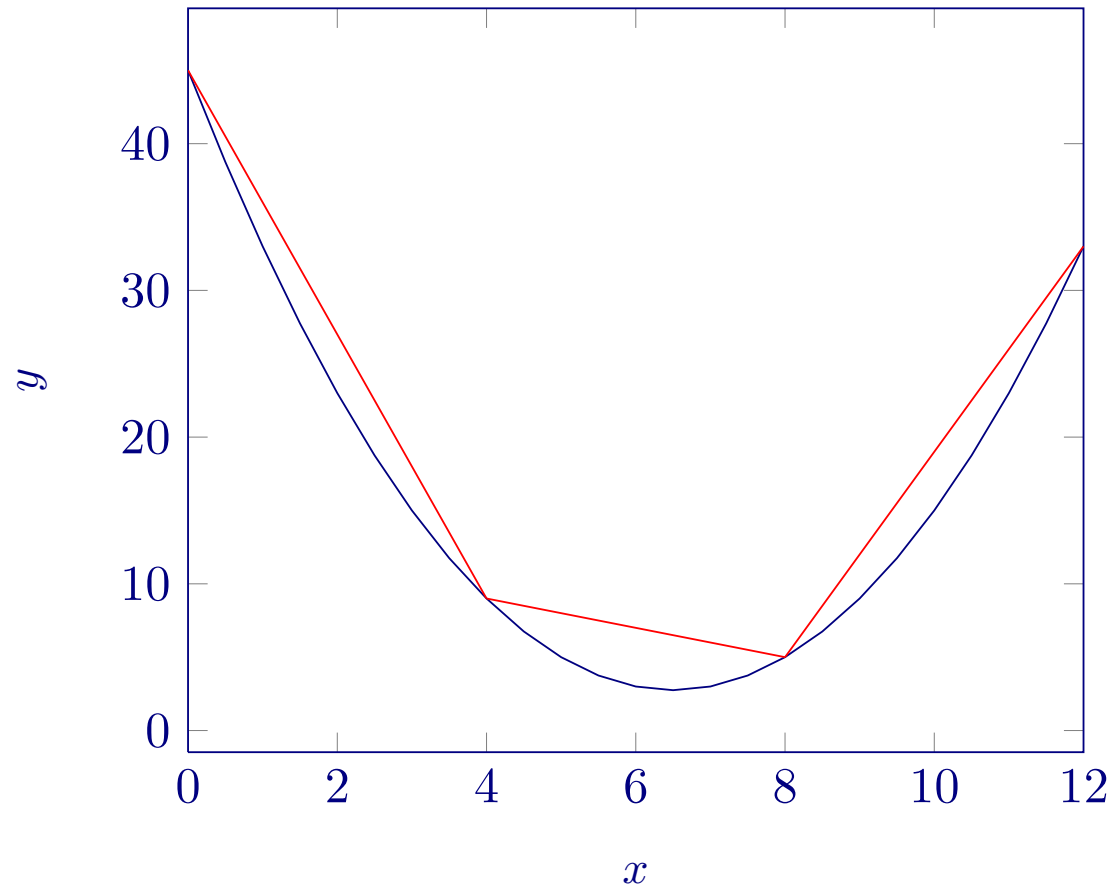
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- ▷ Variants: minimize sum of distances, square of distances, fit a higher-order curve

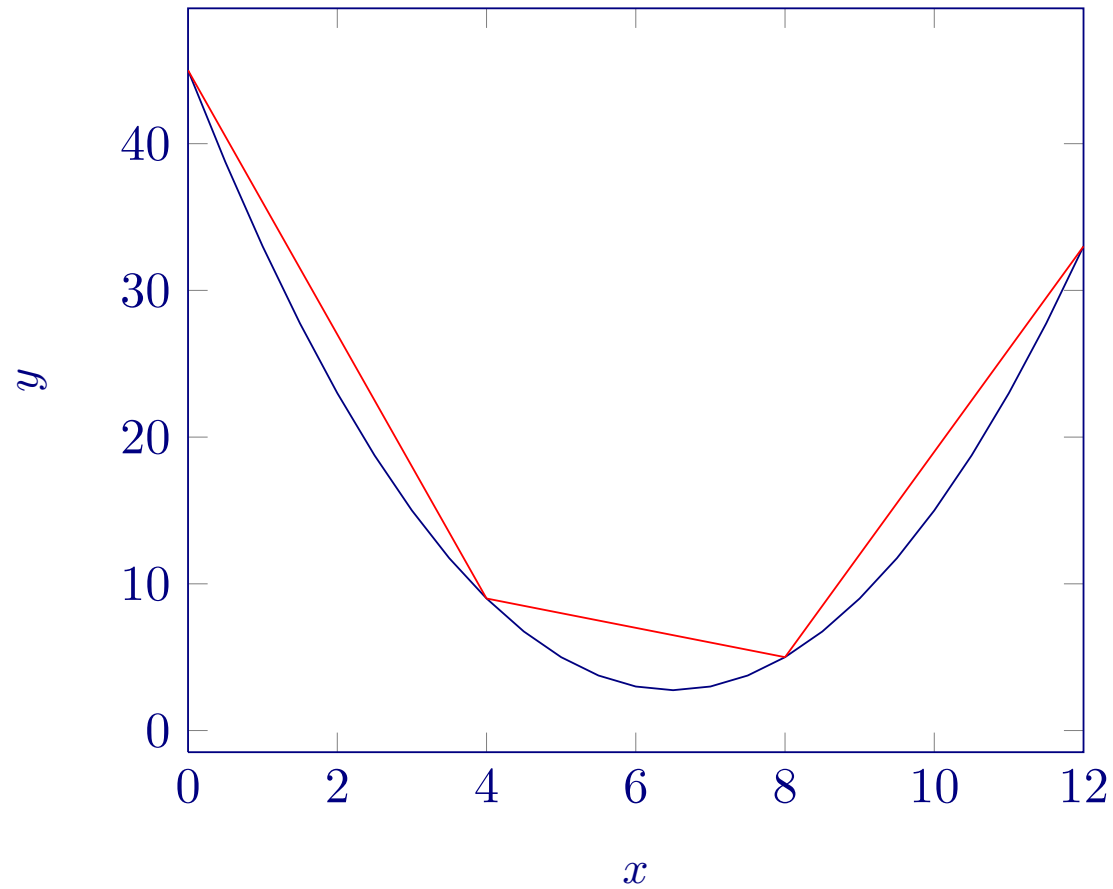
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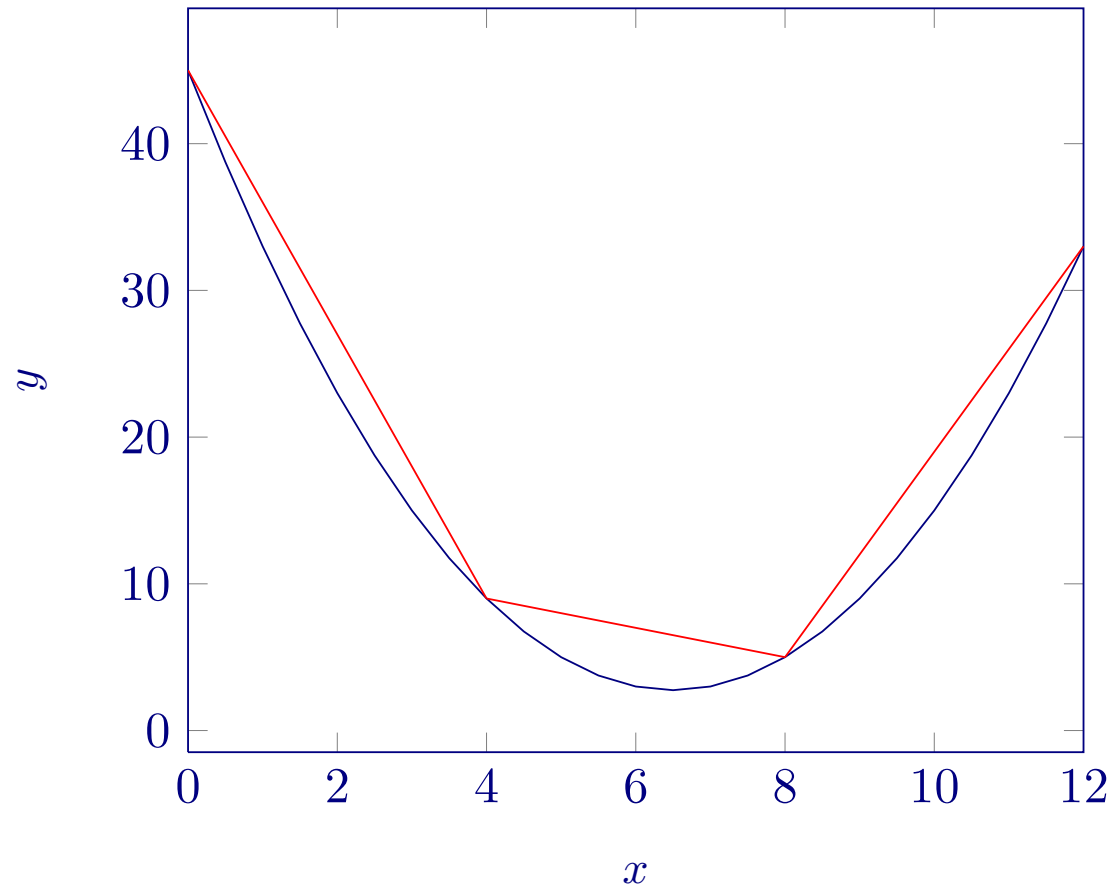


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▷ Variables $\lambda_1, \lambda_2, \lambda_3, \lambda_4 \geq 0$

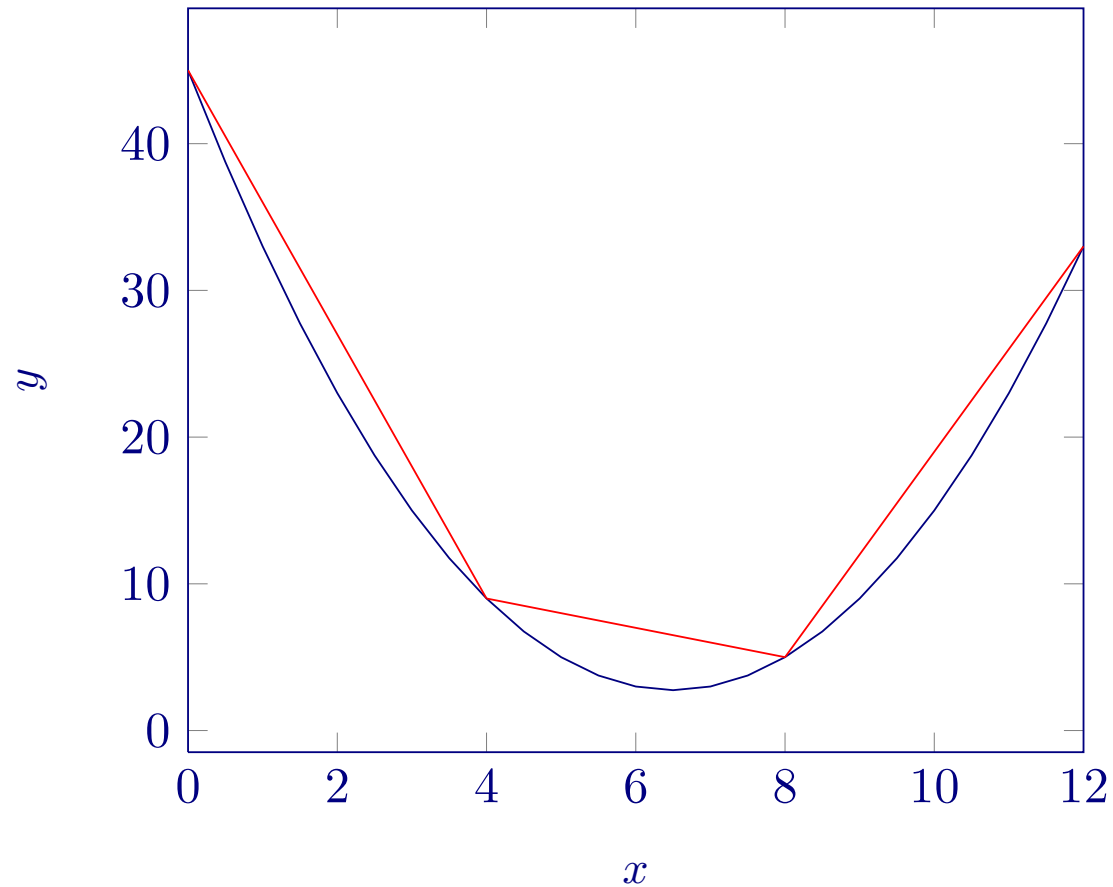
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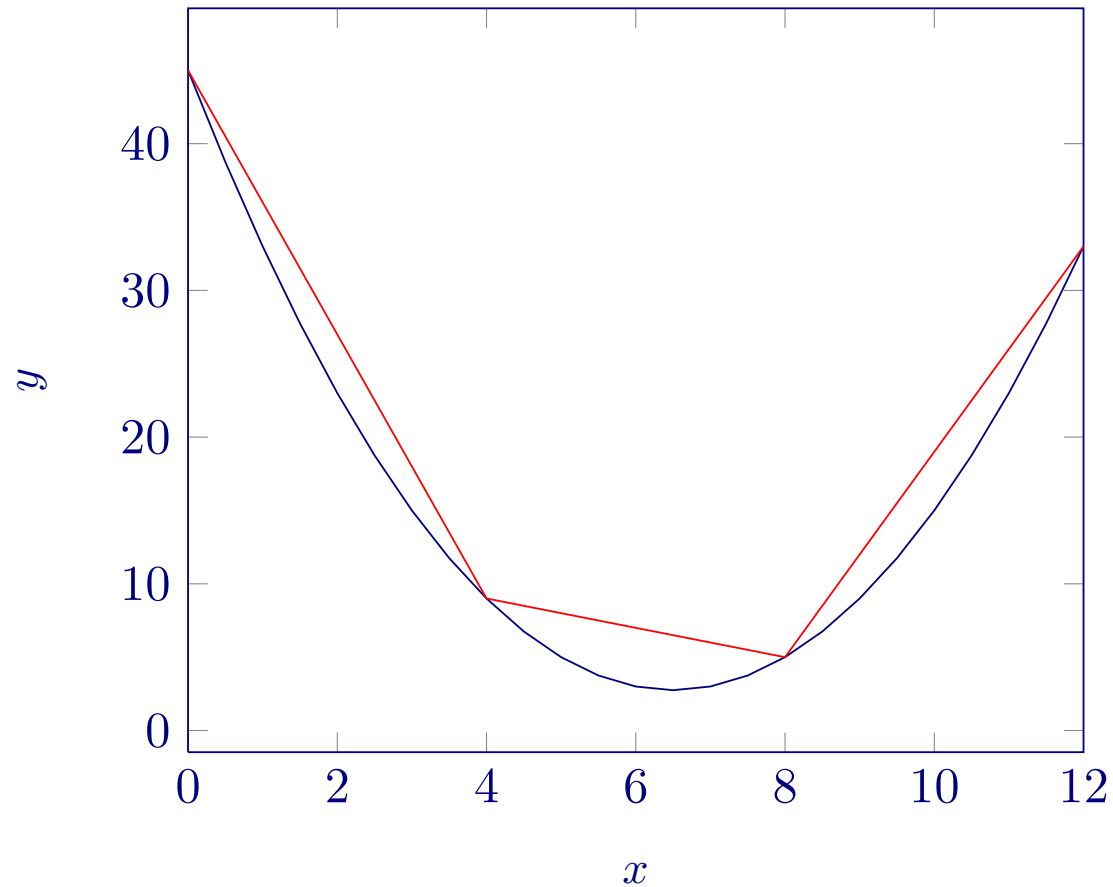
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$$y = 45 \cdot \lambda_1 + 9 \cdot \lambda_2 + 5 \cdot \lambda_3 + 33 \cdot \lambda_4$$

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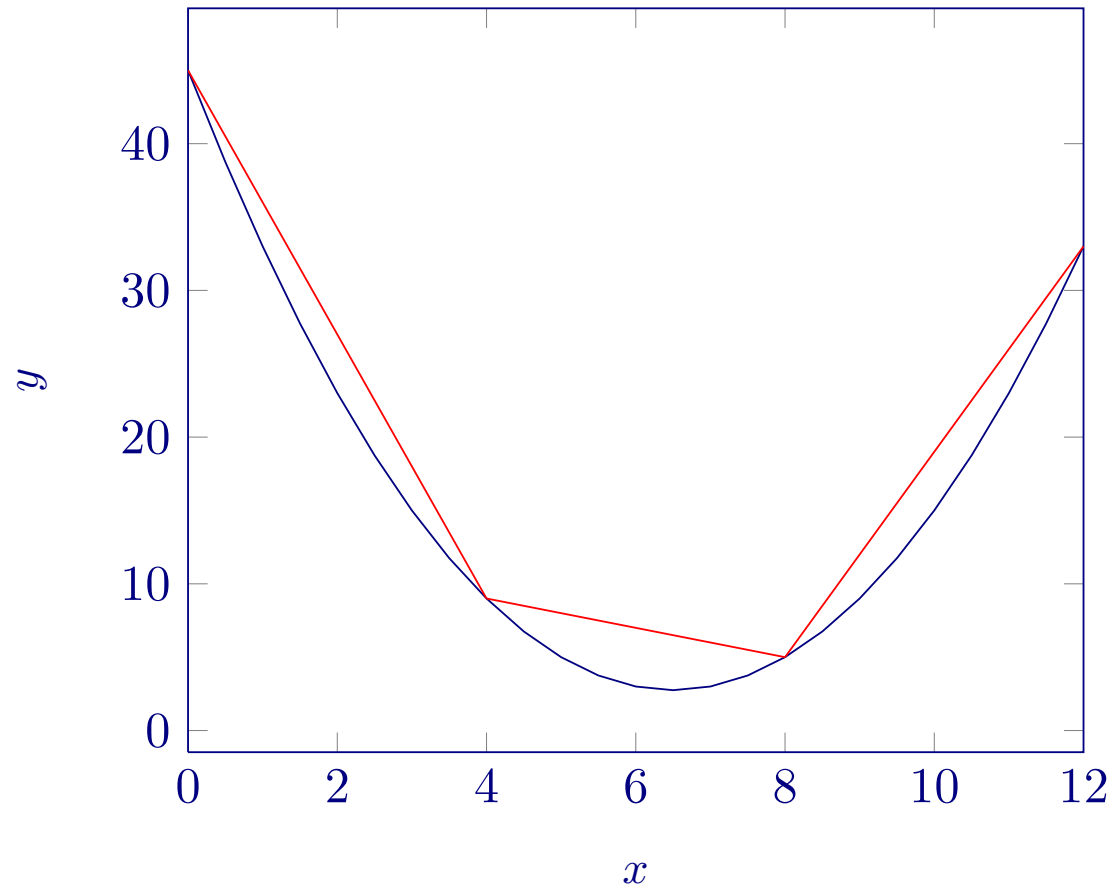
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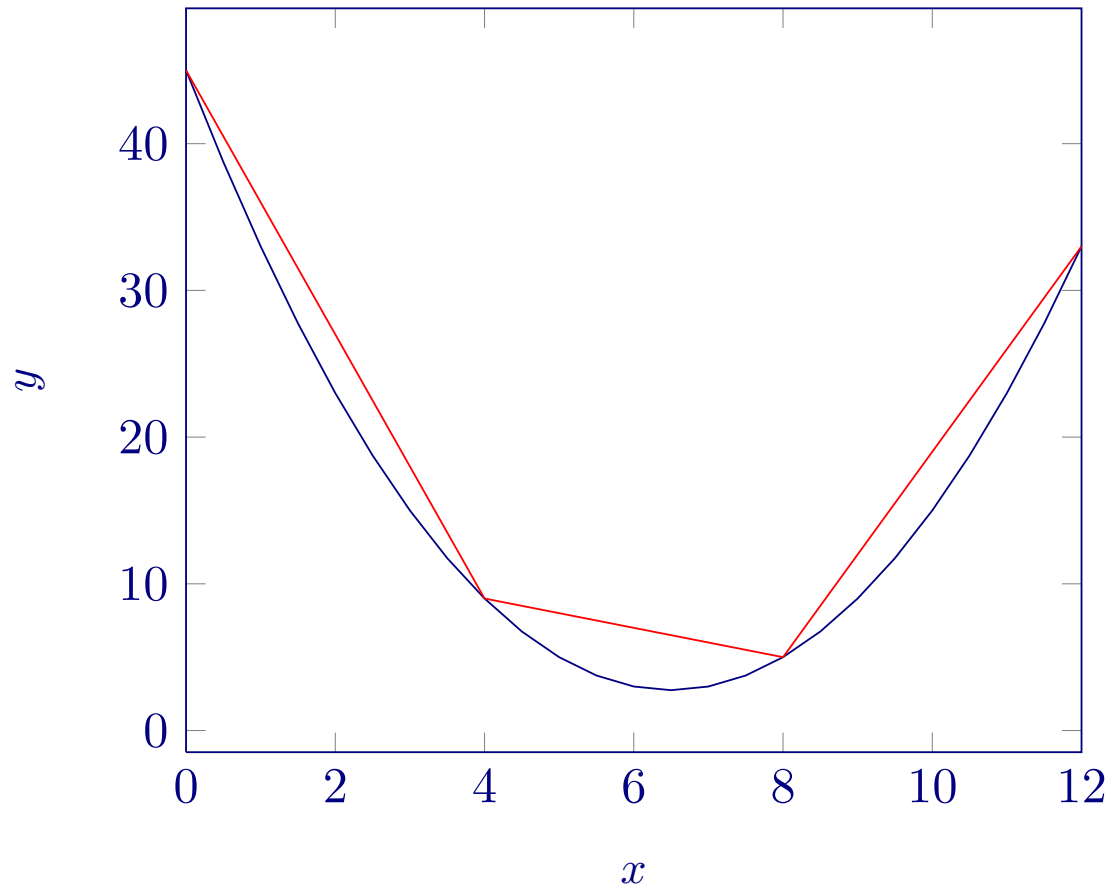
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- ▷ Last constraint can be expressed in integer variables

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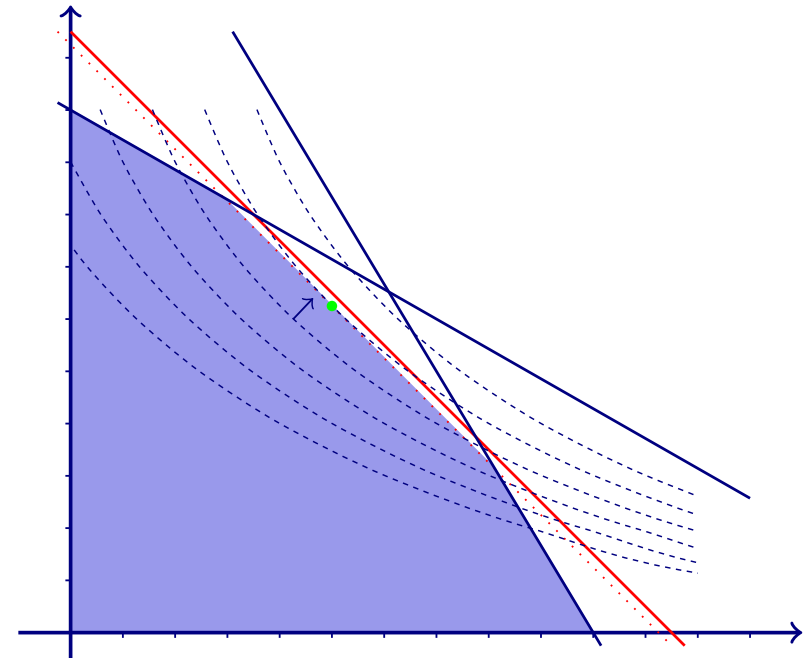
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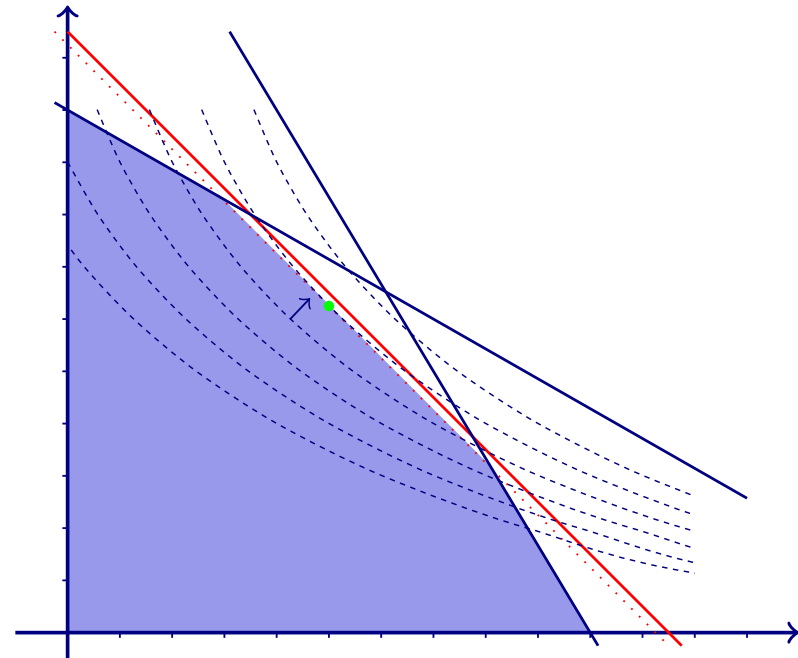
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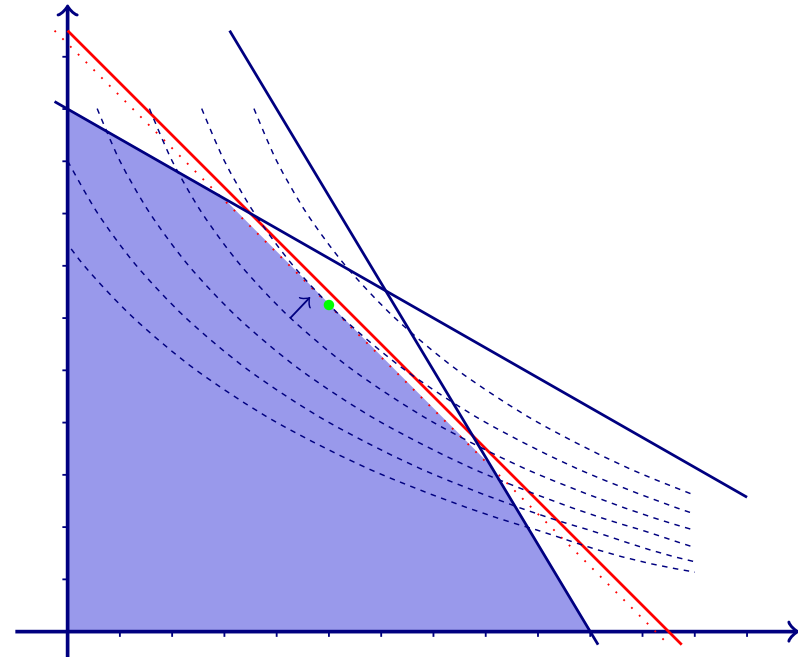
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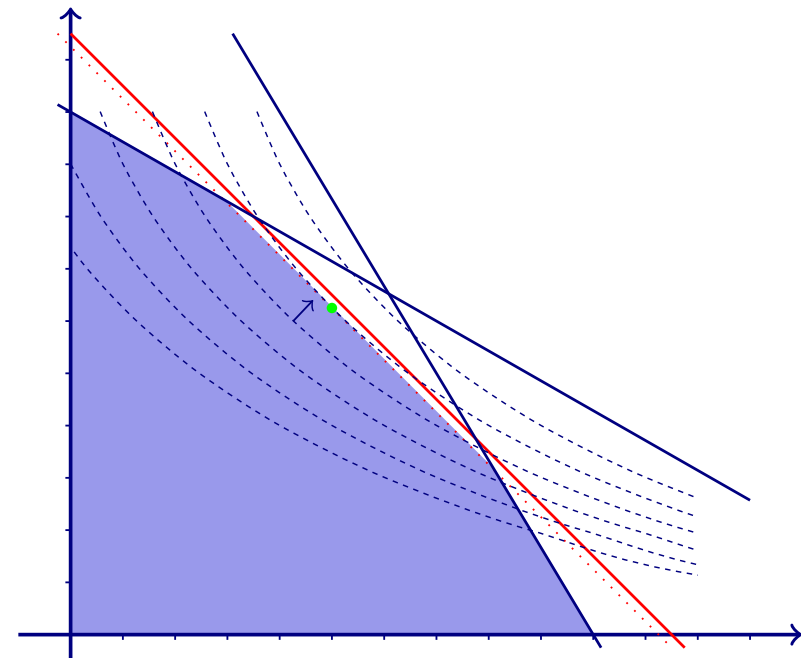
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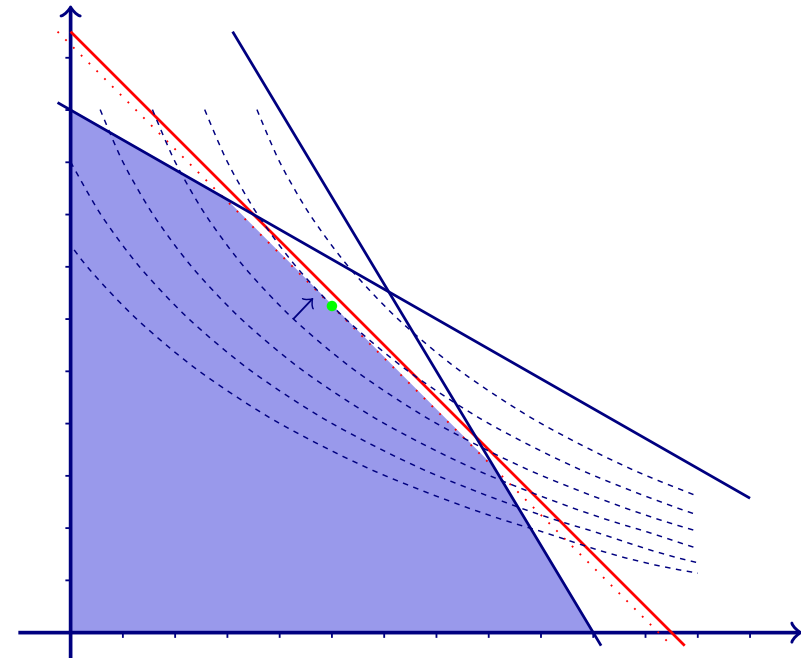
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- ▷ Models, Data and Algorithms
- ▷ Linear Optimization
- ▷ Mathematical Background: Polyhedra, Simplex-Algorithm
- ▷ Sensitivity Analysis; (Mixed) Integer Programming
- ▷ MIP Modelling
- ▷ MIP Modelling: More Examples; Branch & Bound
- ▷ Cutting Planes; Combinatorial Optimization: Examples, Graphs, Algorithms
- ▷ TSP-Heuristics
- ▷ Network Flows
- ▷ Shortest Path Problem
- ▷ Complexity Theory
- ▷ Nonlinear Optimization
- ▷ Scheduling (Jan 25)
- ▷ Lot Sizing (Feb 01)
- ▷ Summary (Feb 08)
- ▷ Oral exam (Feb 15)